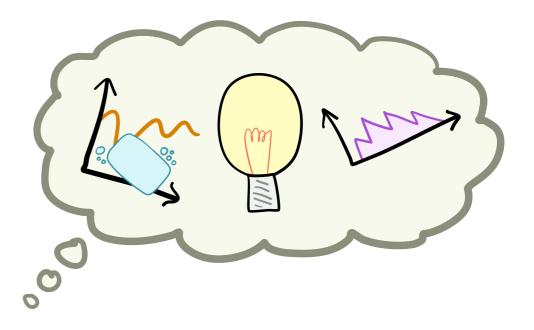
Near-Optimal Scheduling:

Towards a Unified Theory

Ziv Scully
Carnegie Mellon University



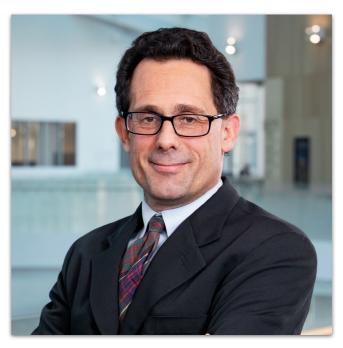
Collaborators



Mor Harchol-Balter (CMU)



Isaac Grosof (CMU)



Alan Scheller-Wolf (CMU)

Adam Wierman (Caltech)
Onno Boxma (TU/e)
Jan-Pieter Dorsman (UvA)
Lucas van Kreveld (UvA)

-Queueing system: jobs waiting for service

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File servers

- *Jobs*: file requests
- Service: load and send contents

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- *Jobs*: SQL queries
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- Service: run on a CPU core

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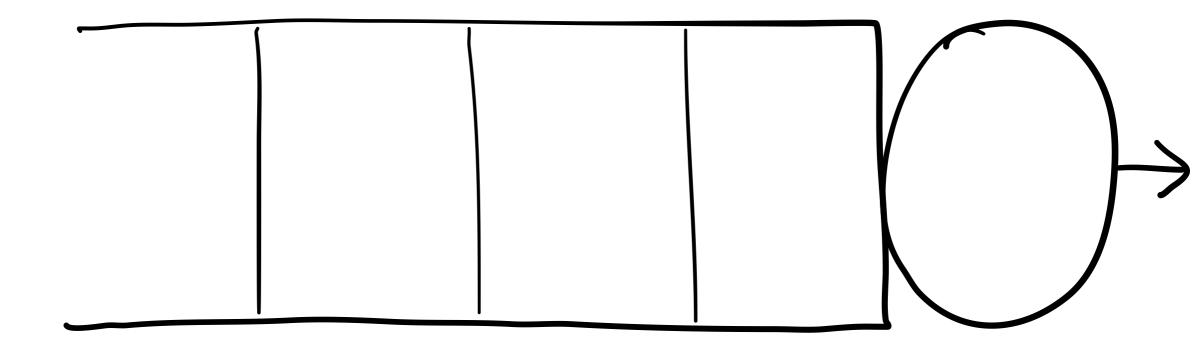
Operating systems

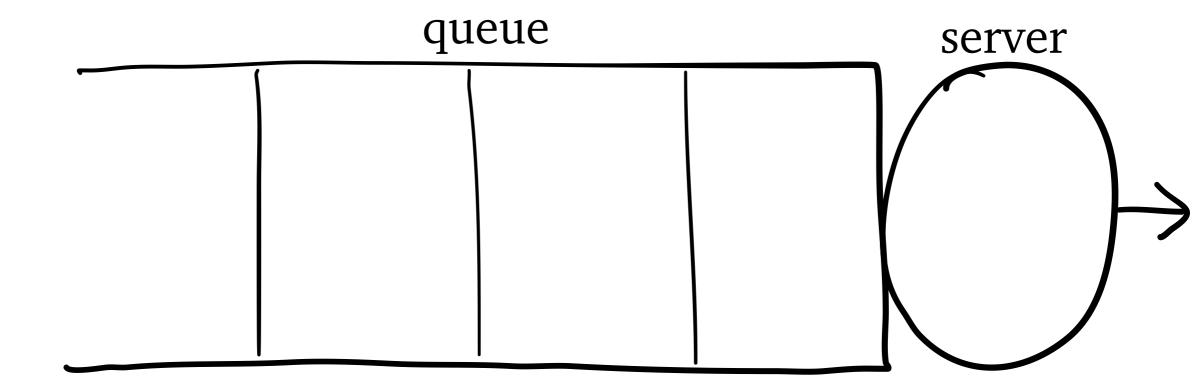
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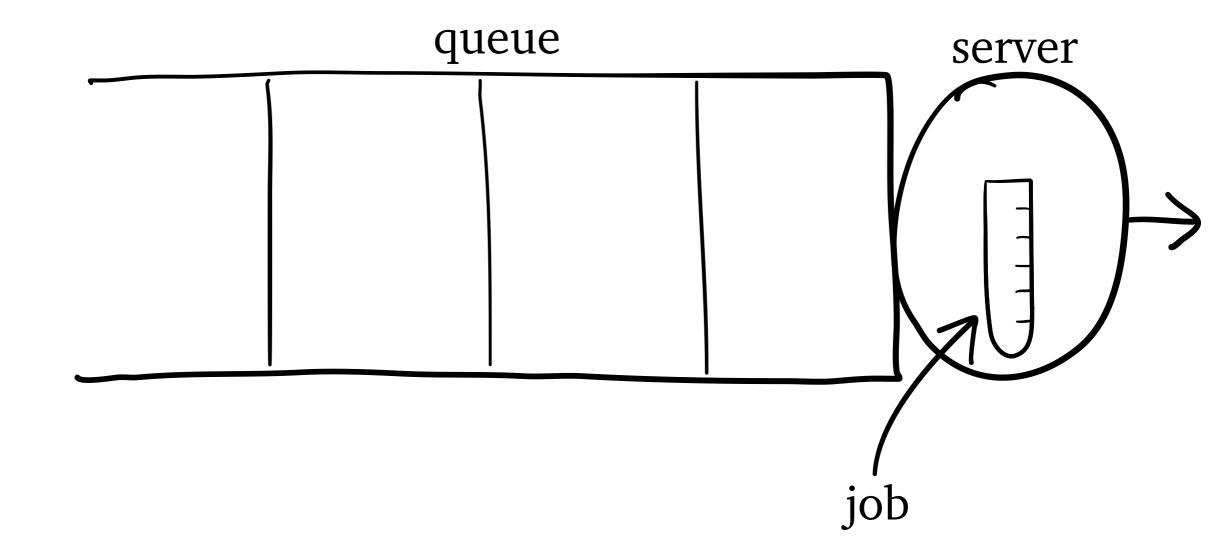
Queueing theory: studies the mathematical essence of queueing systems

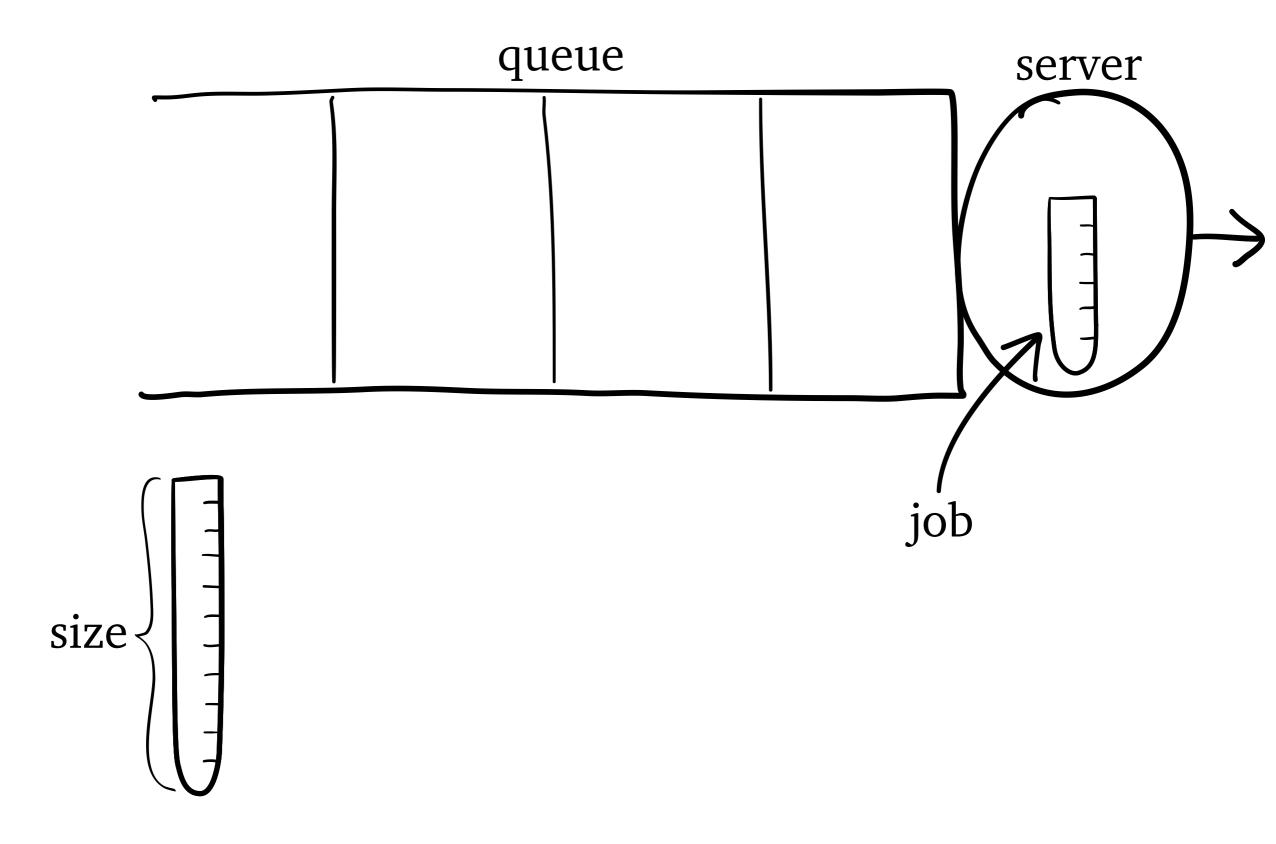


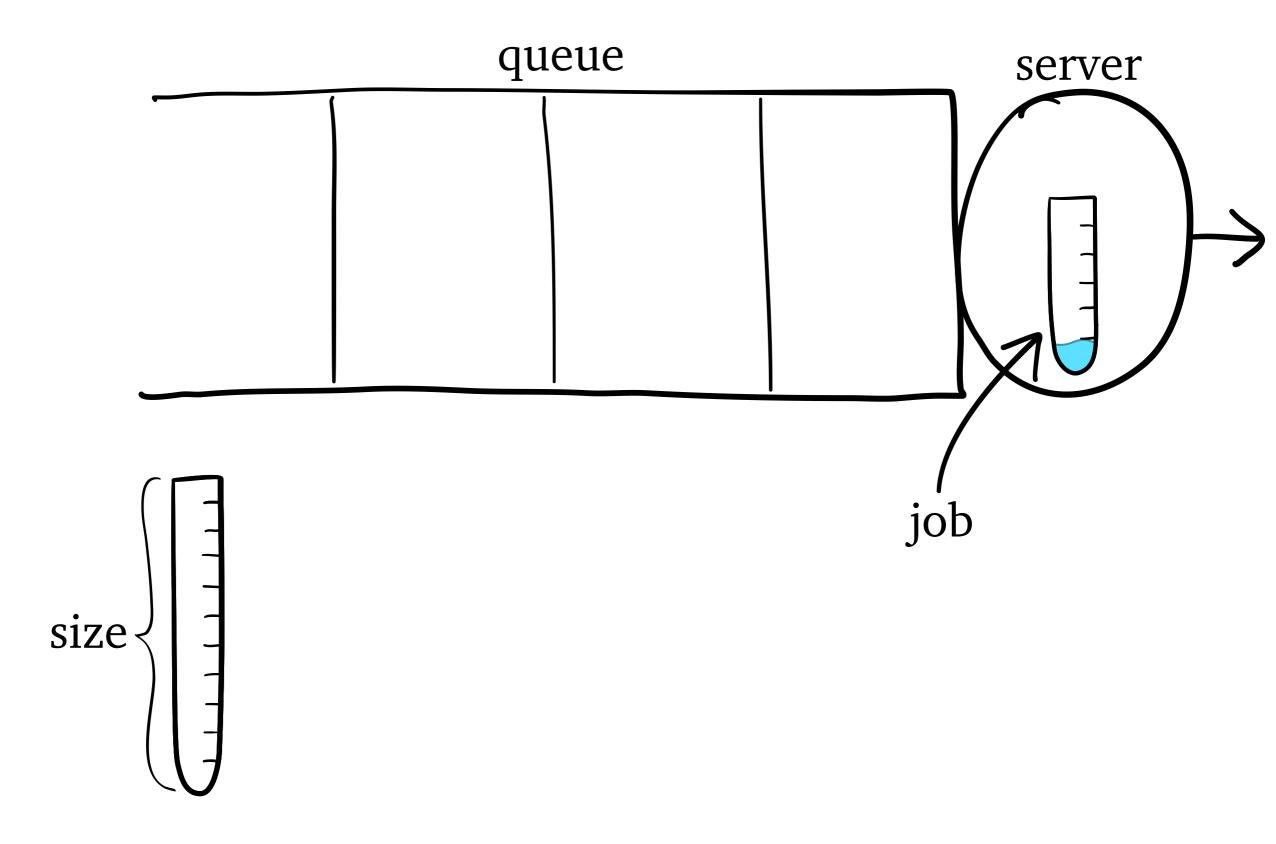


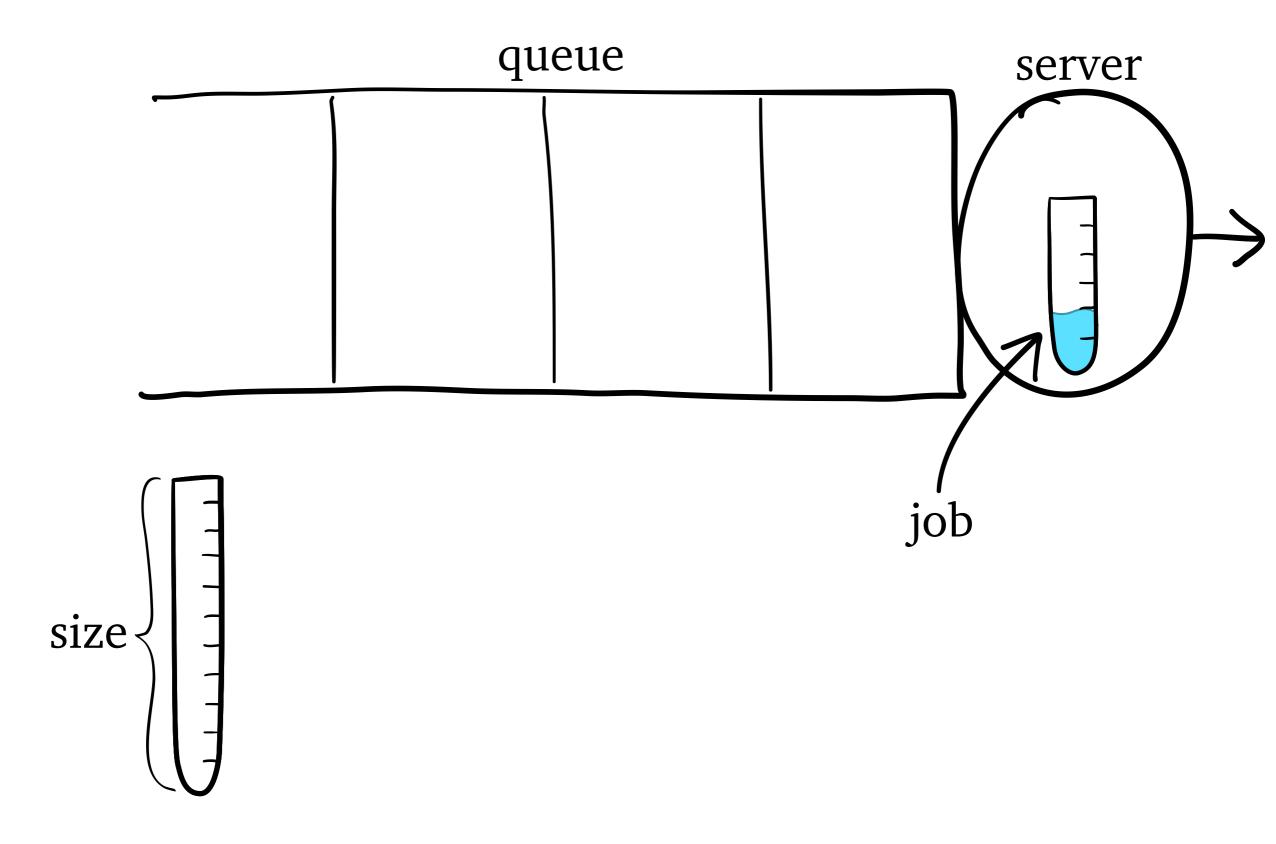


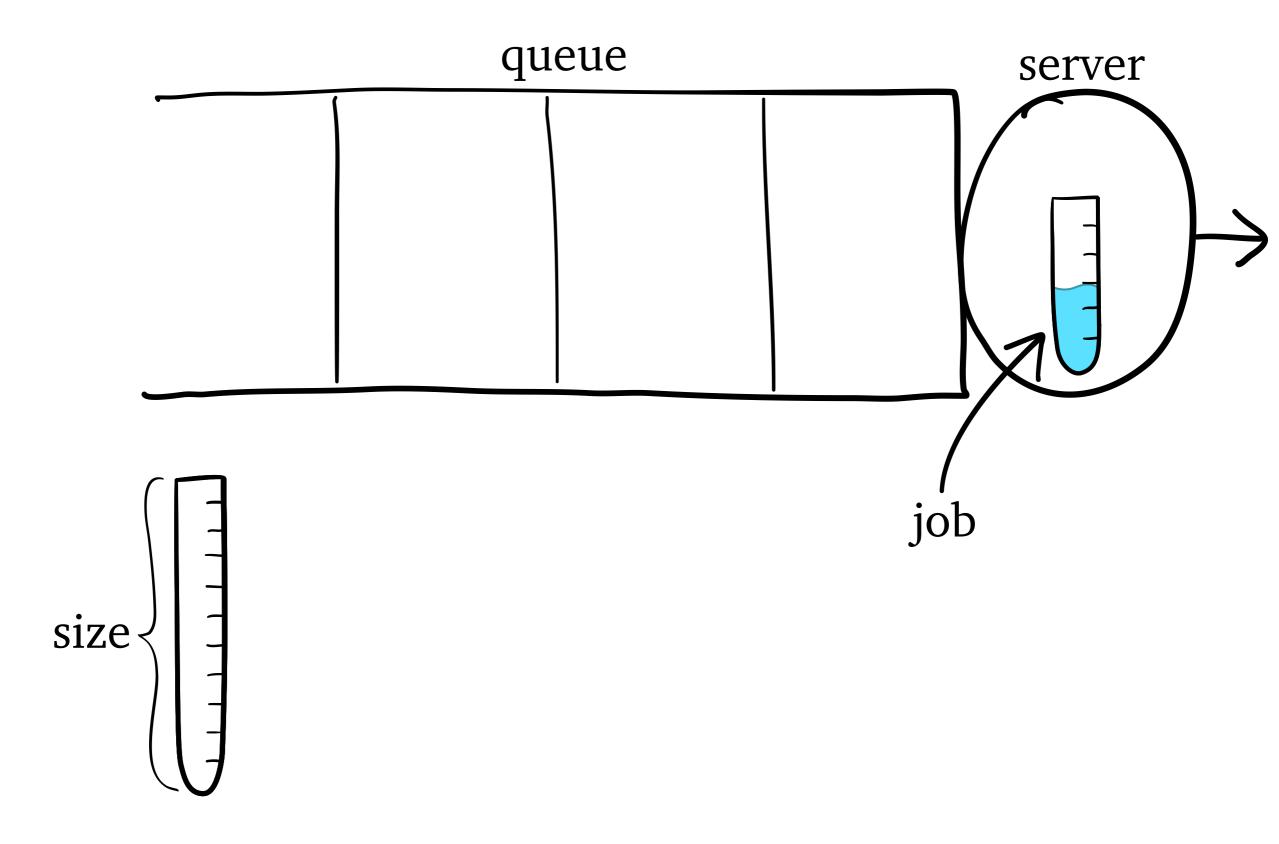


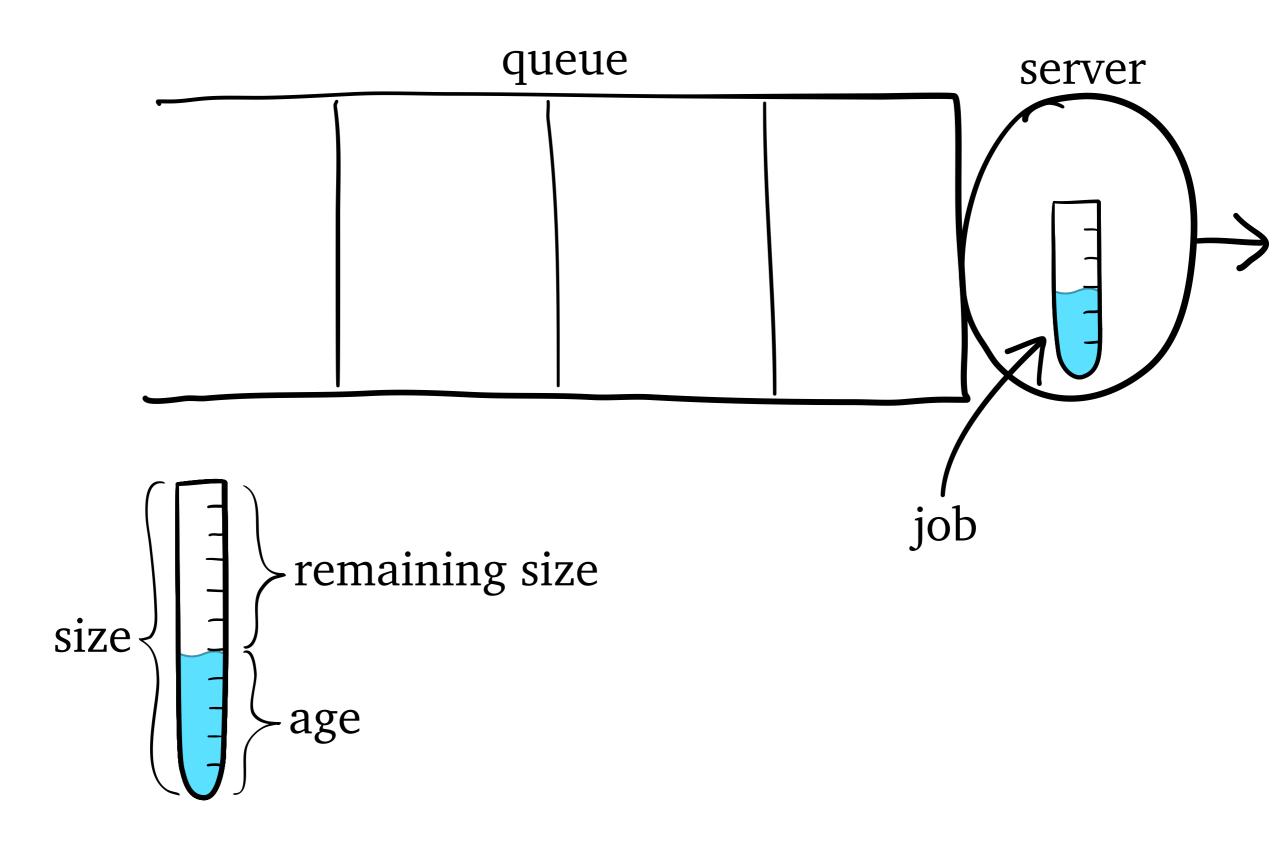


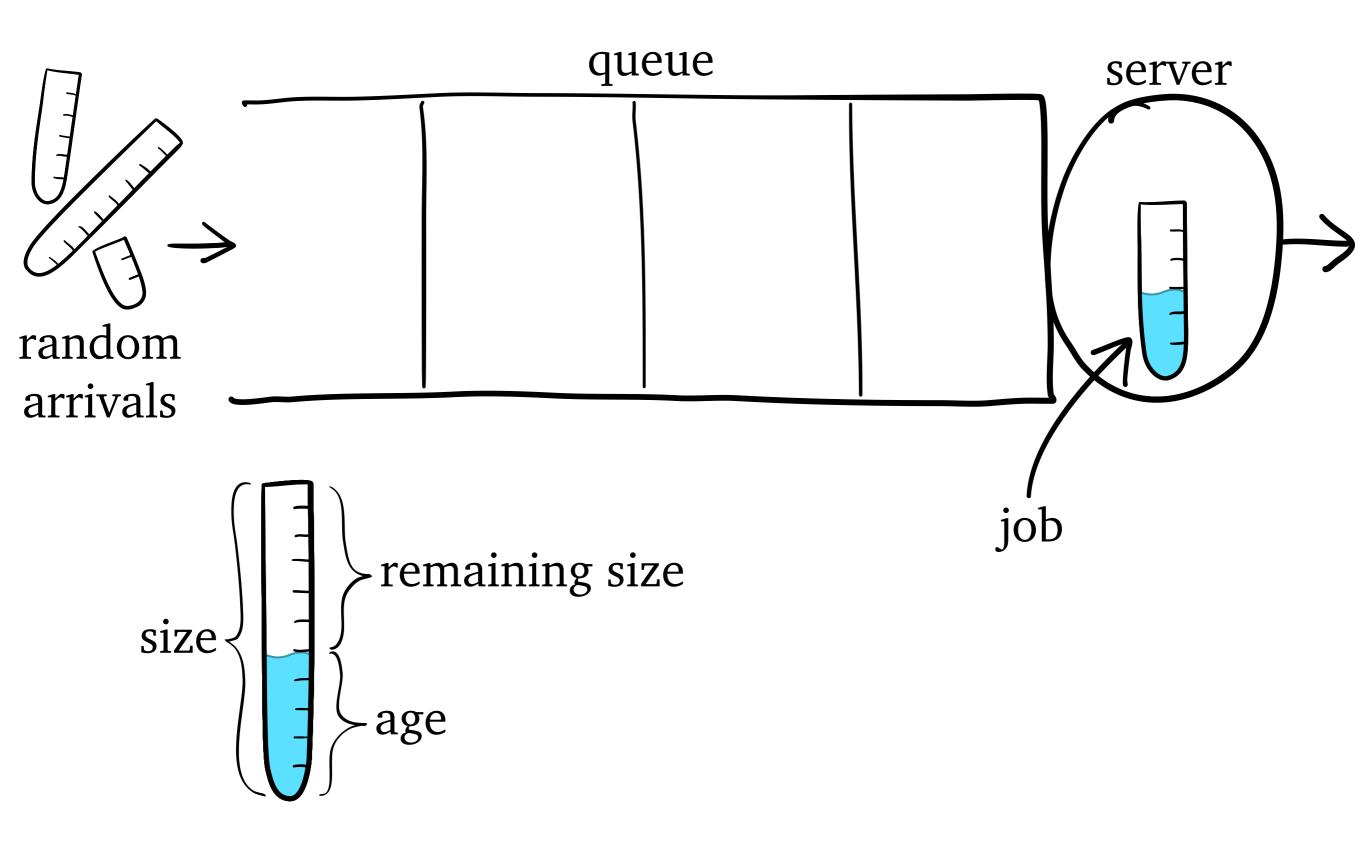


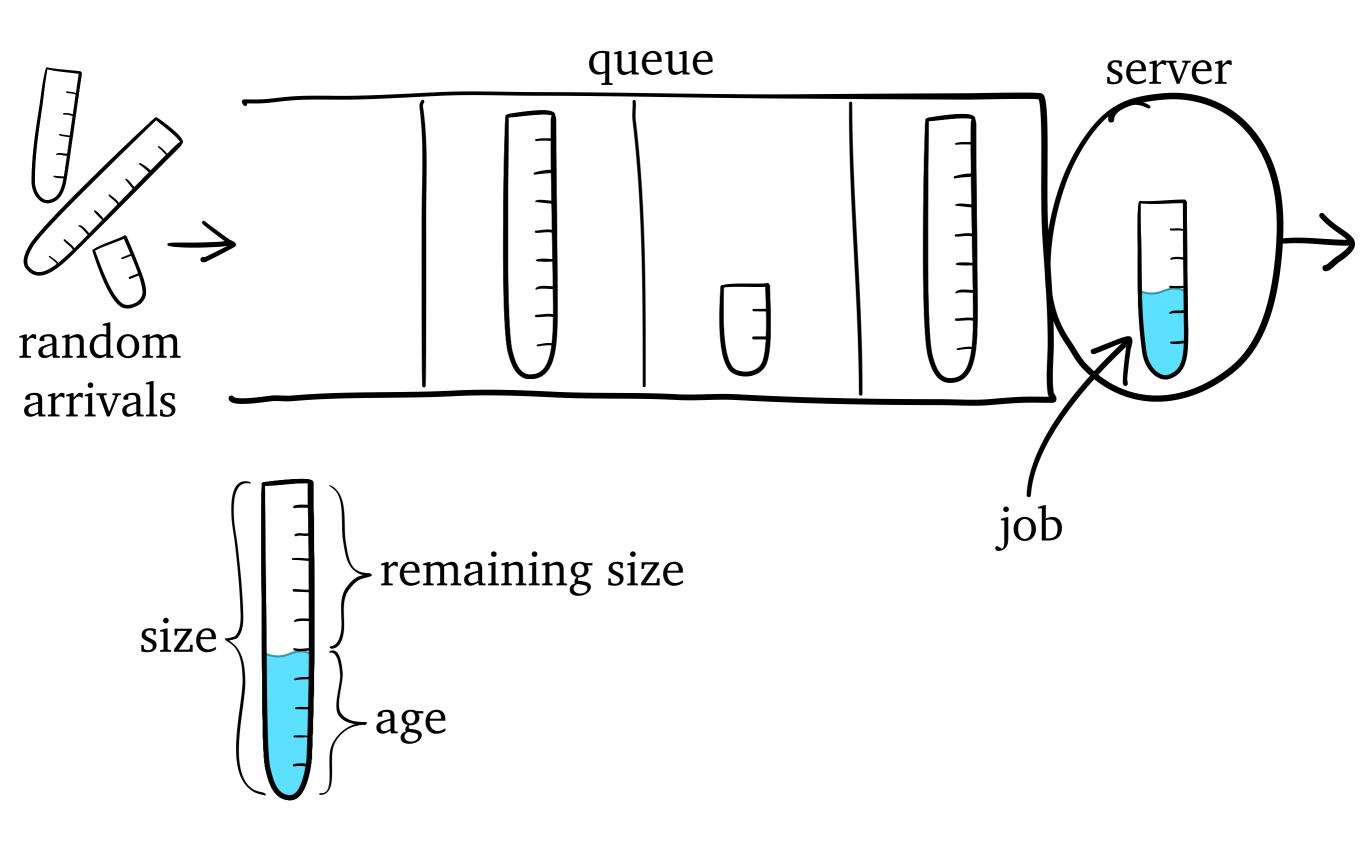


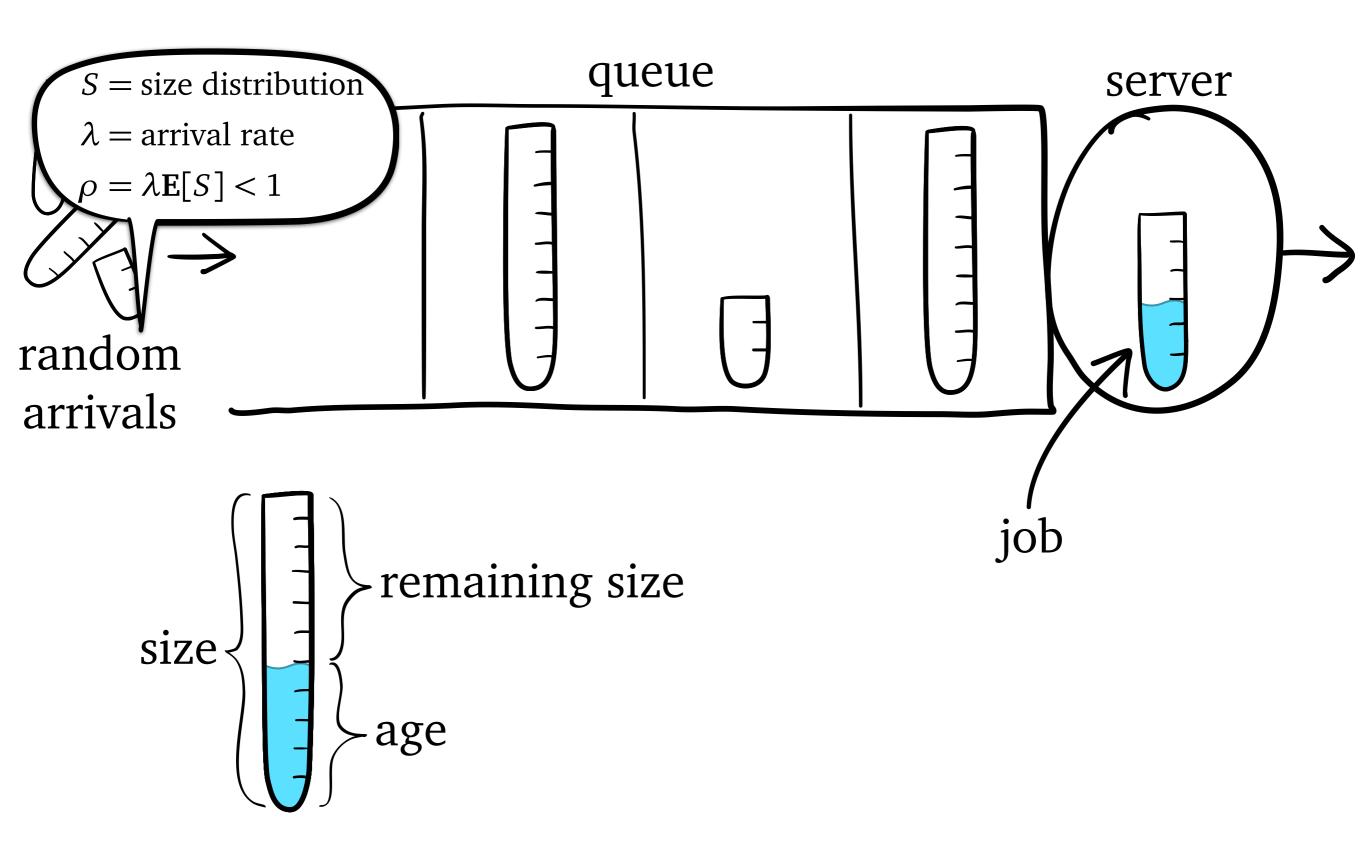


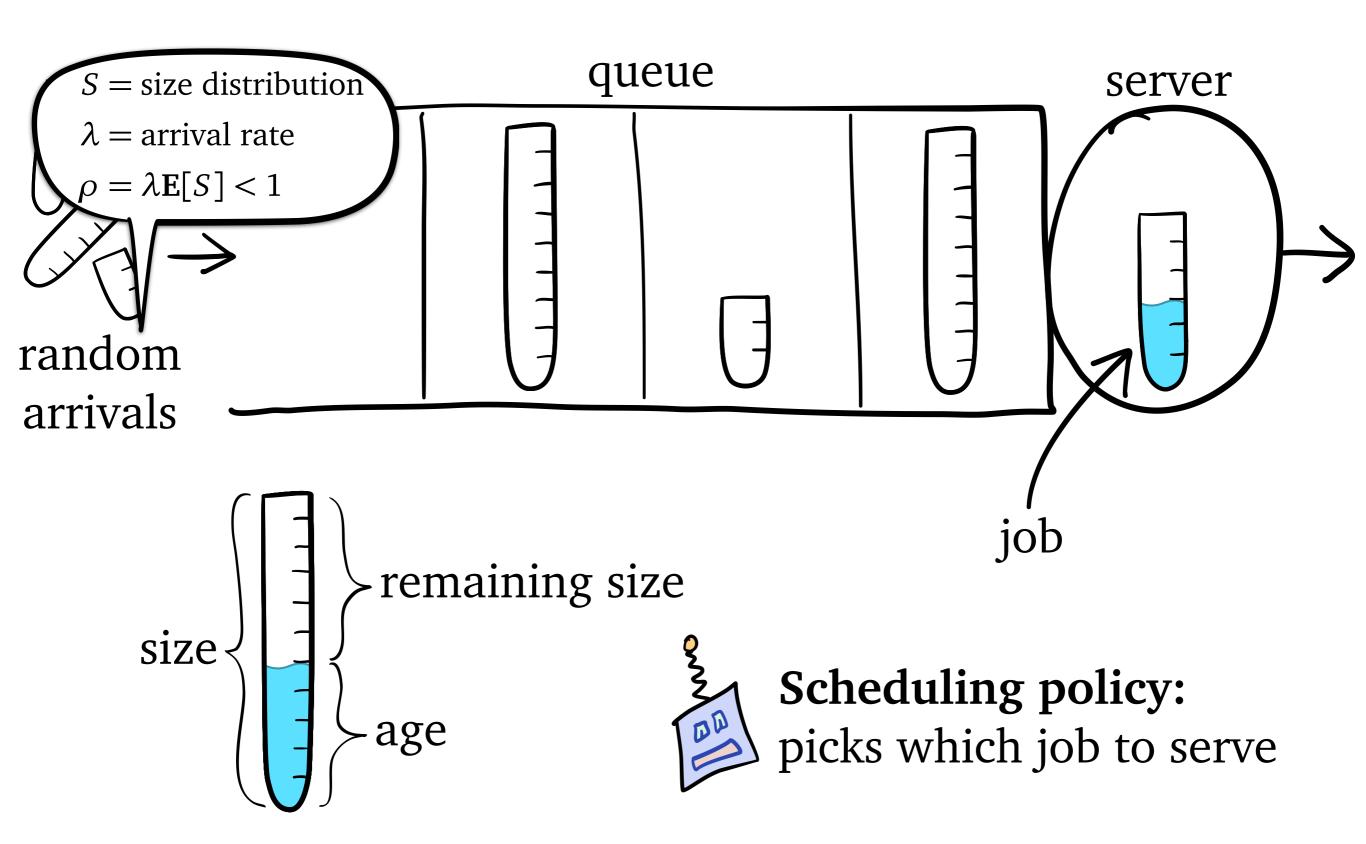


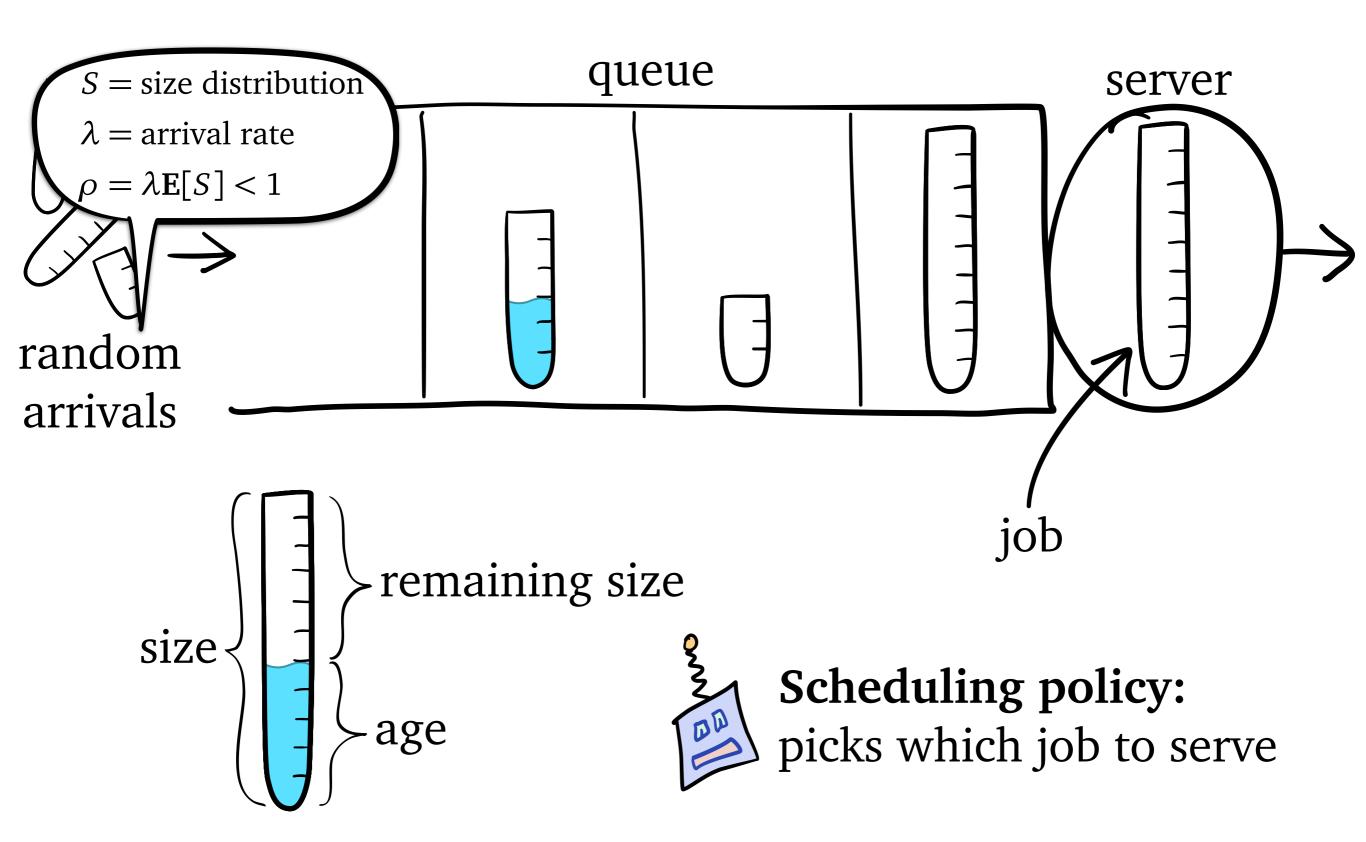


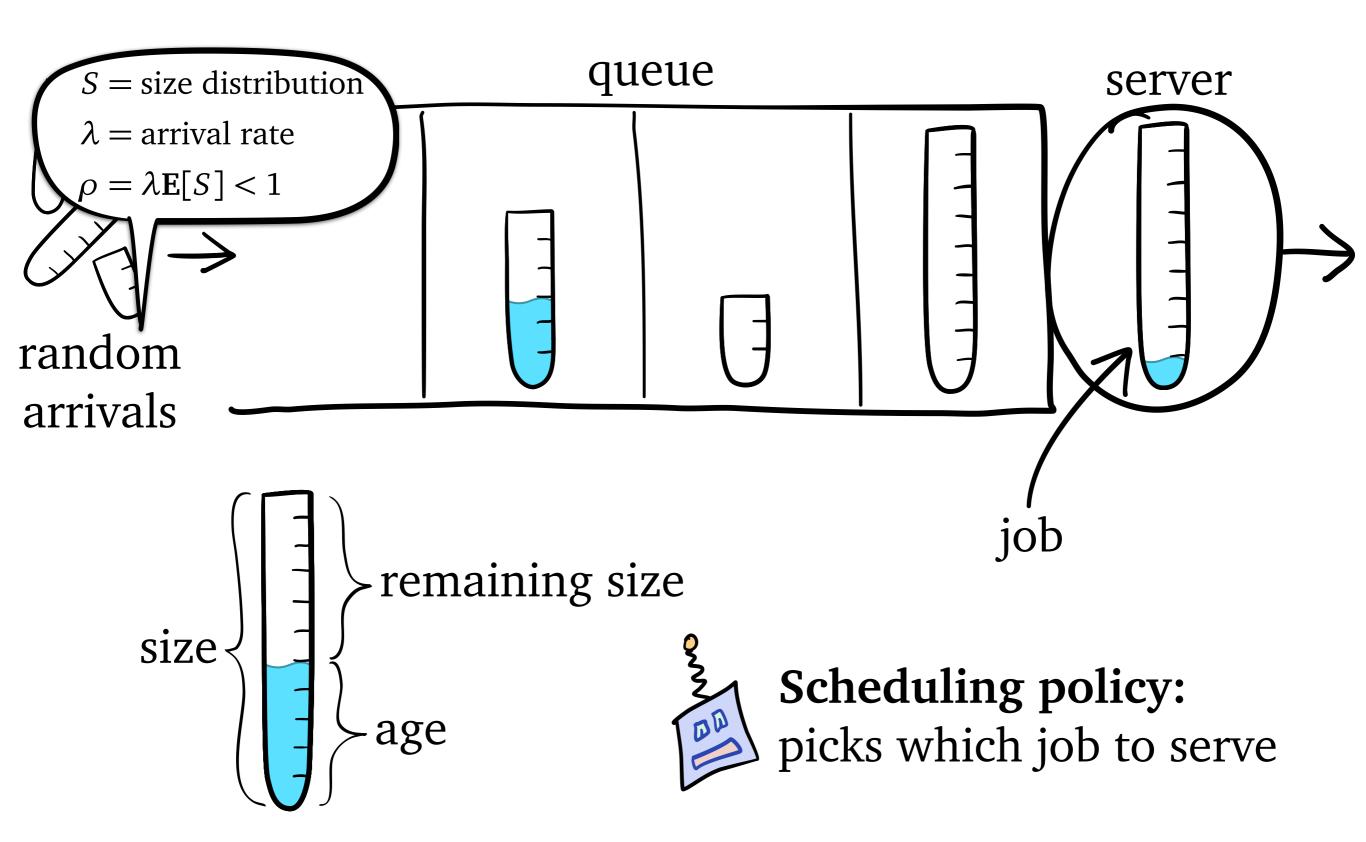


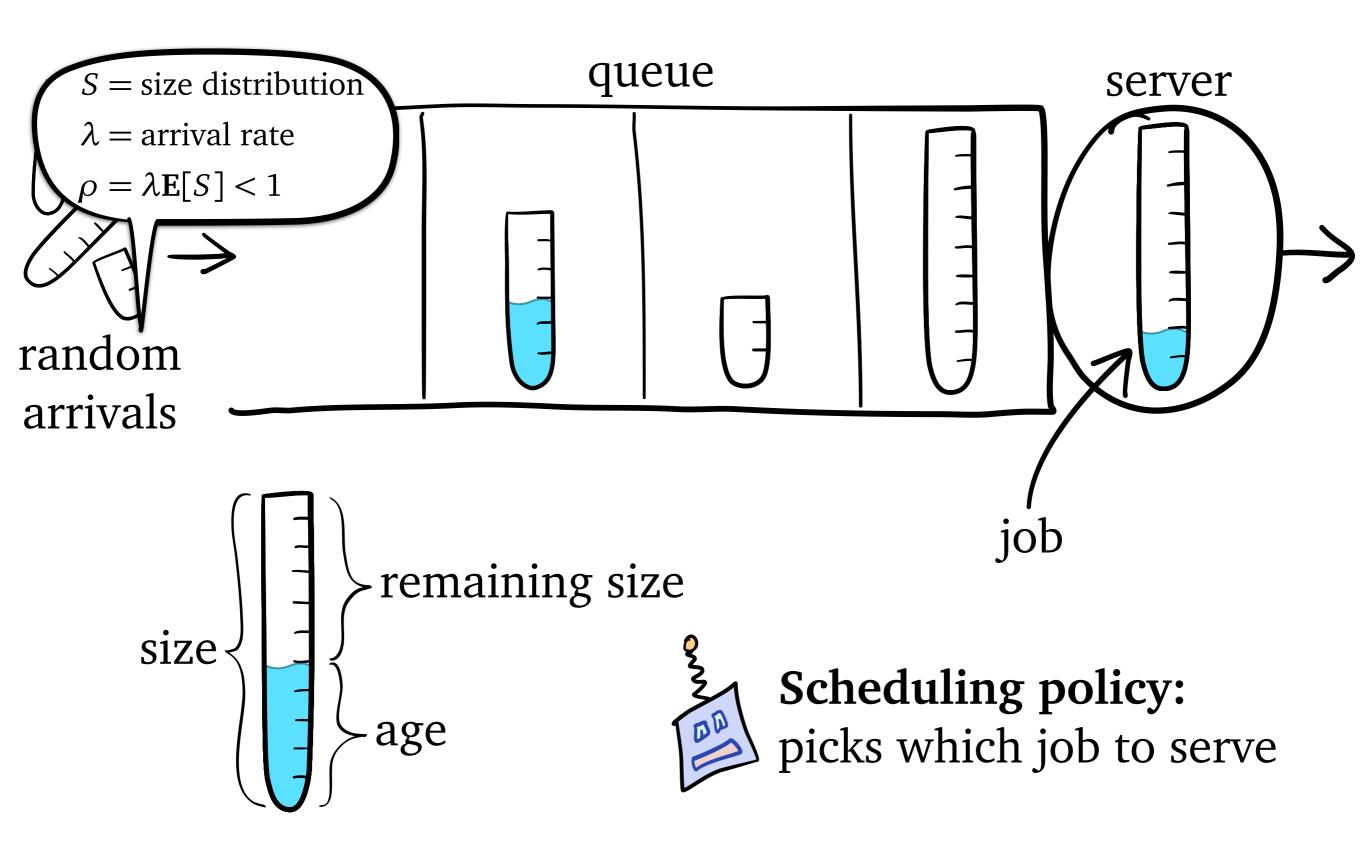


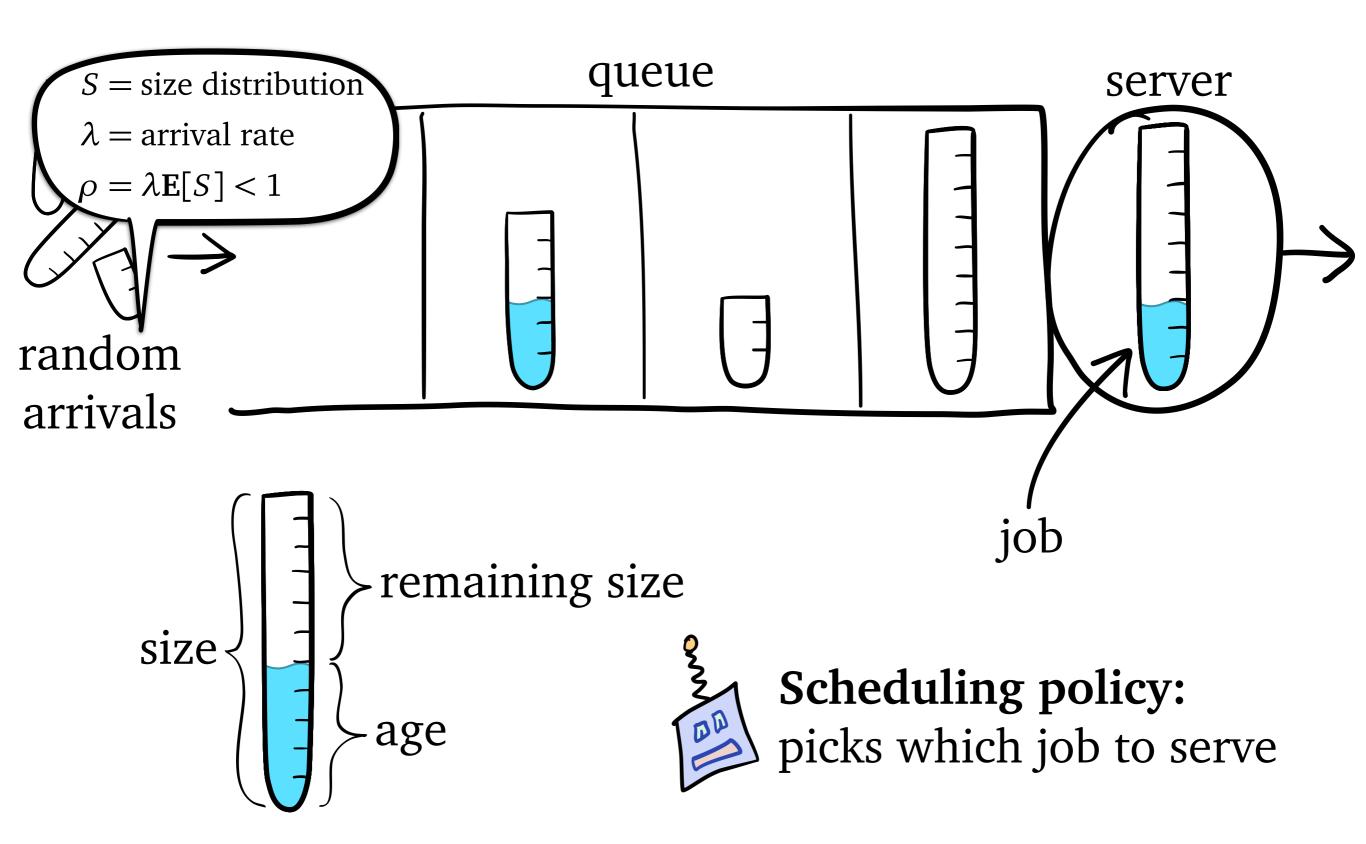


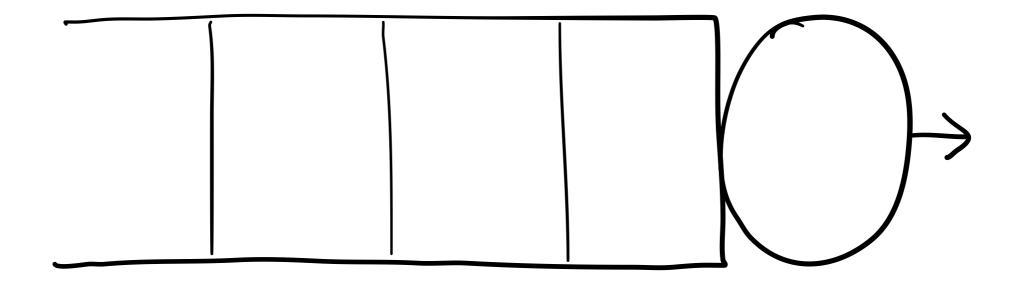


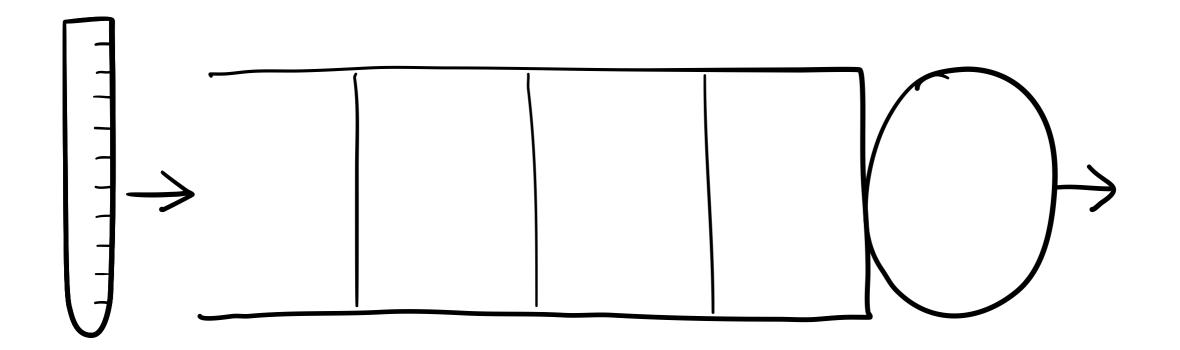


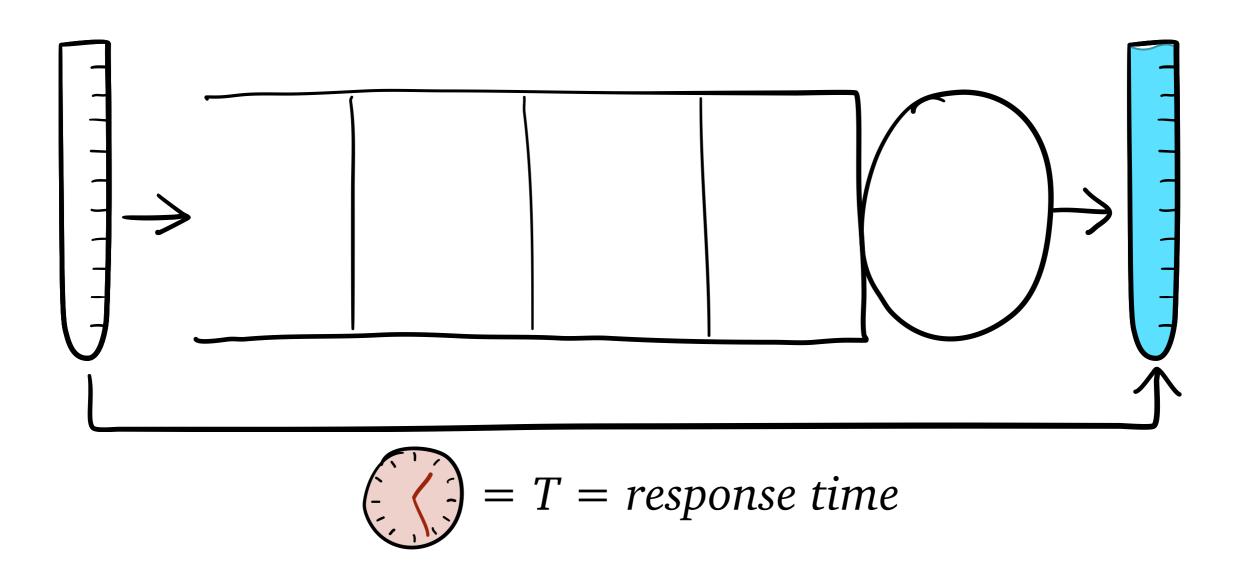


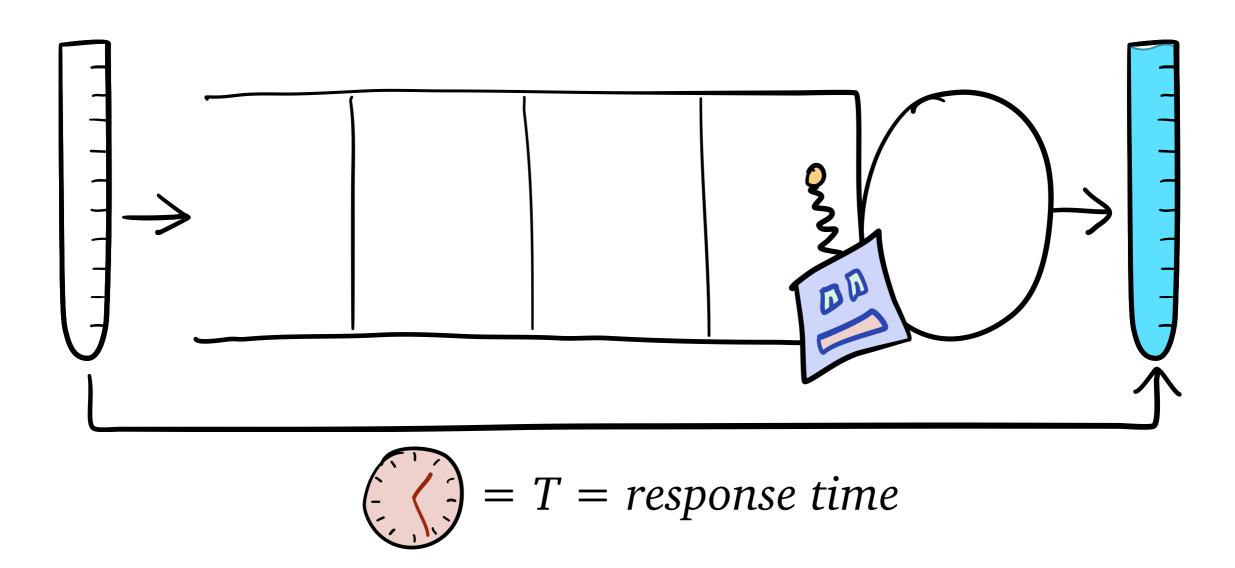


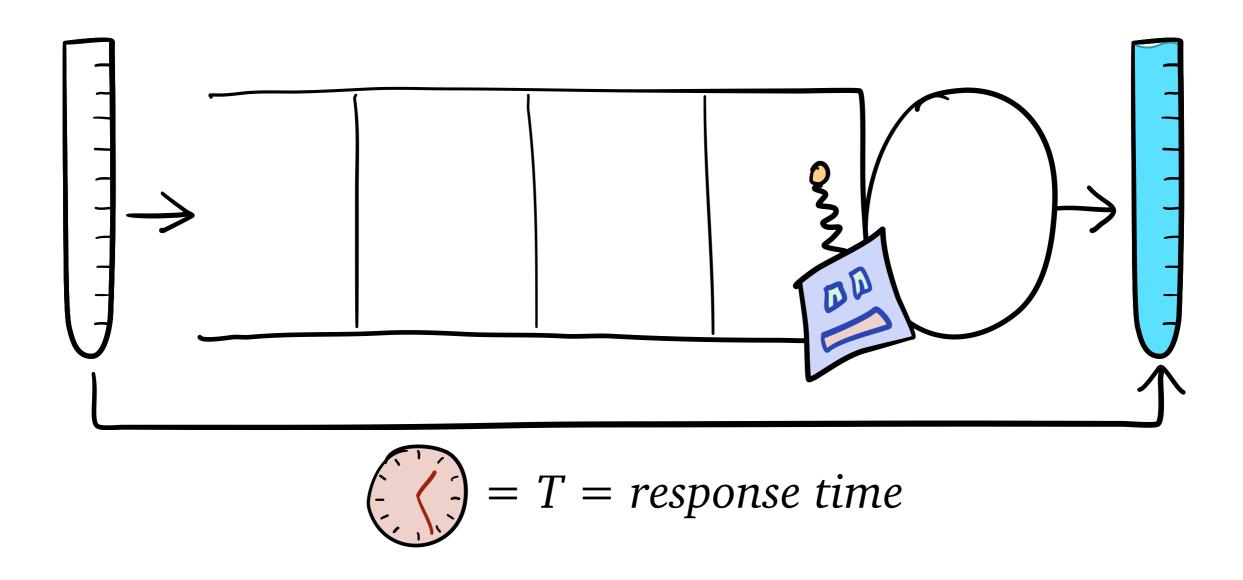




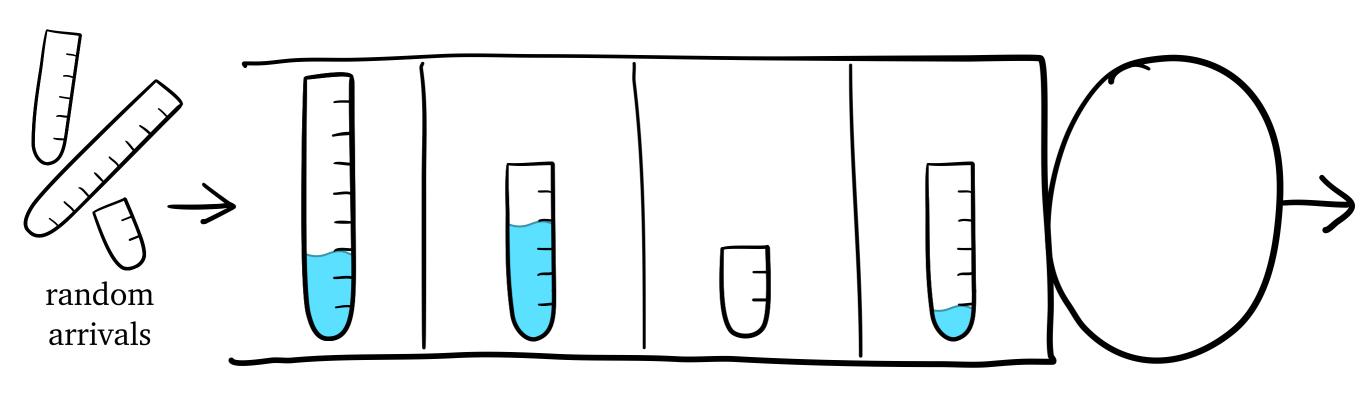


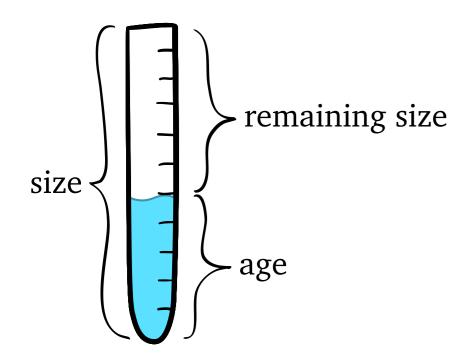


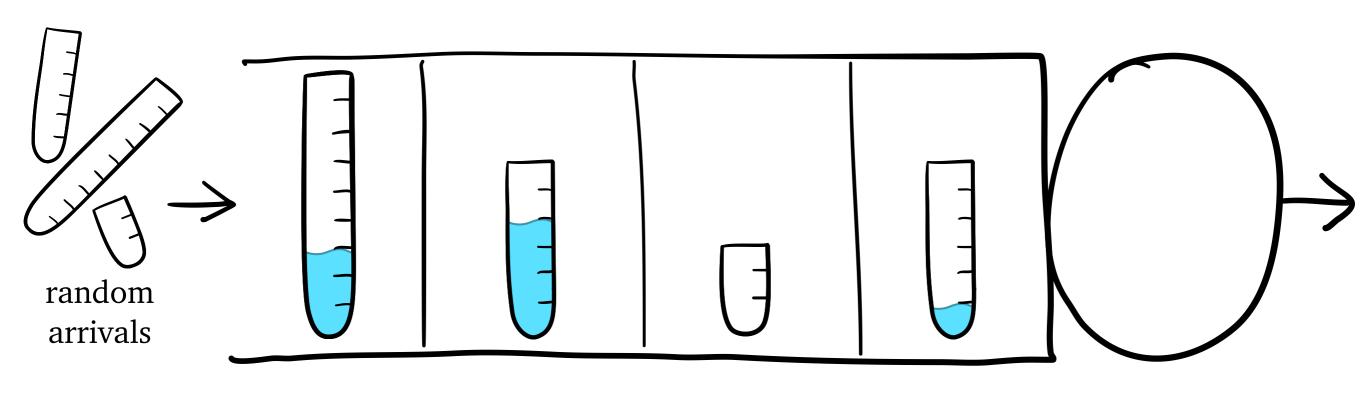


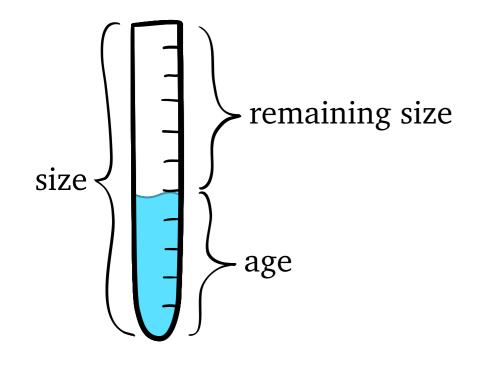


Goal: schedule to minimize *mean* $response\ time\ \mathbf{E}[T]$ and other metrics



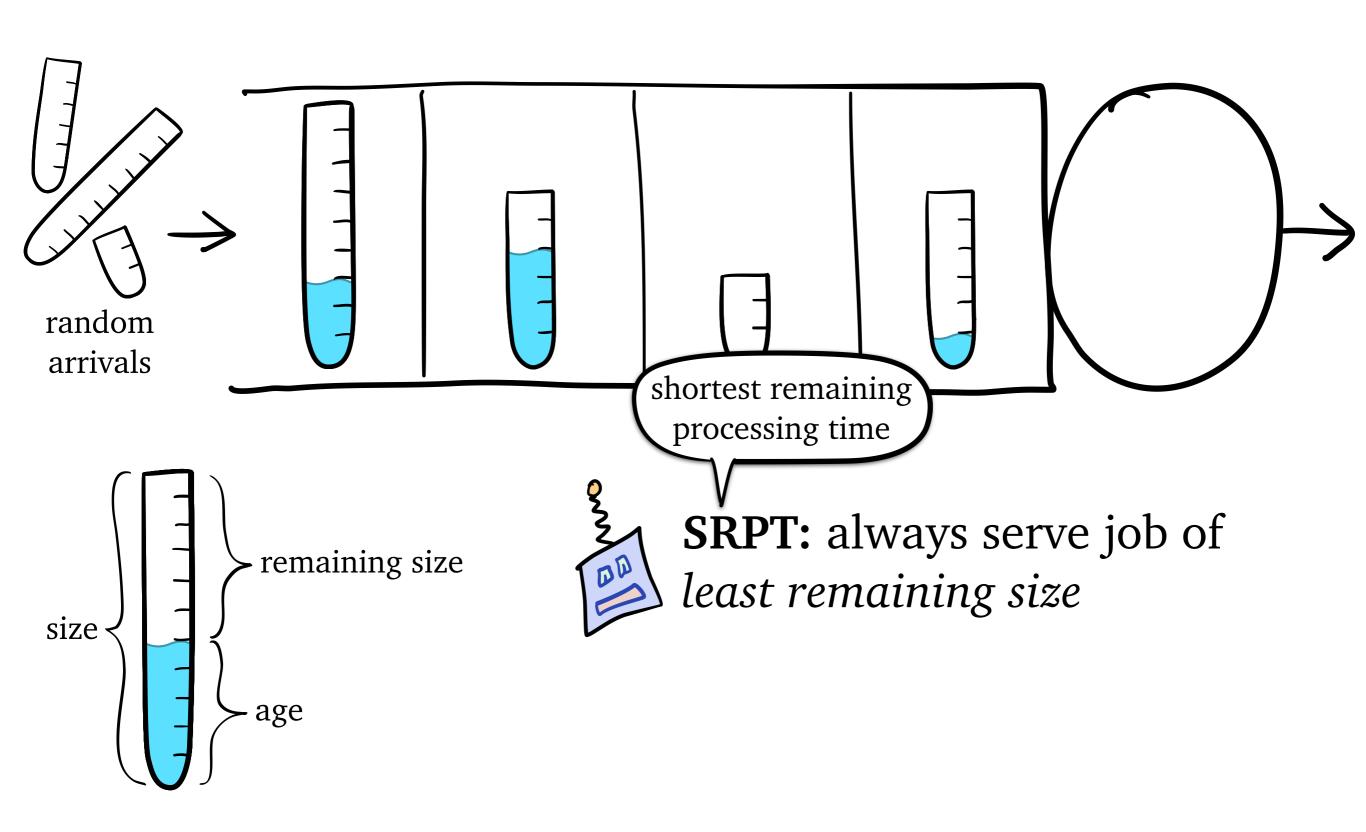


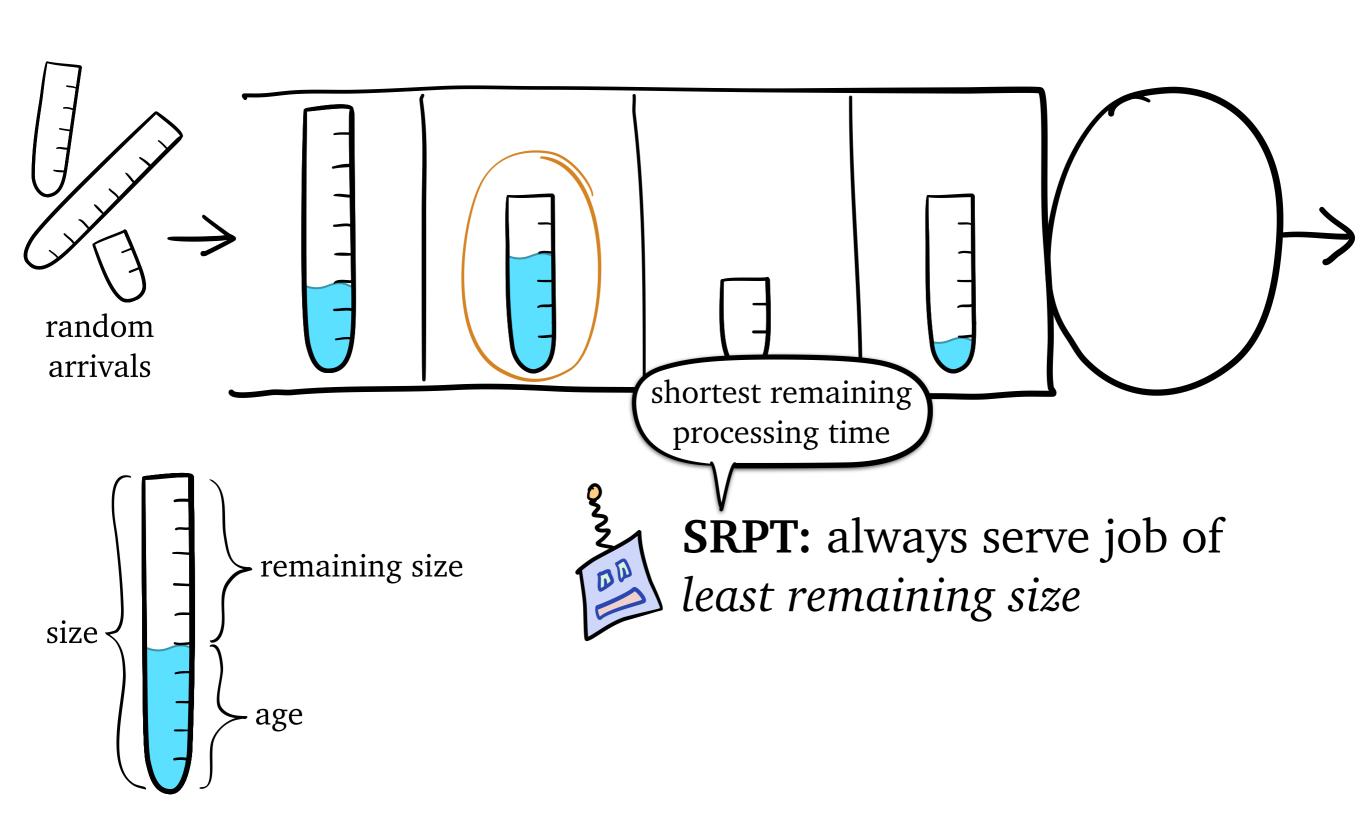


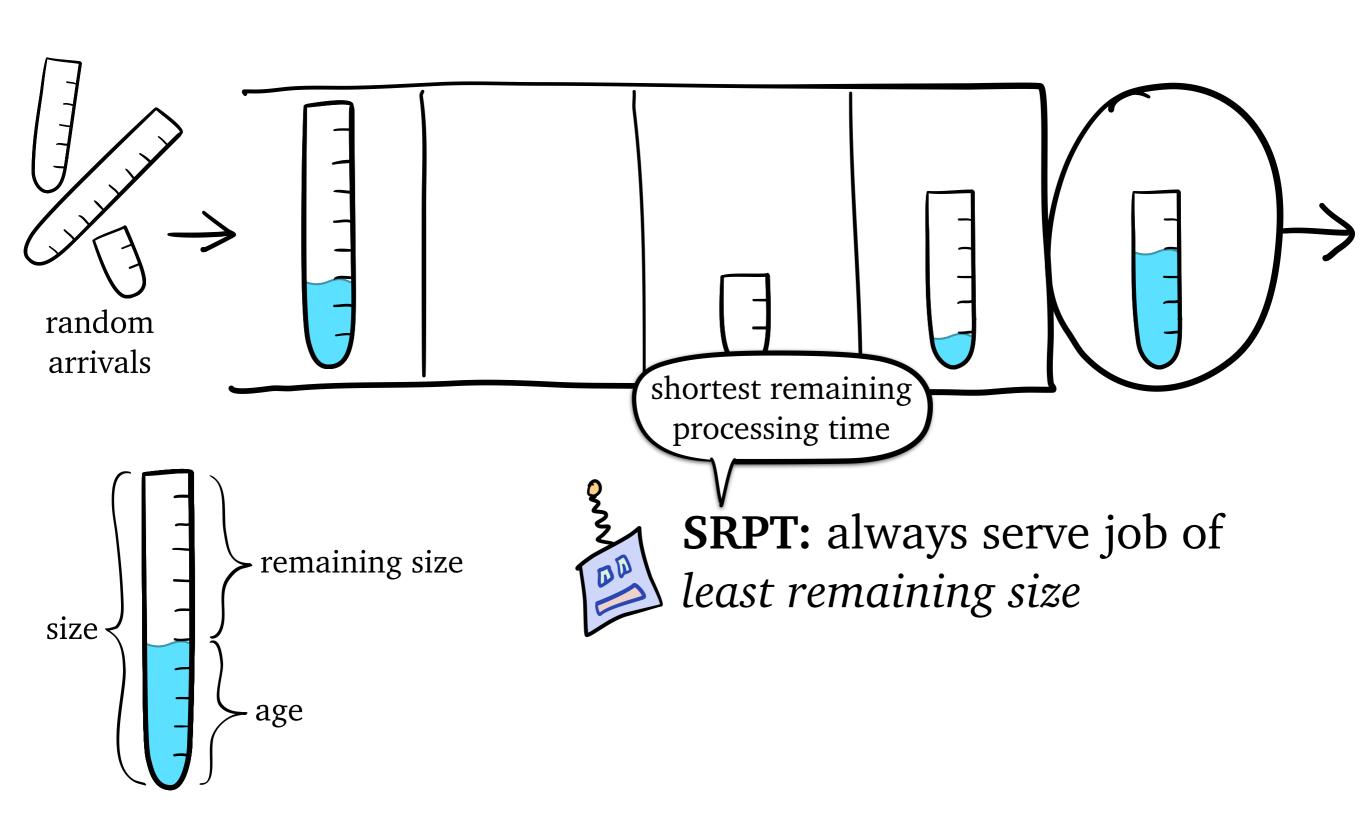




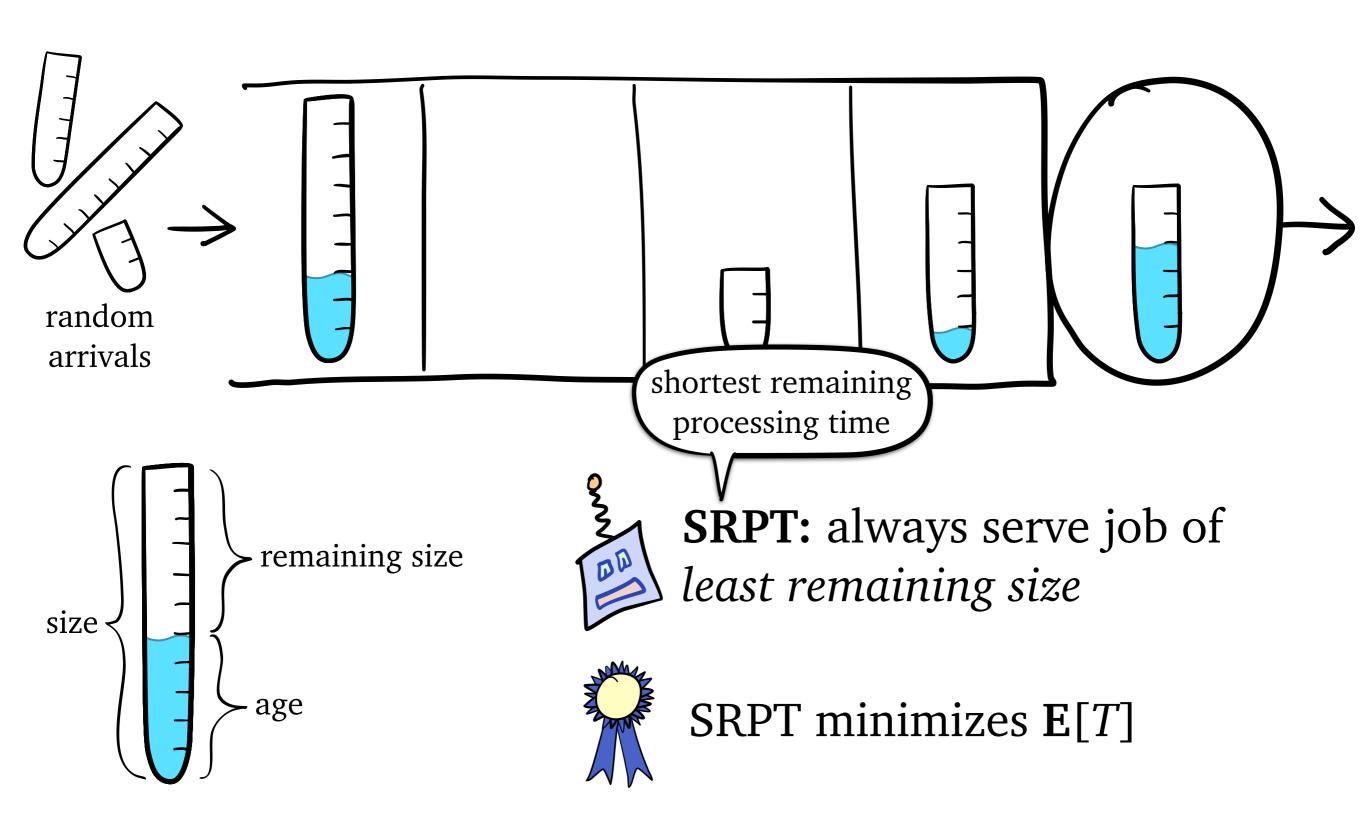
SRPT: always serve job of least remaining size

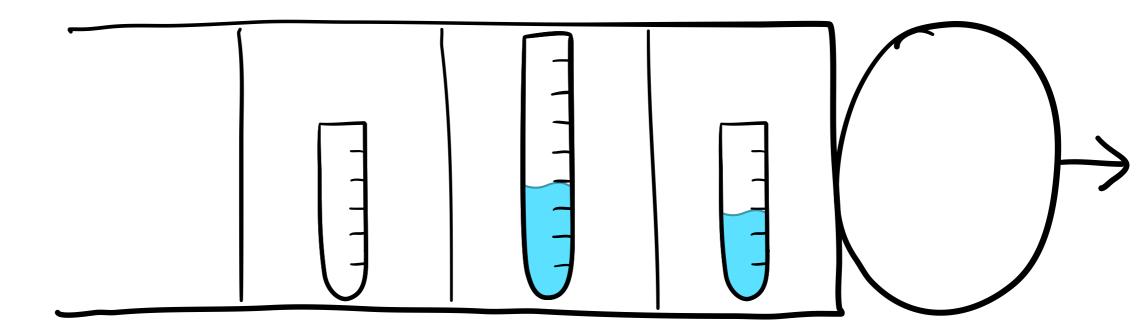


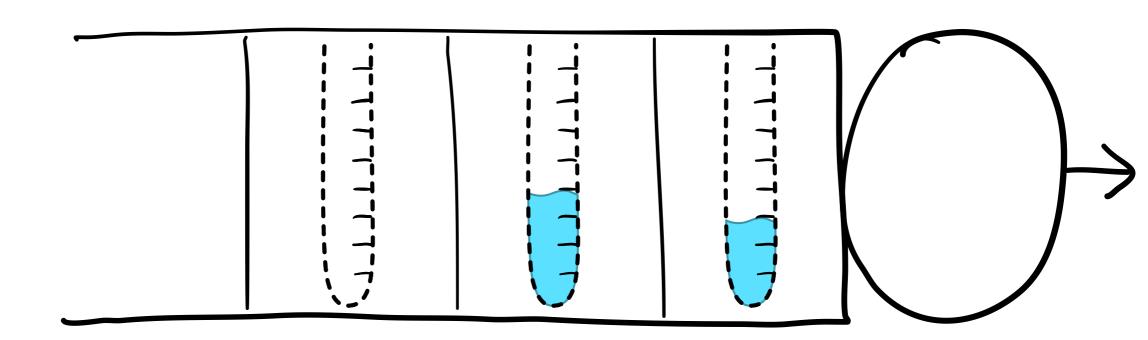


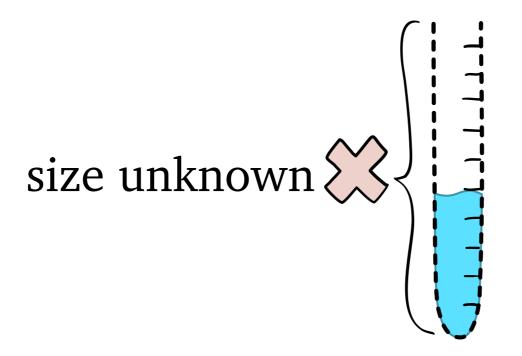


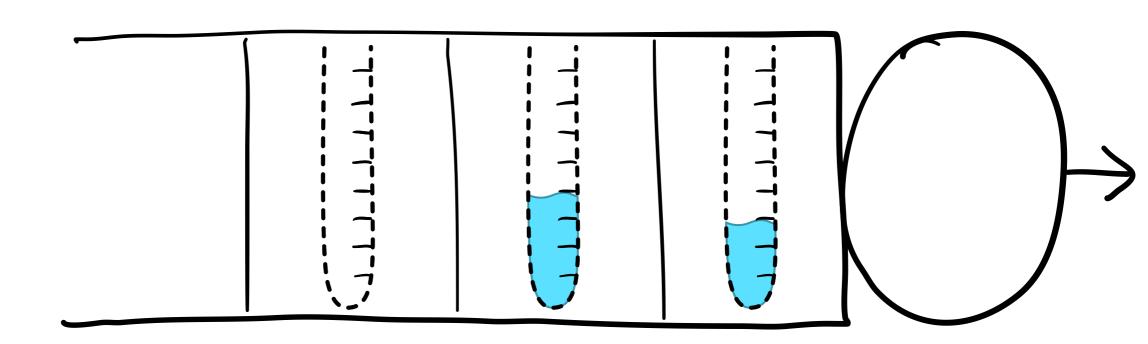
How to Schedule?

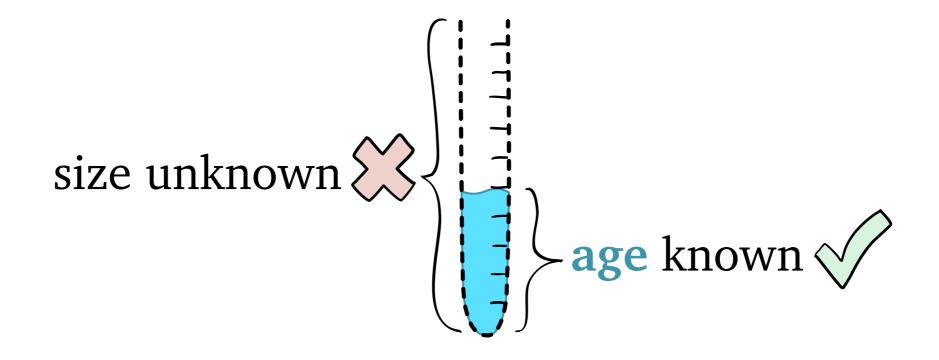


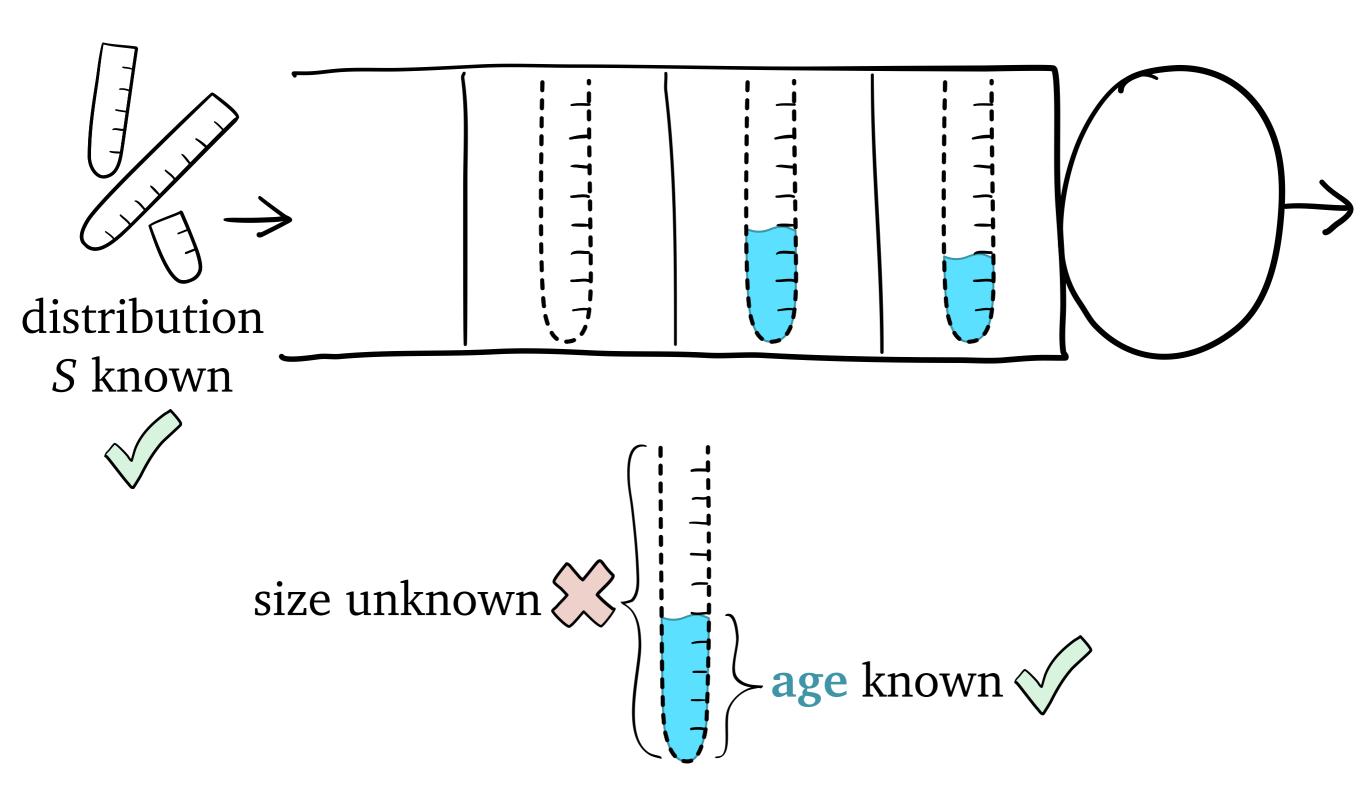


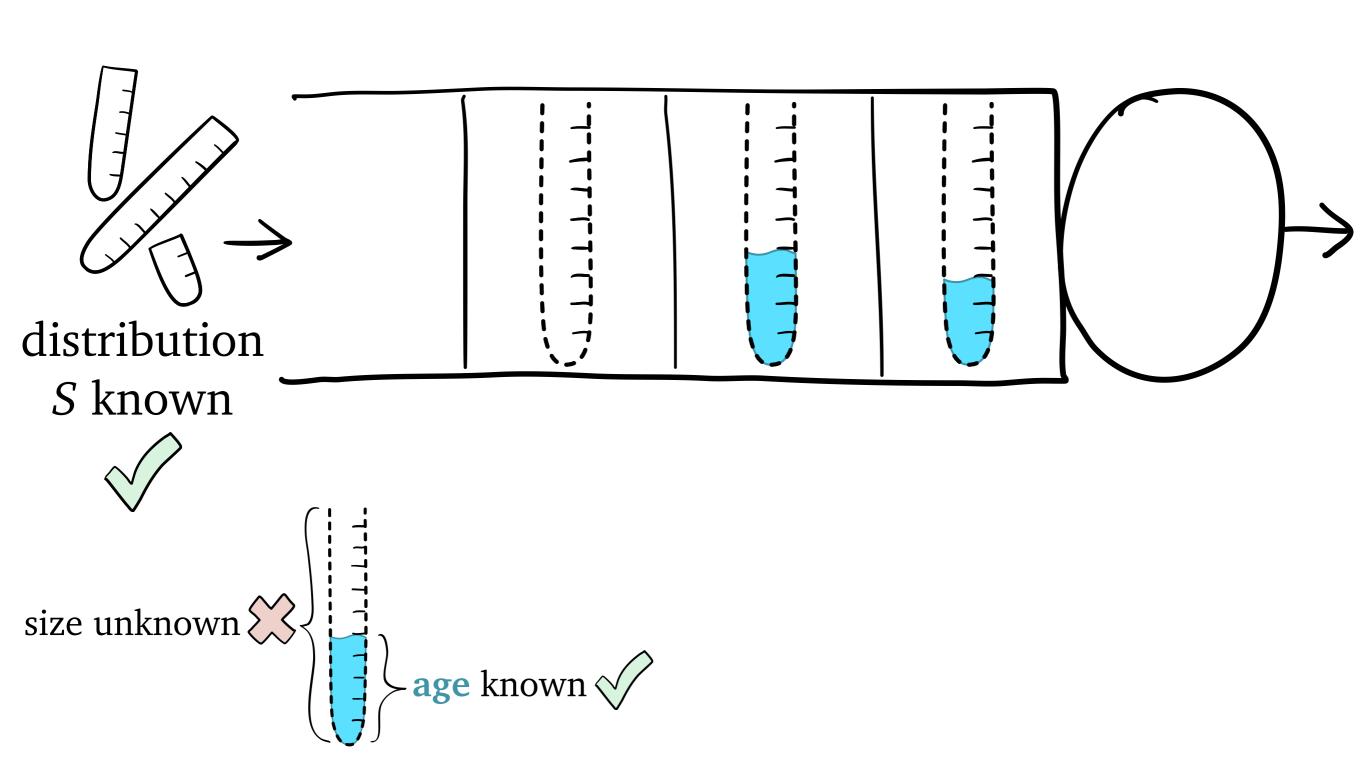


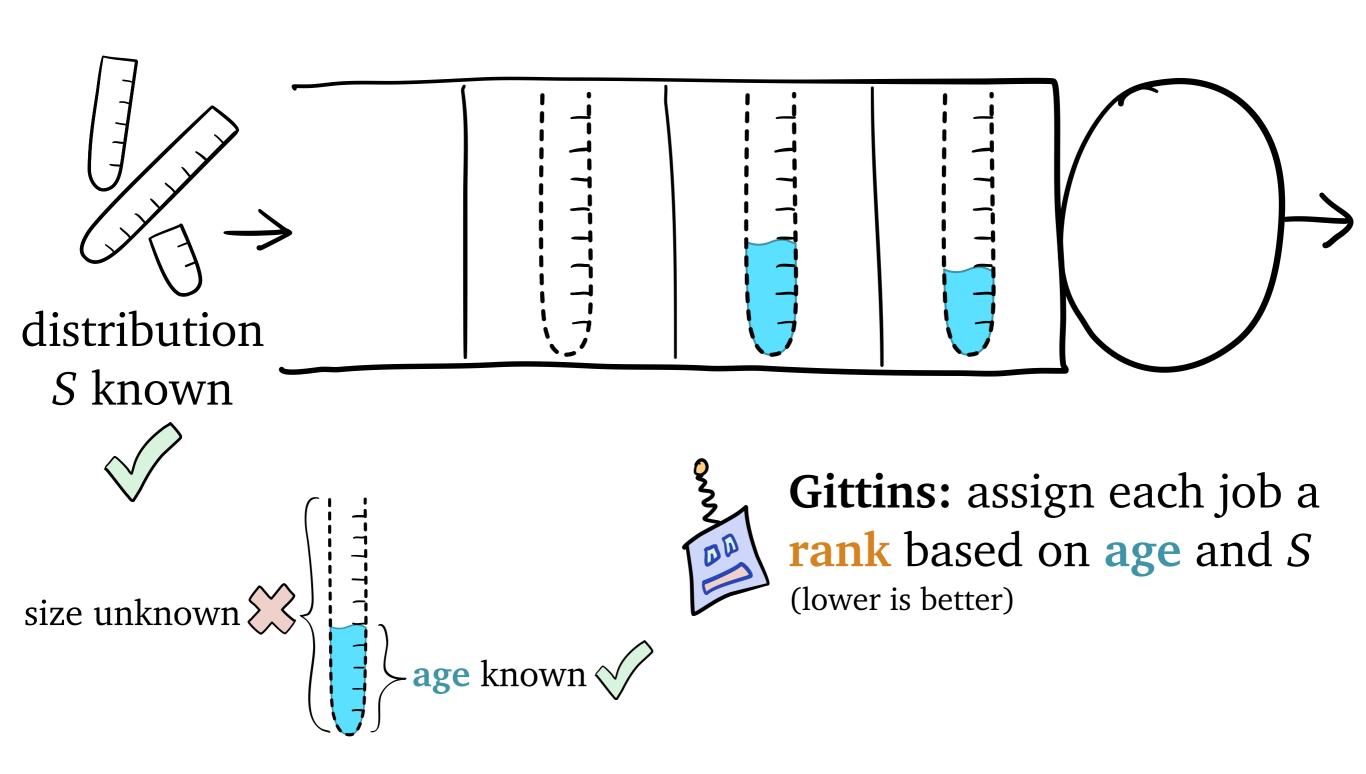


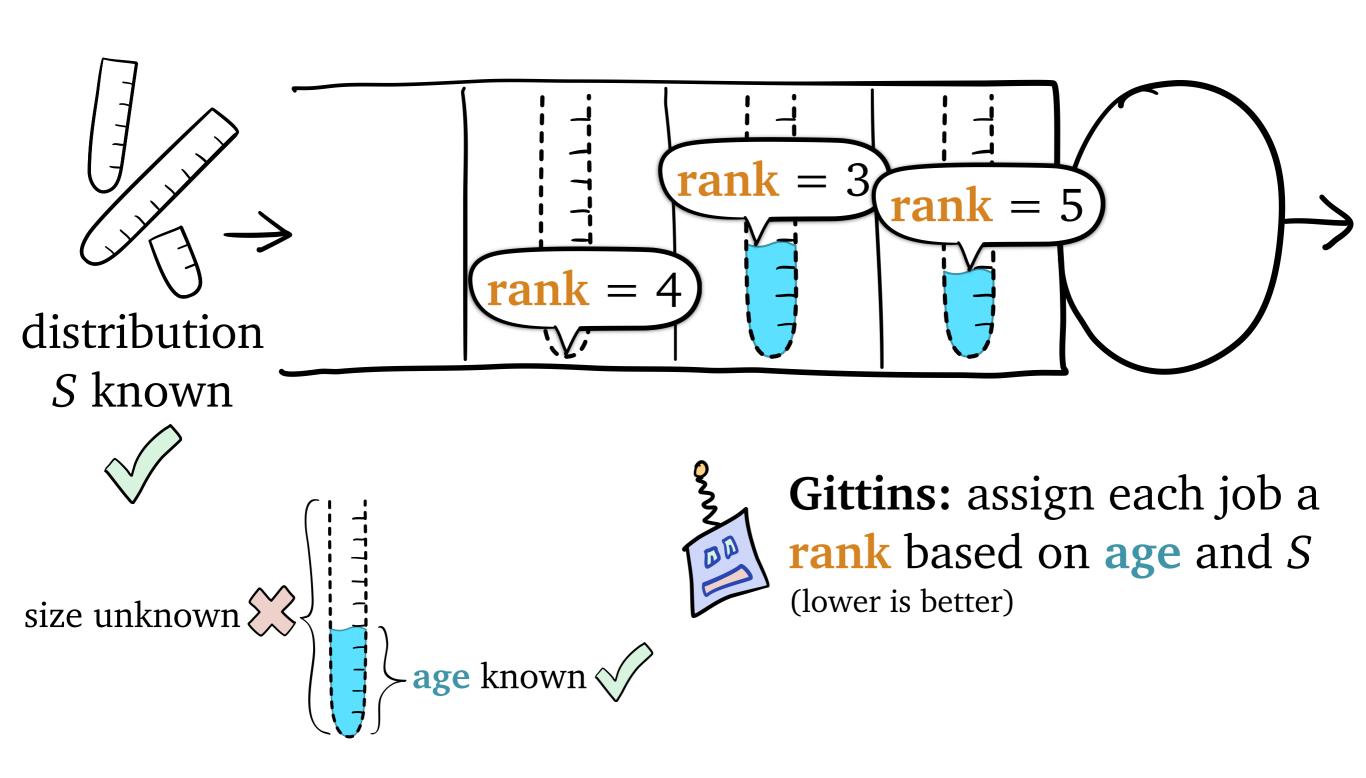


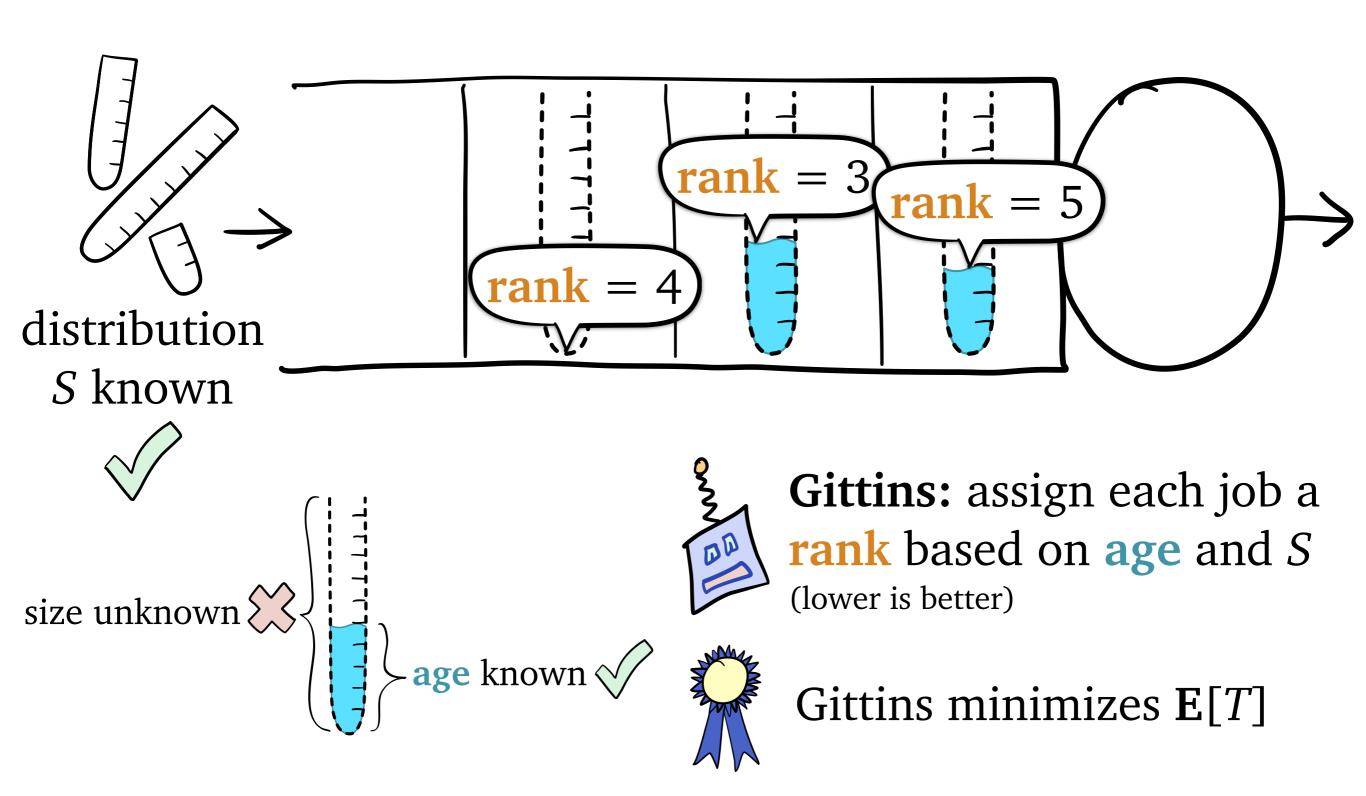












SRPT and Gittins minimize $\mathbf{E}[T]$

known job sizes

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sizes unknown, partially known, known (subsumes SRPT), ...

SRPT and Gittins minimize $\mathbf{E}[T]$

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SRPT and Gittins minimize $\mathbf{E}[T]$

Why not use Gittins?

Gittins Assumption	Computer System Reality

Gittins Assumption Single server

Computer System Reality

Gittins Assumption Computer System Reality Single server Multiple servers

Computer System Reality

Single server

Multiple servers

Complicated implementation not a problem

Computer System Reality

Single server

Buigue beiver

Complicated implementation not a problem

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Simple implementation preferred

Computer System Reality

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Preemption *unrestricted* with *no cost*

Computer System Reality

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Arbitrarily many priority levels

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Want to optimize other response time metrics

Hard

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Hard ... and in this talk!

new queueing-theoretic tools for solving practical scheduling problems

Goals

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Simple implementation preferred

Preemption *restricted* and/or *costly*

Limited number of priority levels

New Tools

Goals

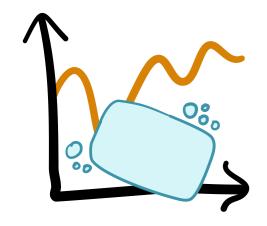
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New Tools



SOAP

analyzes a huge variety of scheduling heuristics

Goals

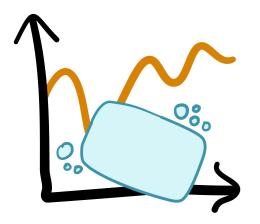
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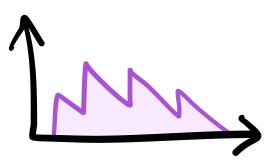
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r-Work

provides a new, deeper understanding of Gittins

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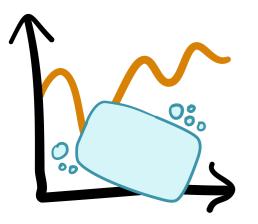
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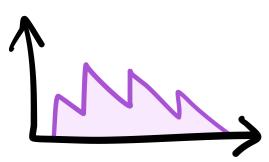
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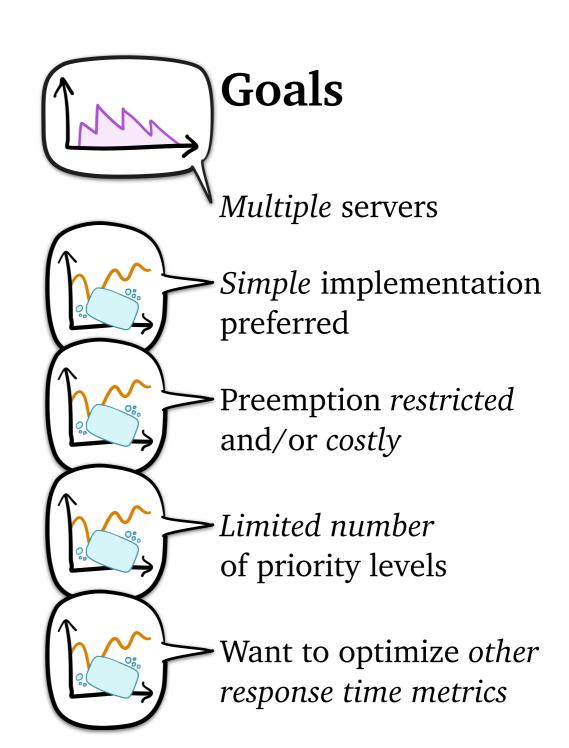
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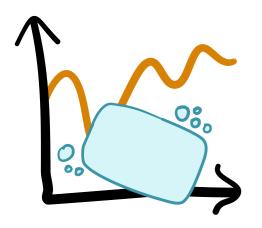


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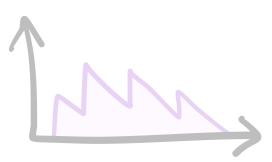


New Tools



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SOAP policies:

broad class of scheduling policies

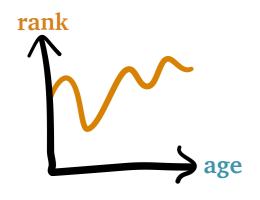


SOAP policies:

broad class of scheduling policies

SOAP analysis:

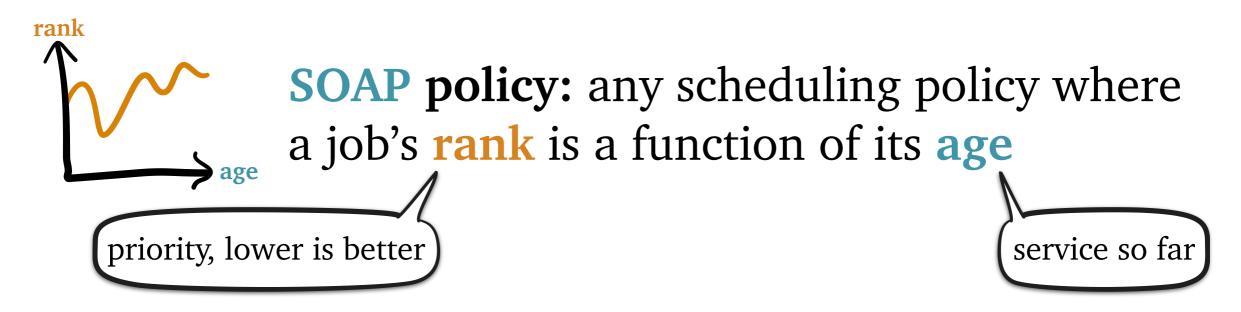
analyze response time of any SOAP policy



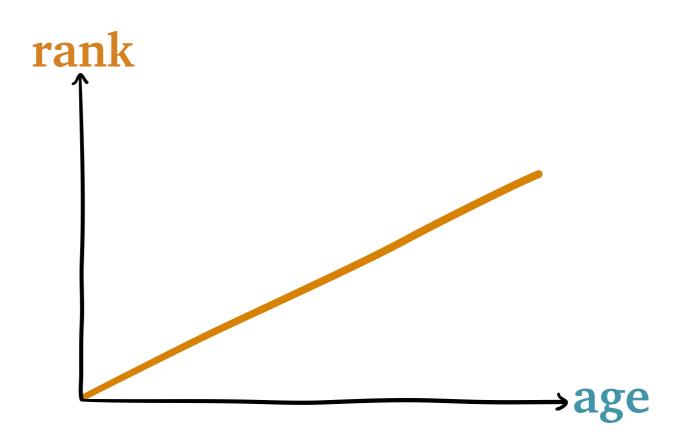
SOAP policy: any scheduling policy where a job's rank is a function of its age

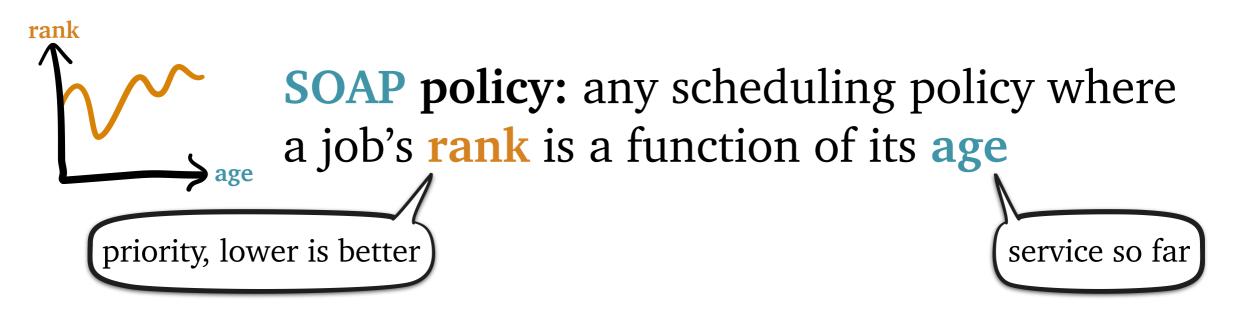




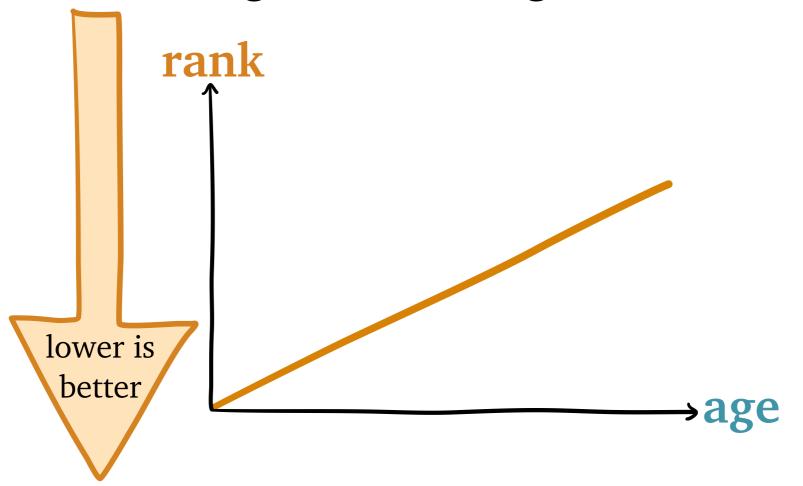


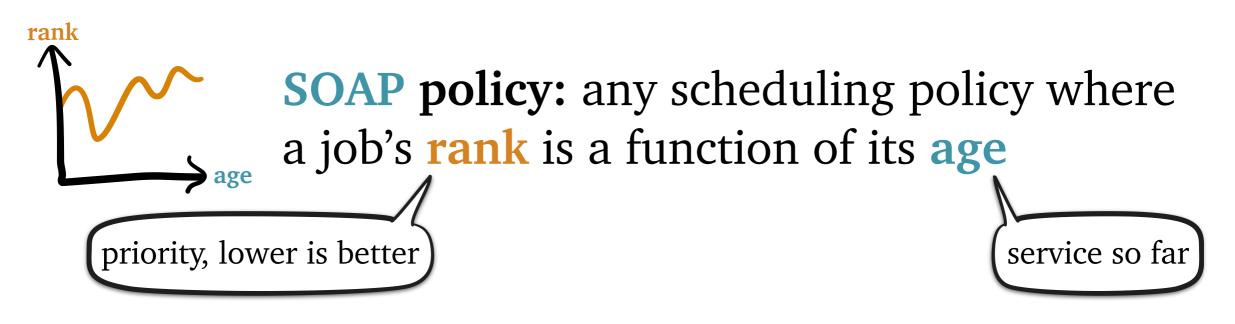
Foreground-Background (FB)

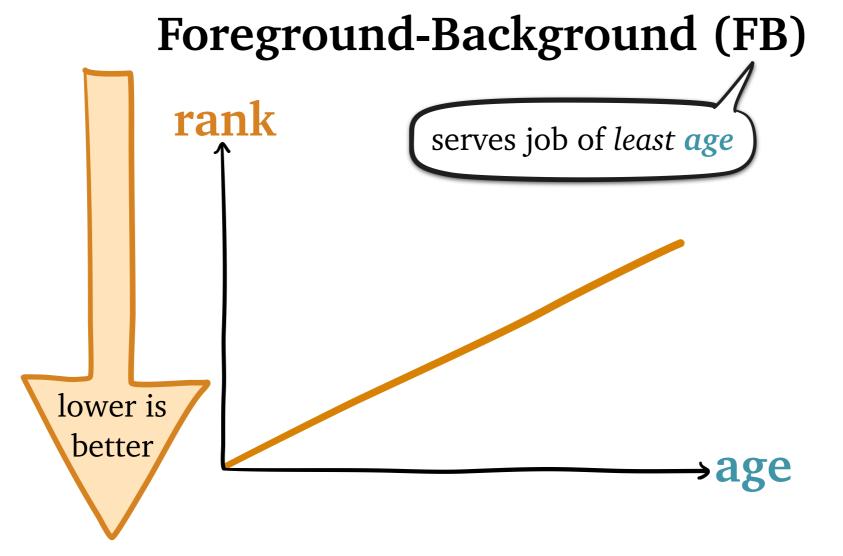


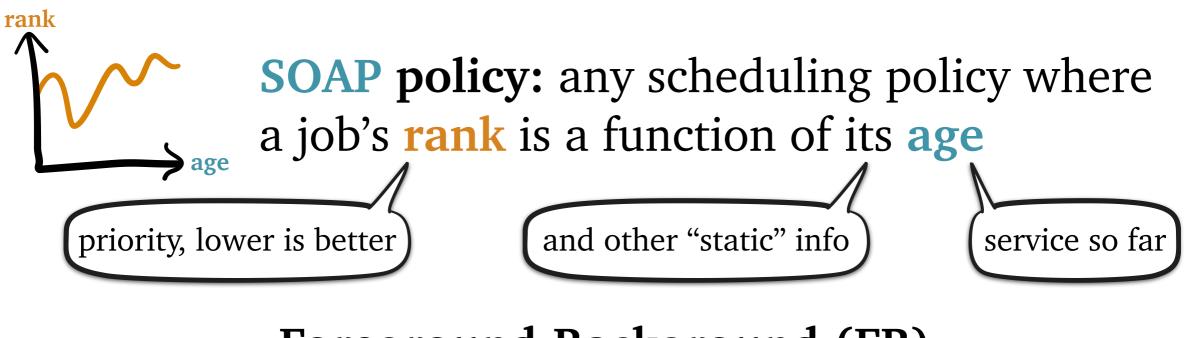


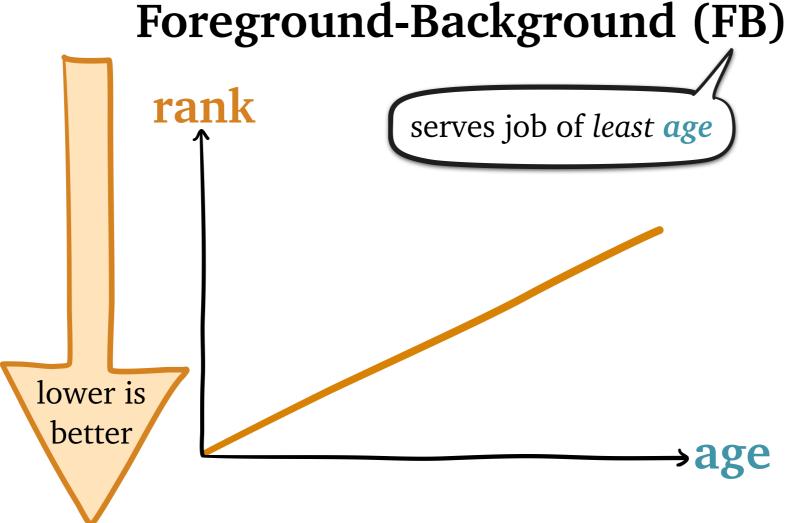
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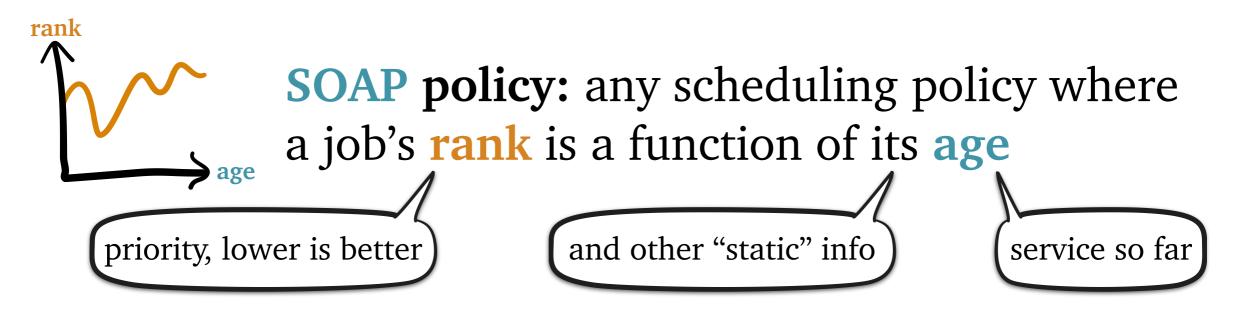


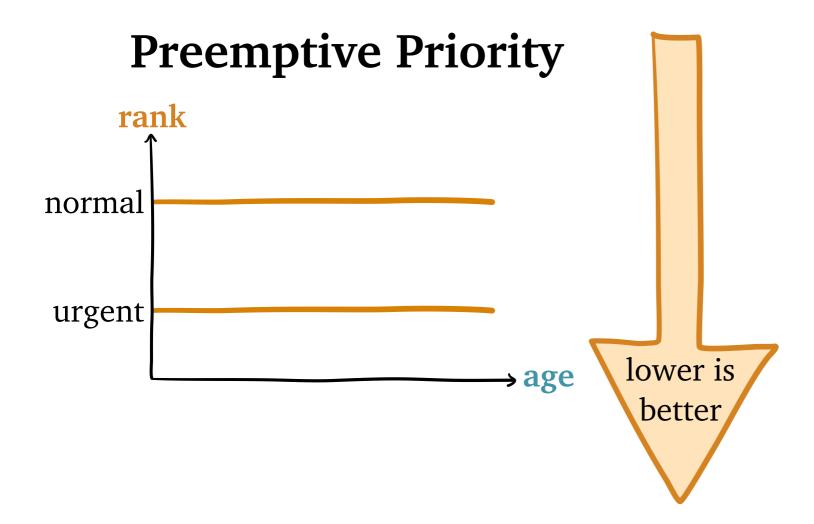


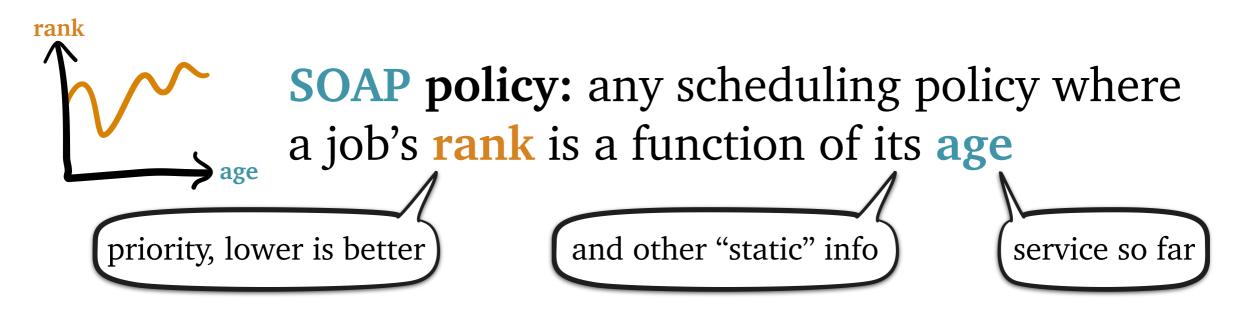


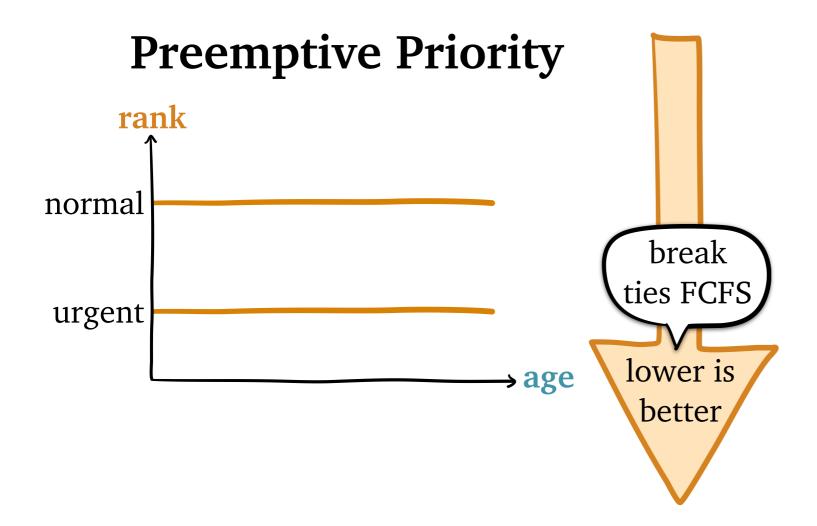


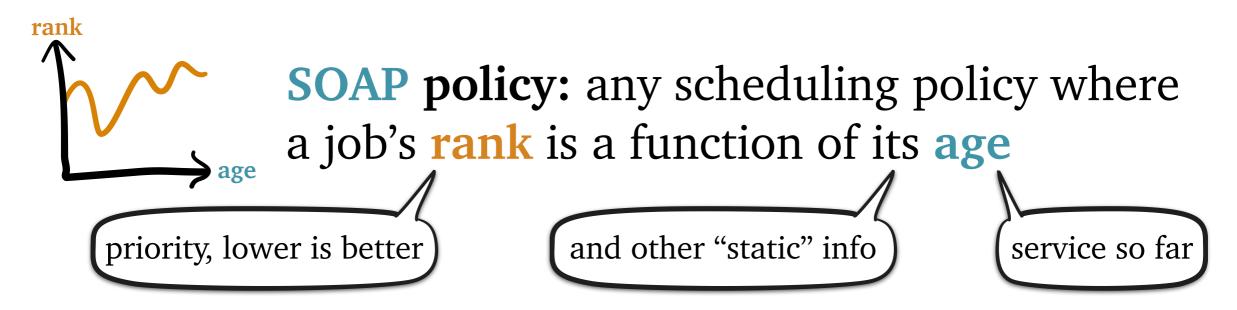


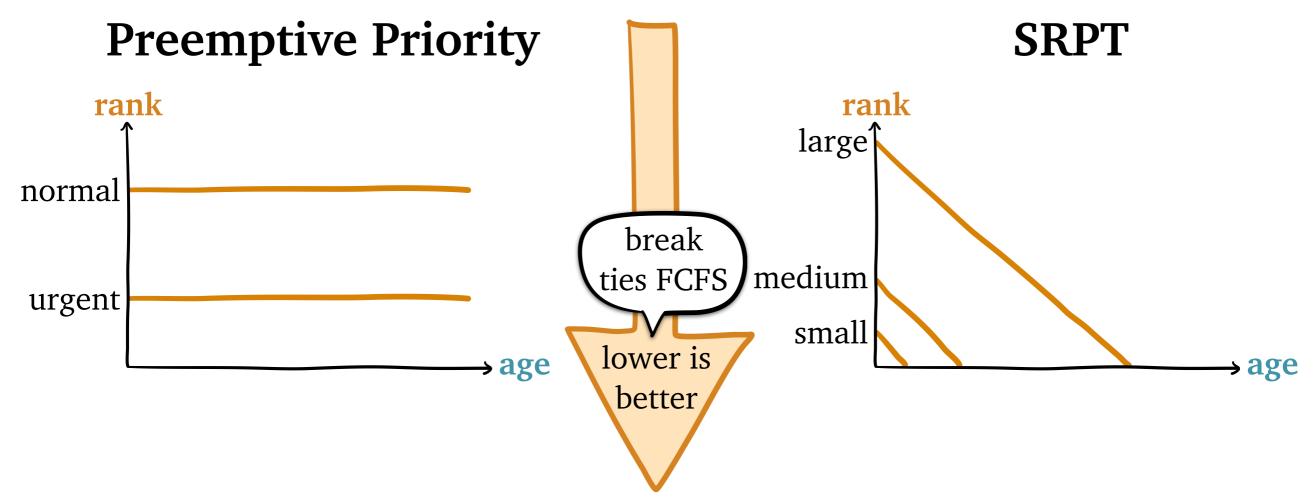




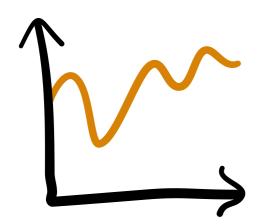






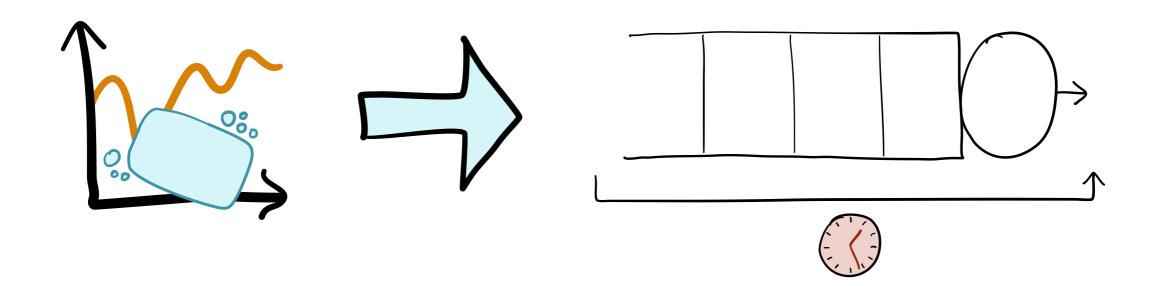


Given any rank function...

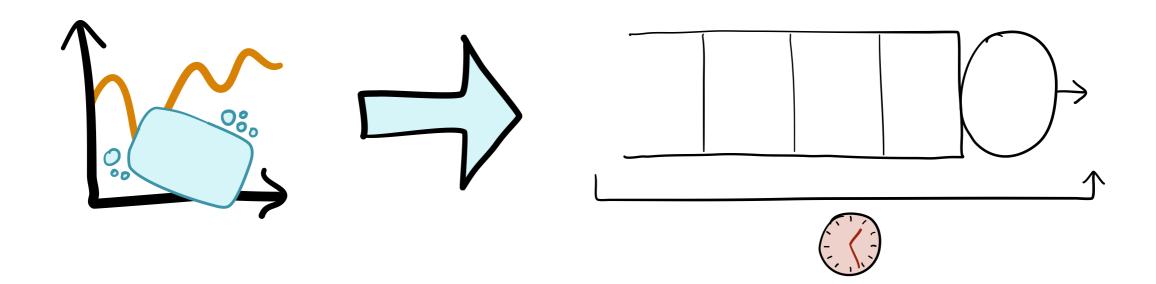


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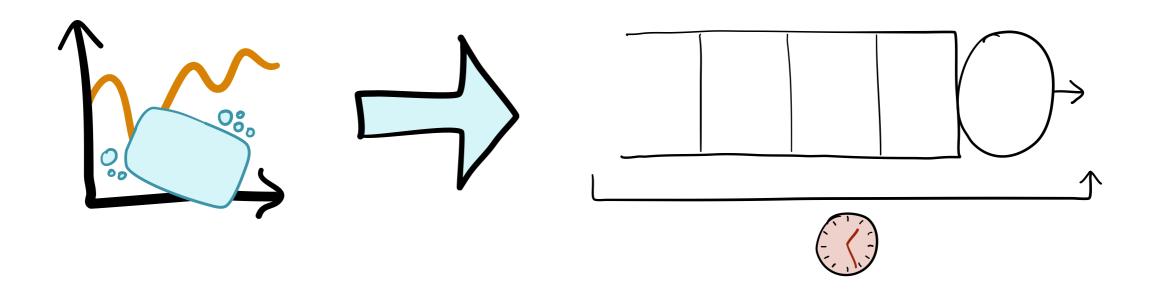
... **SOAP** analyzes its response time



Given any rank function.. (exact formula!)
... SOAP analyzes its response time

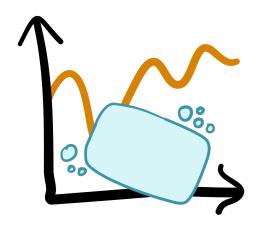


Given any rank function.. (exact formula!)
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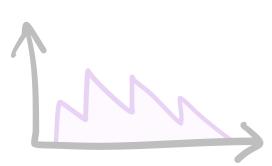
[Scully, Harchol-Balter, & Scheller-Wolf, SIGMETRICS 2018]

New Tools



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Goals

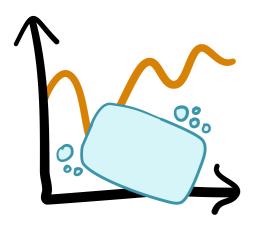
Multiple servers

Simple implementation preferred

Preemption restricted and/or costly

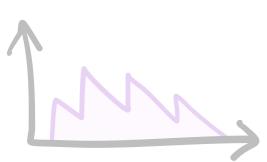
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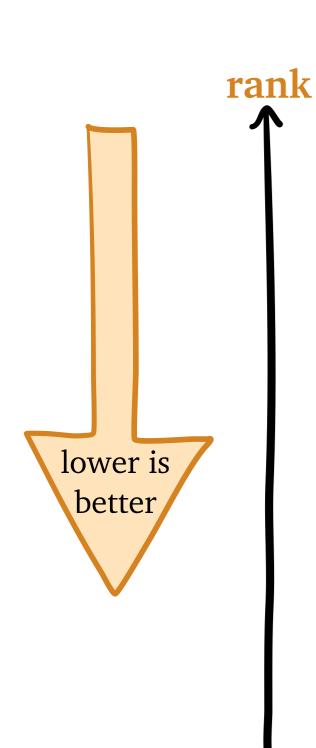
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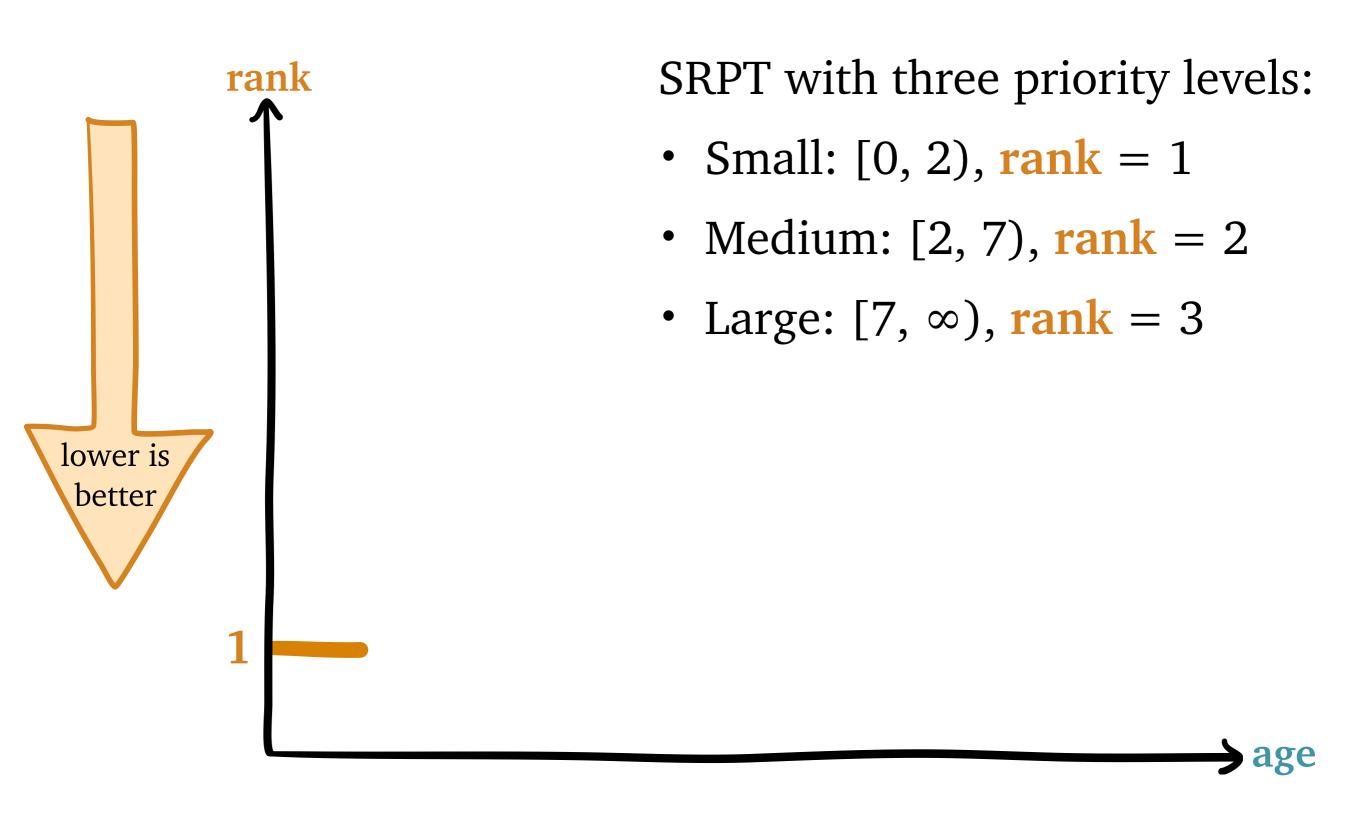
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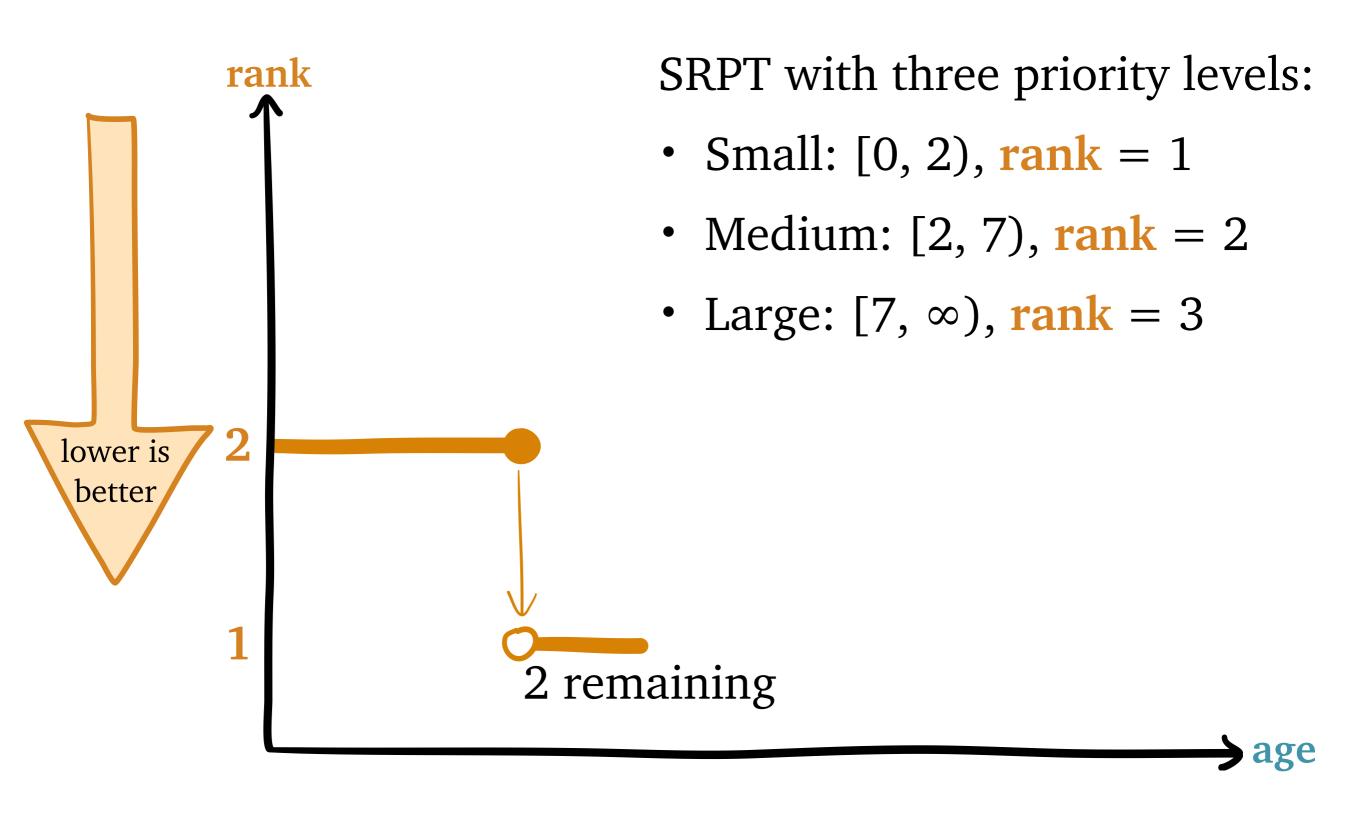


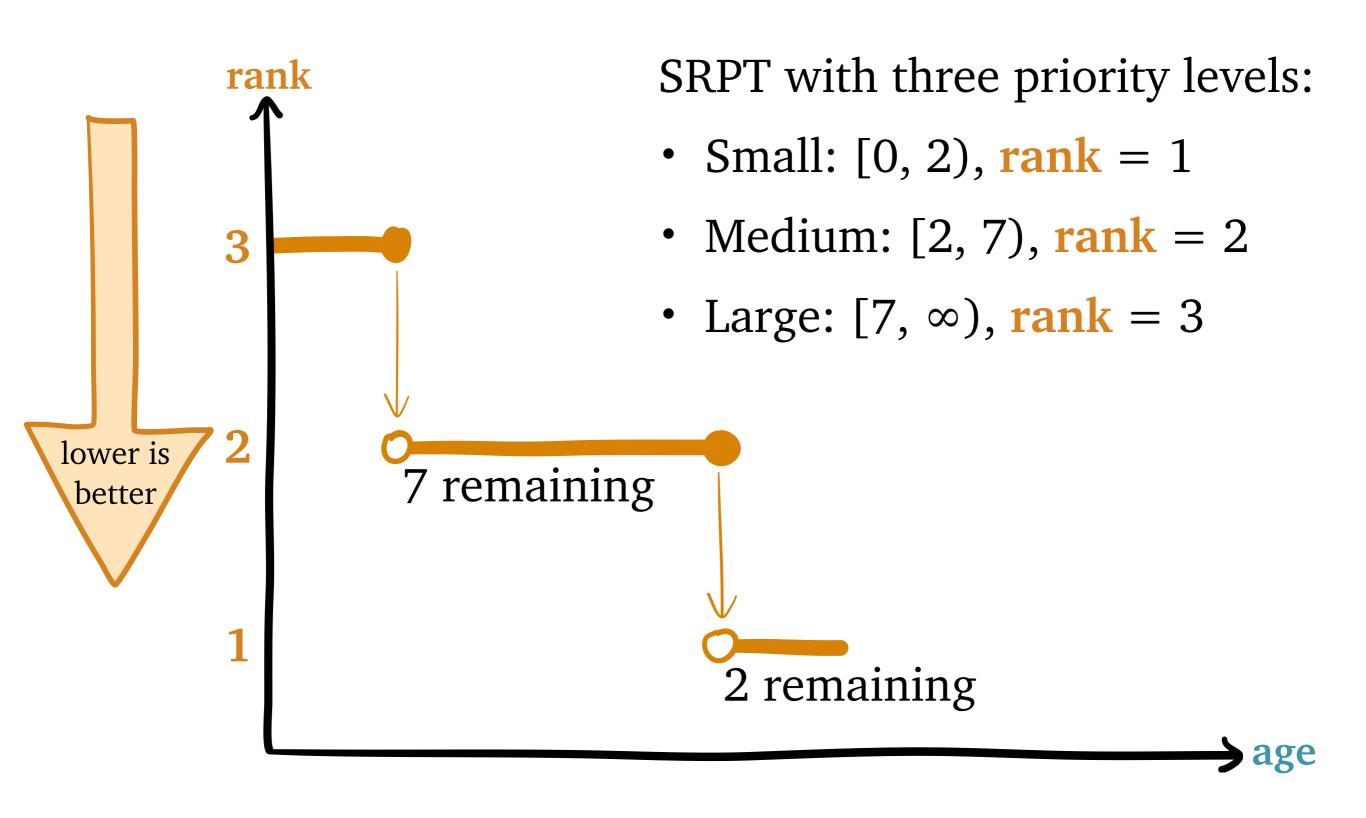


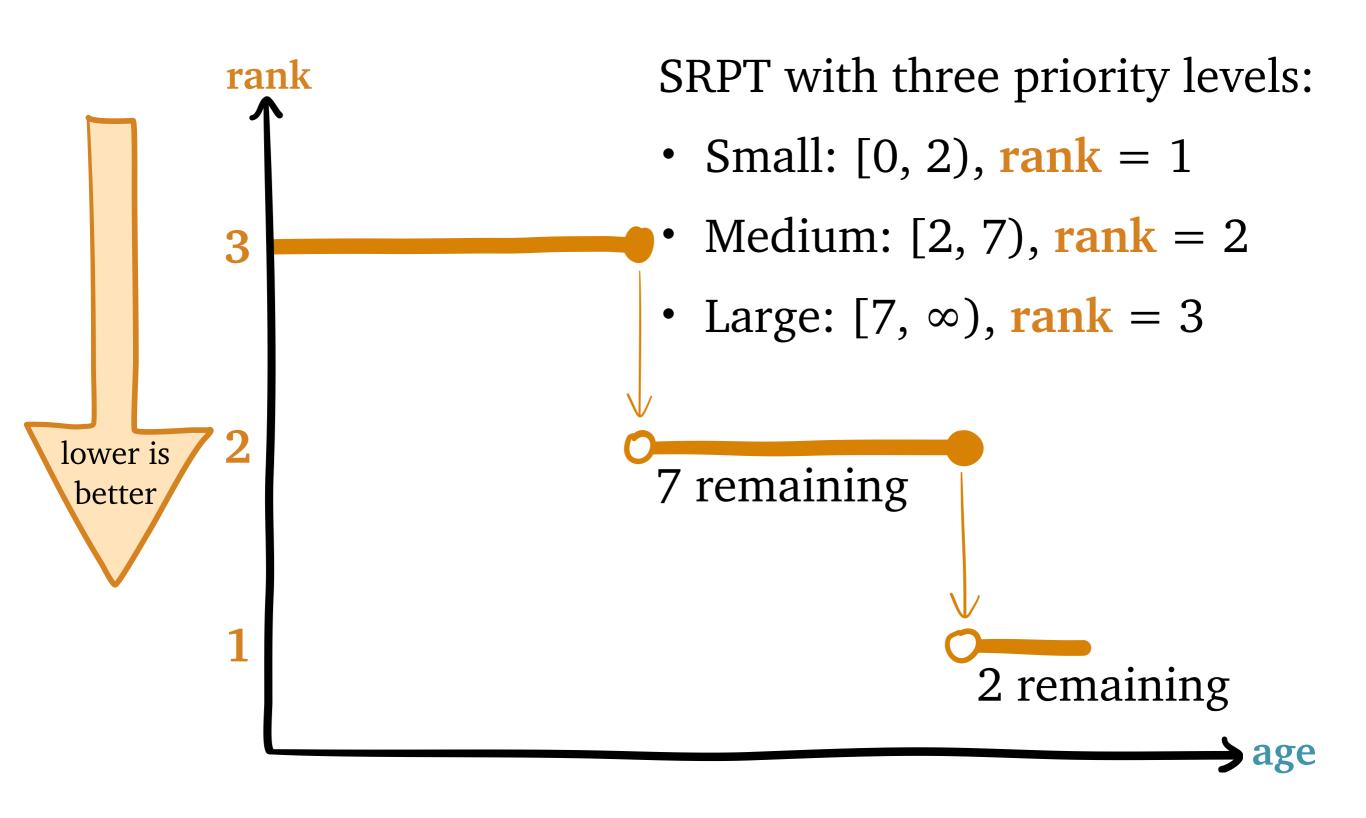
SRPT with three priority levels:

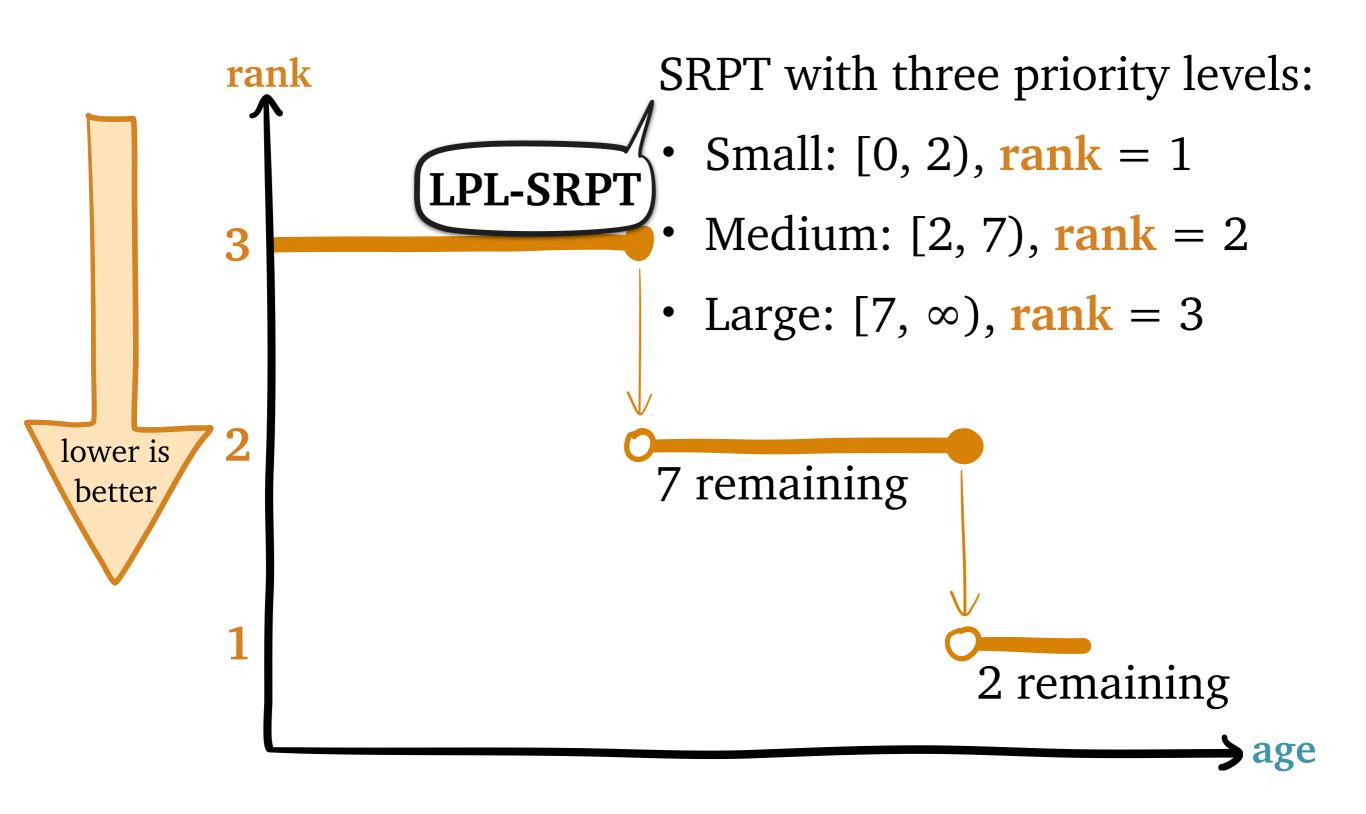
- Small: [0, 2), rank = 1
- Medium: [2, 7), rank = 2
- Large: $[7, \infty)$, rank = 3



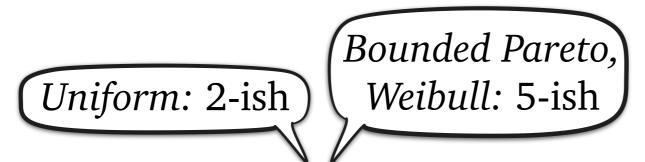








- How many levels do we need?
- How do we choose size cutoffs?
- Can we do better than LPL-SRPT?



- How many levels do we need?
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Load-balancing heuristic suffices

- How many levels do we need?
- How do we choose size cutoffs?
- Can we do better than LPL-SRPT?



Load-balancing heuristic suffices

- How many levels do we need?
- How do we choose size cutoffs?
- Can we do better than LPL-SRPT?

Yes! LPL-PSJF often better

Uniform: 2-ish Bounded Pareto, Weibull: 5-ish

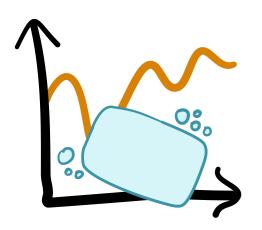
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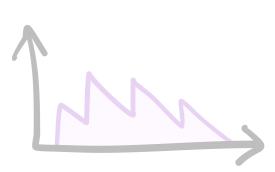
[Scully & Harchol-Balter, in preparation]

New Tools



SOAP

analyzes a huge variety of scheduling heuristics



r-Work

provides a new, deeper understanding of Gittins

Goals

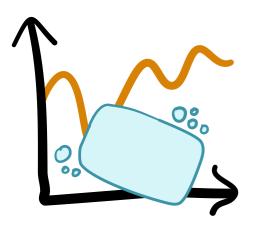
Multiple servers

Simple implementation preferred

Preemption restricted and/or costly

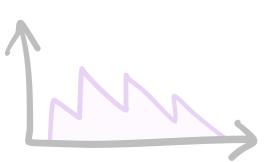
Limited number of priority levels

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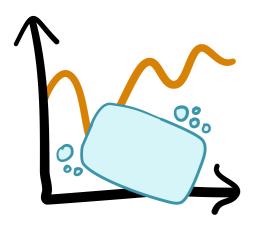
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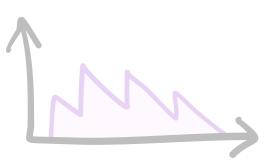
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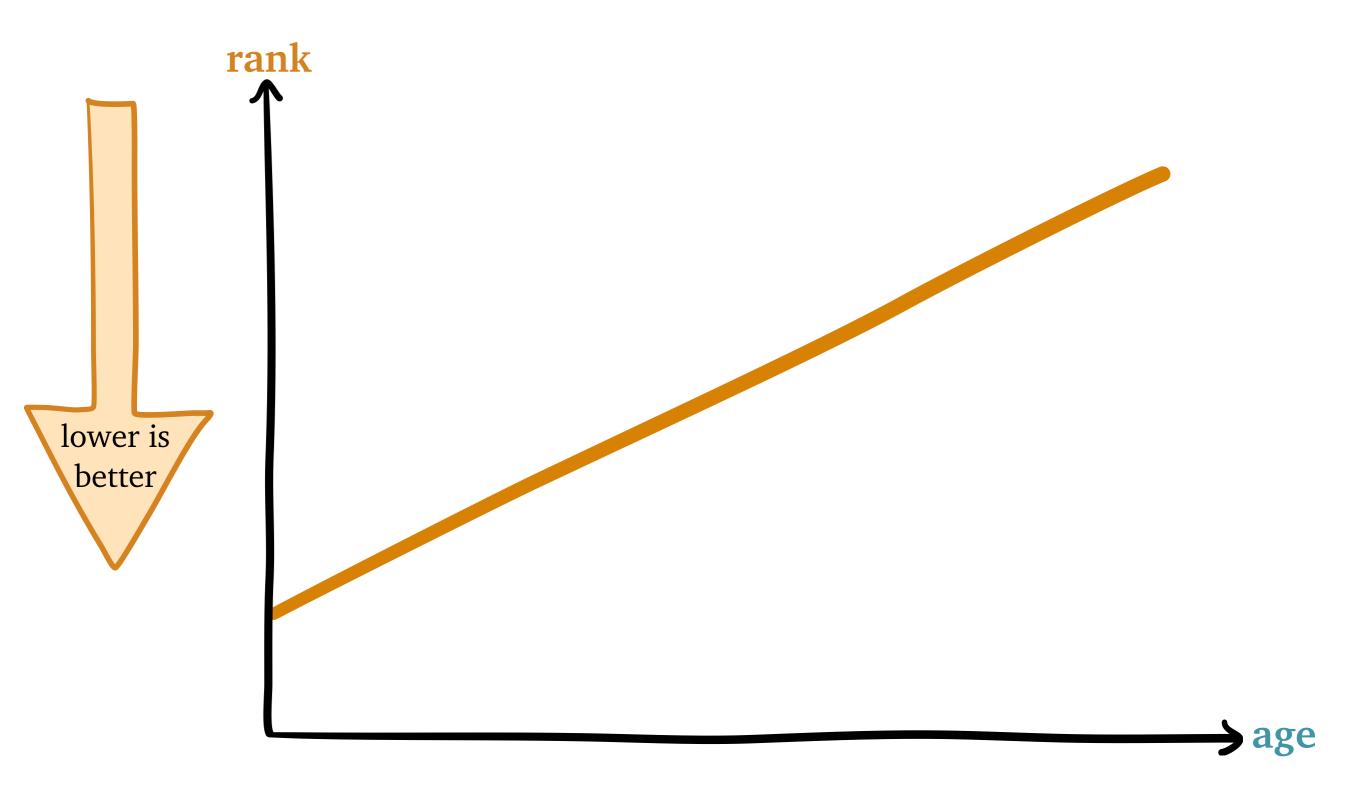
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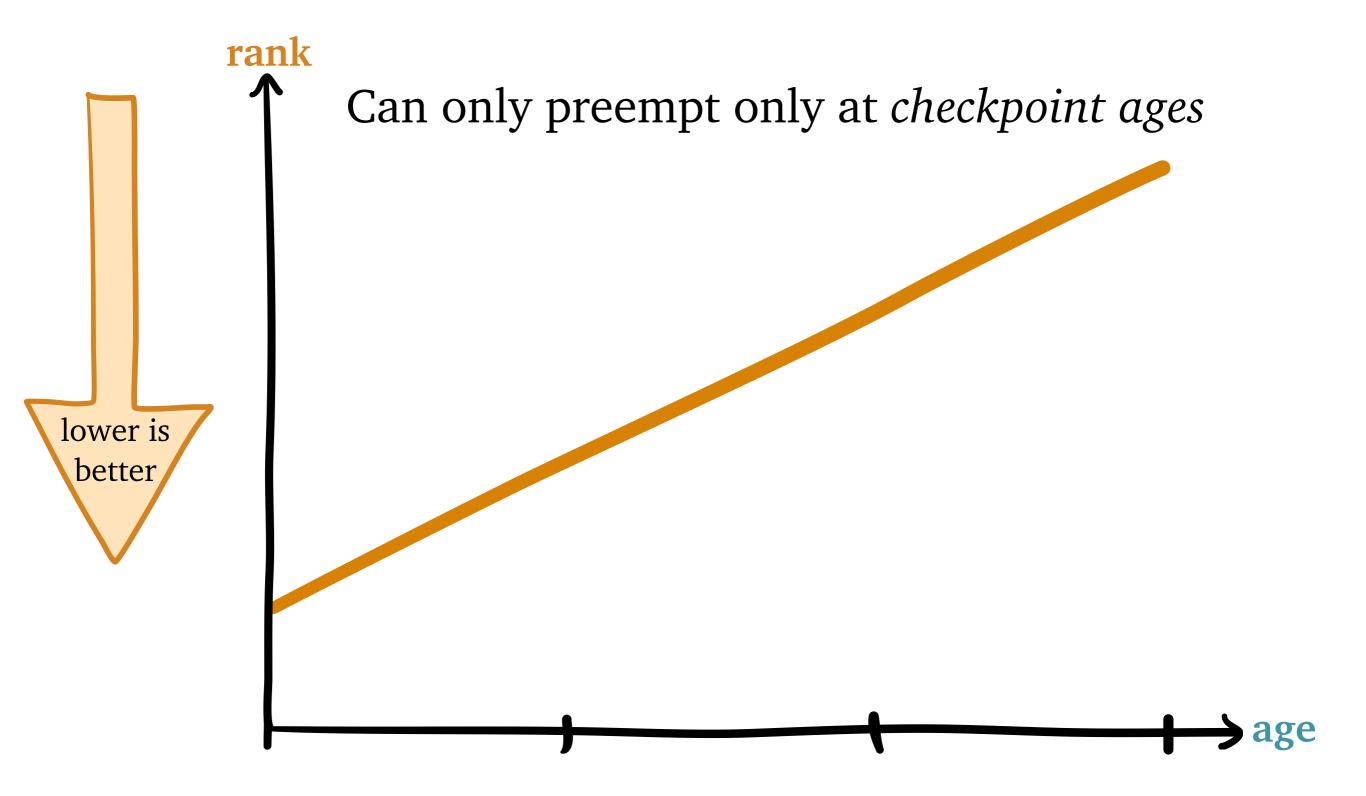
Preemption *restricted* and/or *costly*



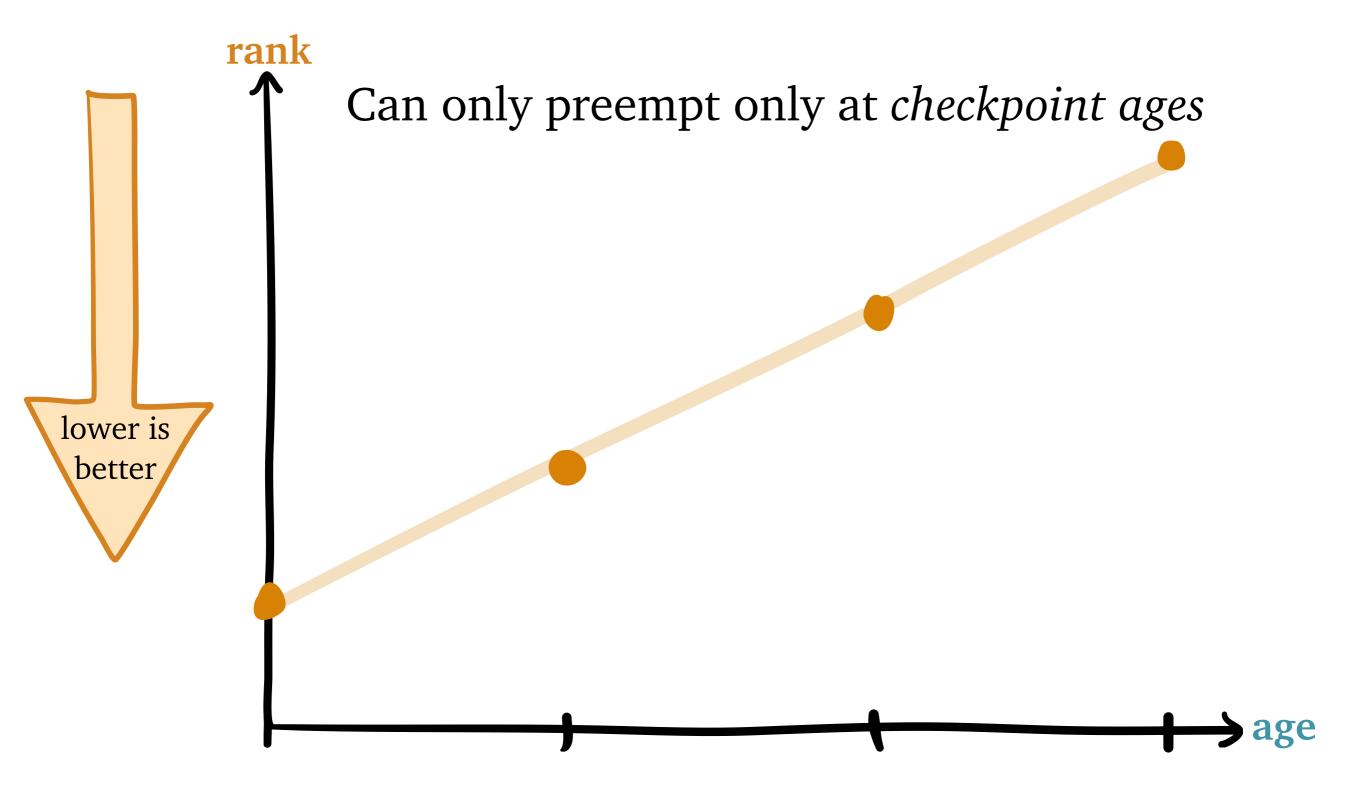
Preemption Checkpoints



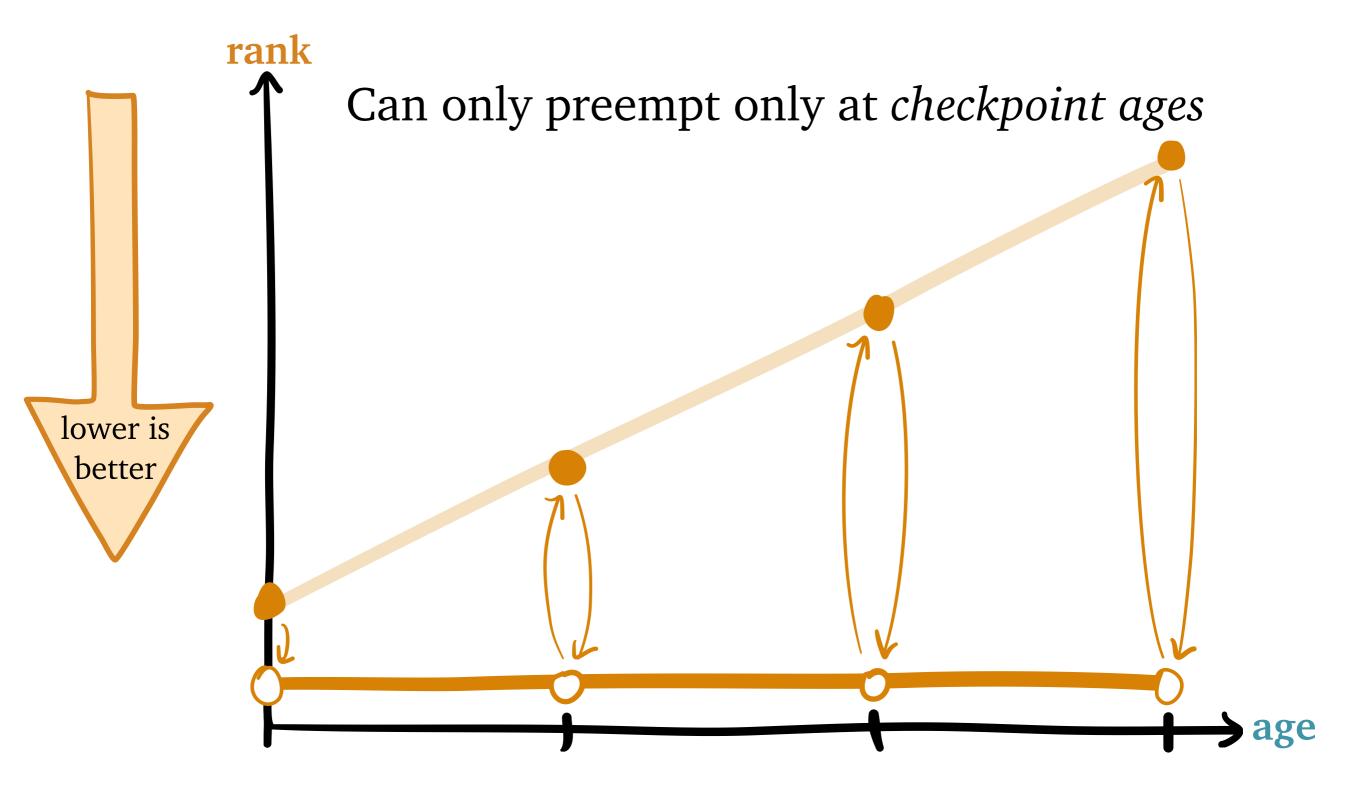
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Preemption Checkpoints



Suppose each checkpoint incurs an overhead What is the optimal gap Δ between checkpoints?

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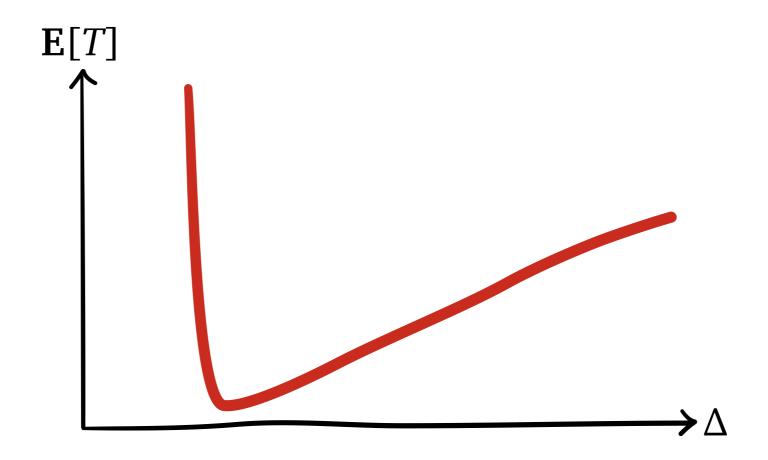
• large Δ : less overhead

Suppose each checkpoint incurs an overhead

- large Δ : less overhead
- small Δ : better scheduling

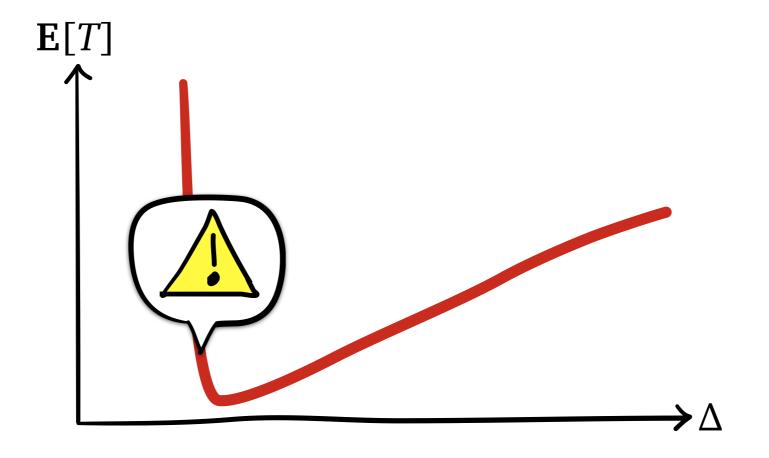
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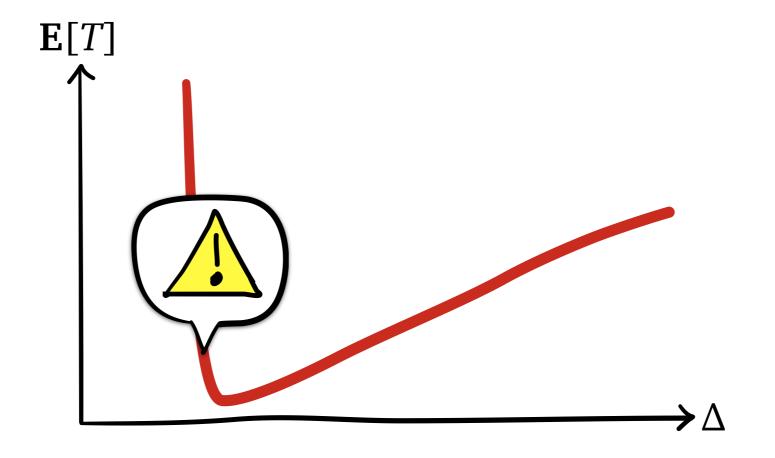
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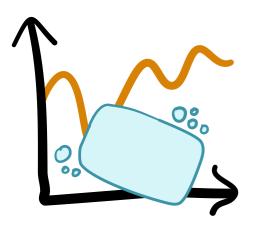
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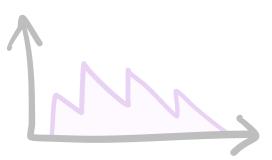
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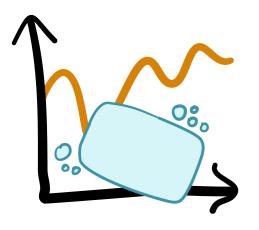
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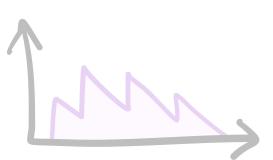


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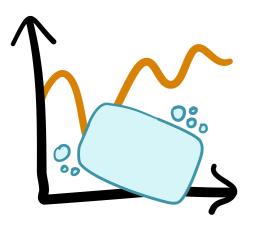


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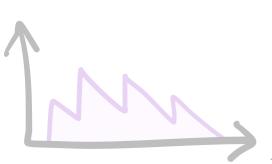
Limited number of priority levels

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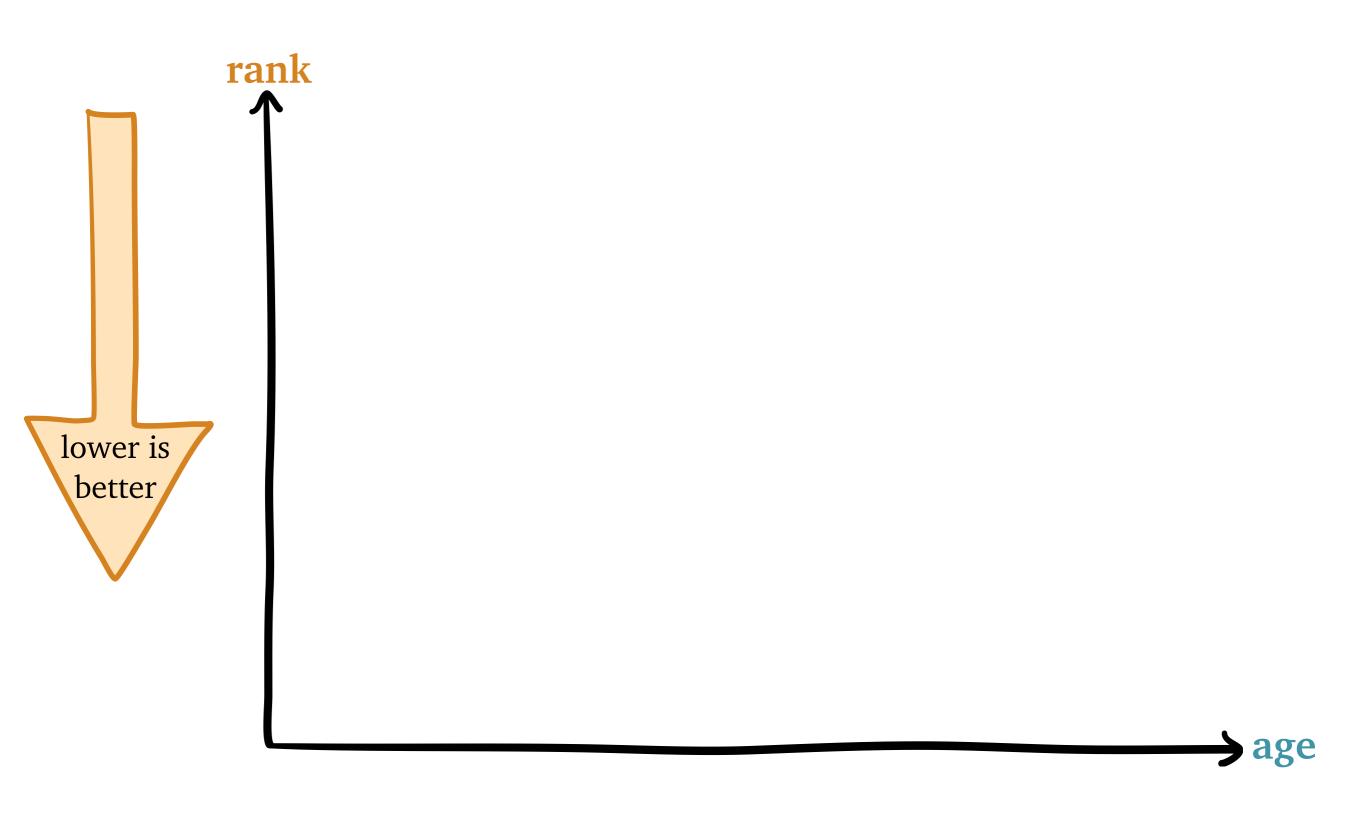
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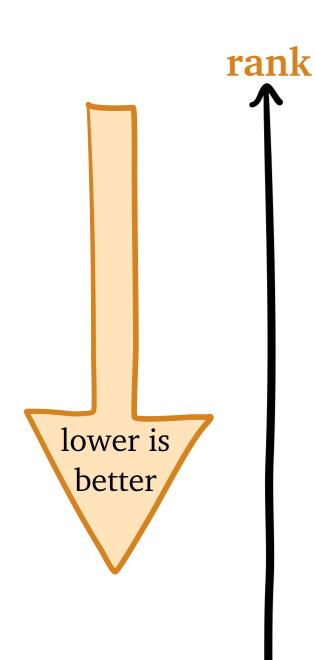


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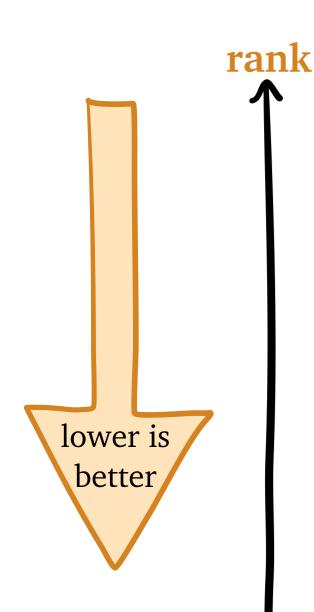


Limited number of priority levels





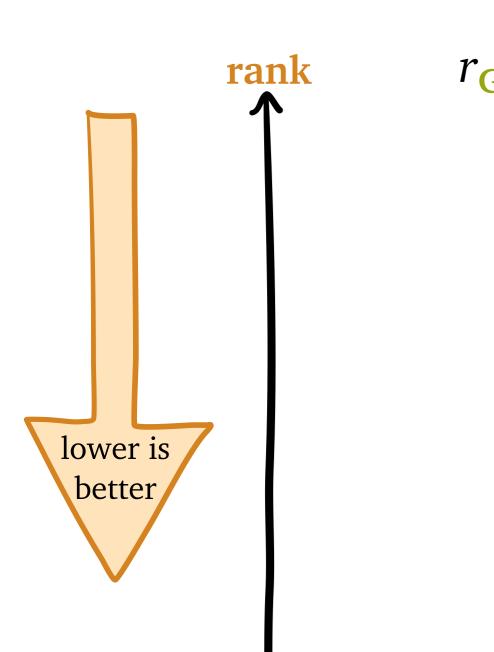
$$S = \begin{cases} 1 & \text{w.p. } \frac{1}{3} \\ 6 & \text{w.p. } \frac{1}{3} \\ 14 & \text{w.p. } \frac{1}{3} \end{cases}$$



$$r_{\text{Gittins}}(a) = \inf_{b>a} \frac{\mathbf{E}[\min\{S-a,b\} \mid S>a]}{\mathbf{P}[S \leq b \mid S>a]}$$

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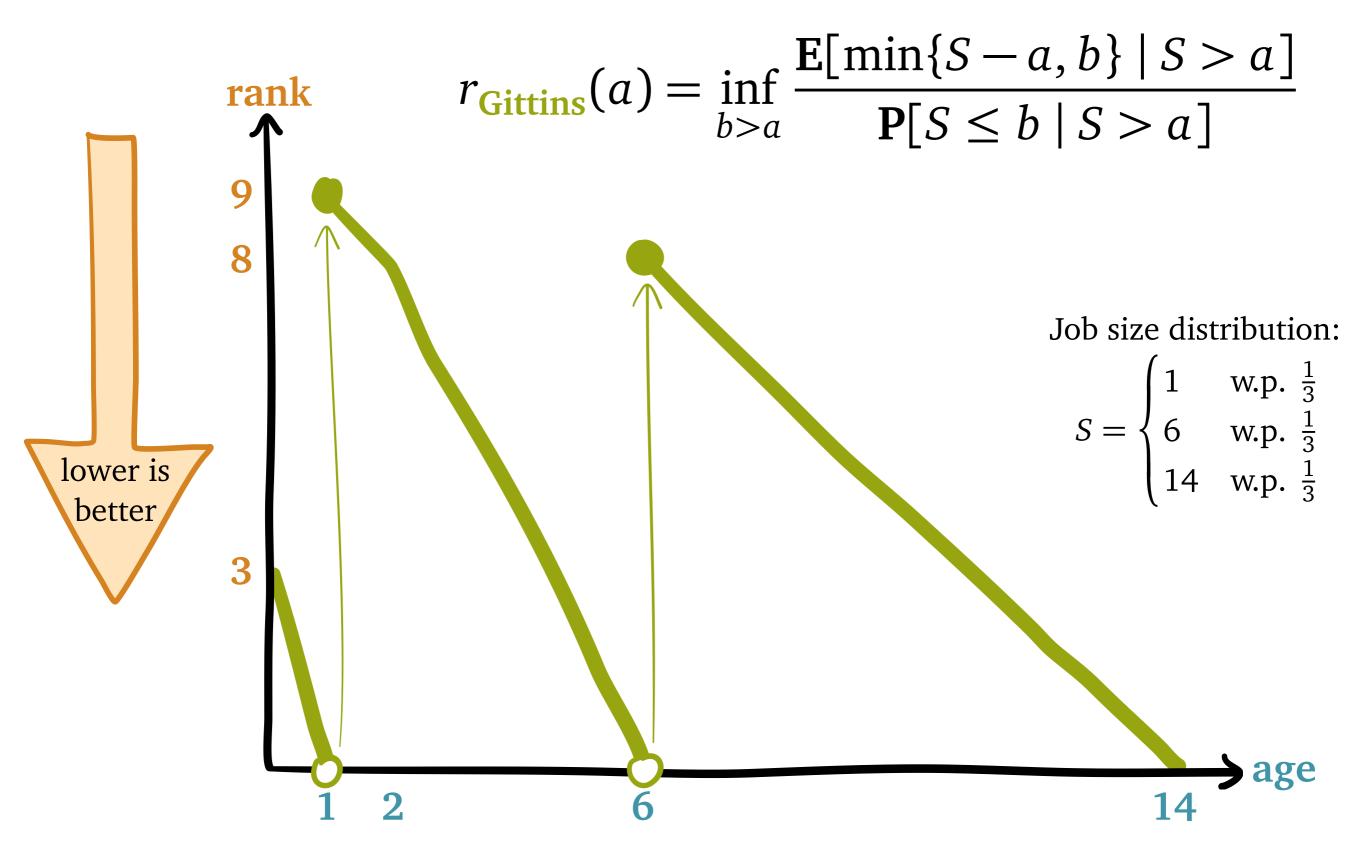


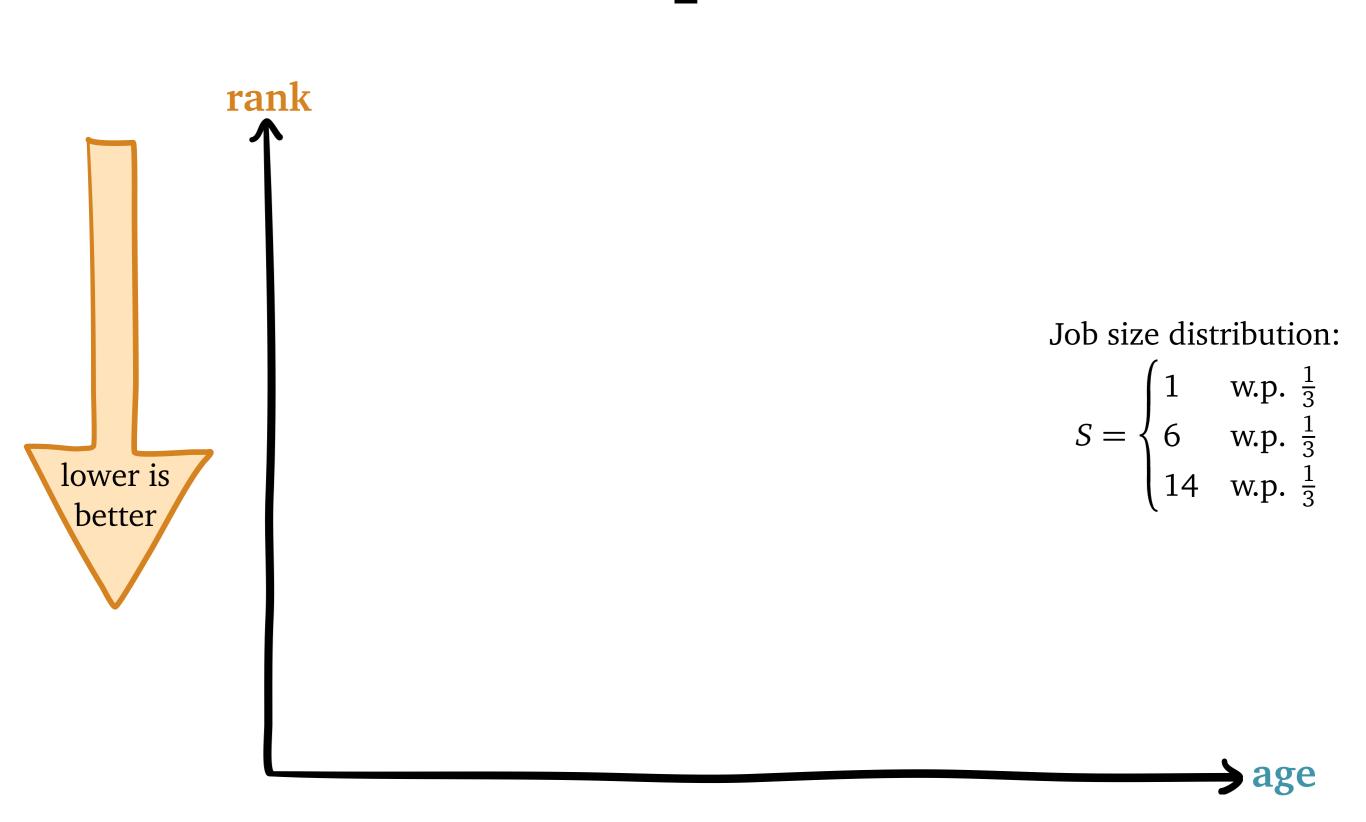


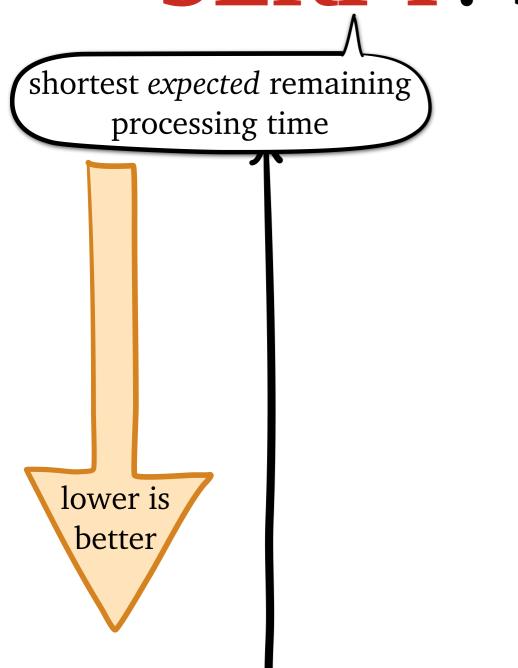
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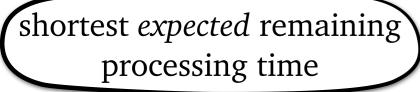








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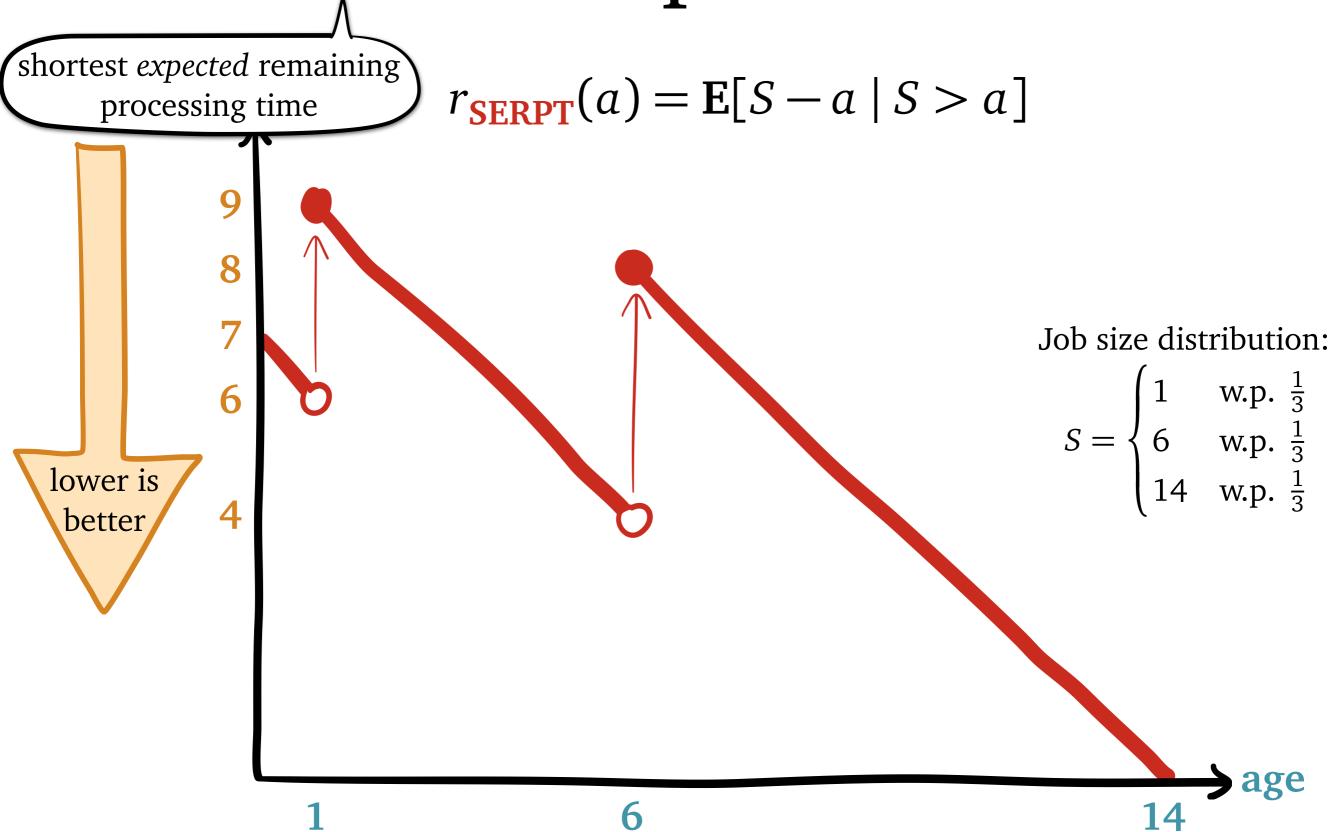
lower is

better

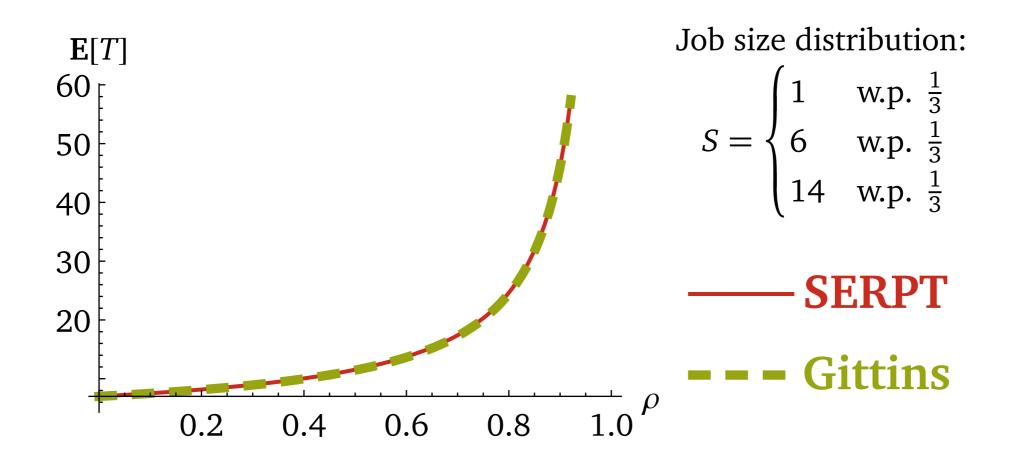
$$r_{\text{SERPT}}(a) = \mathbf{E}[S - a \mid S > a]$$

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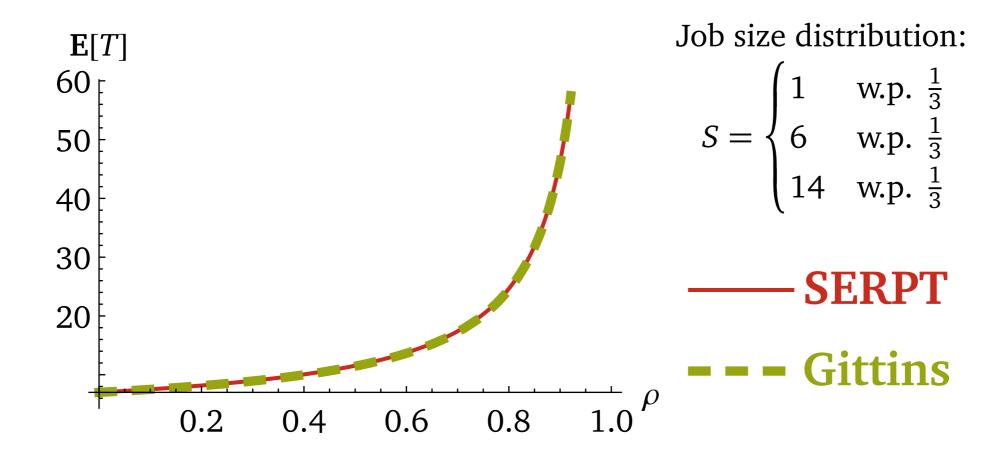




Can SERPT Replace Gittins?



Can SERPT Replace Gittins?





- Gittins is hard to compute
- **SERPT** has no E[T] guarantee



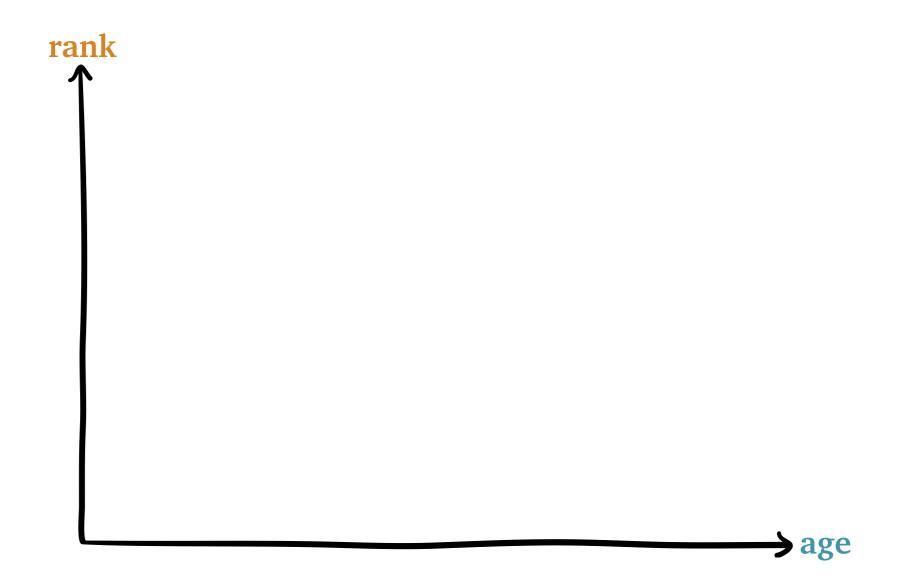
I wish for a policy with...

- simple definition like **SERPT**
- provable guarantee on E[T] like Gittins

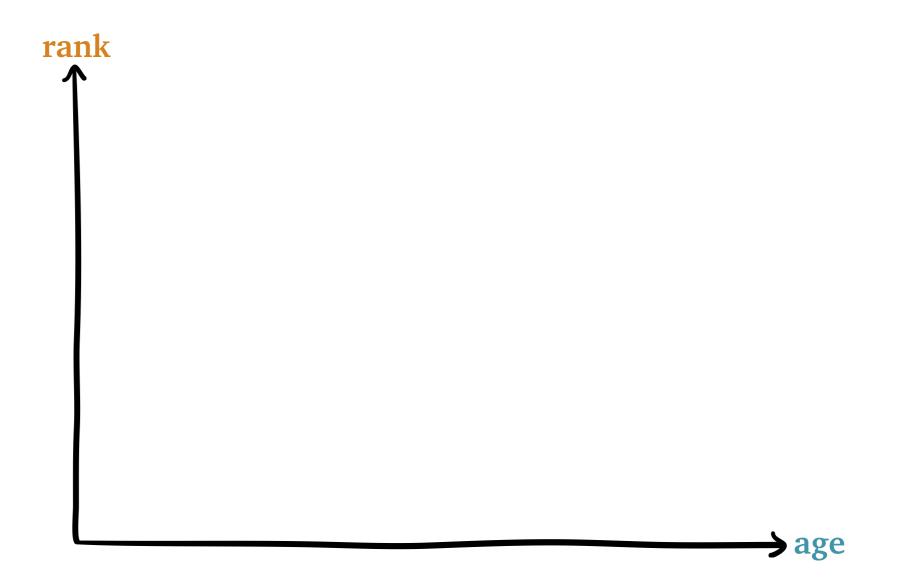


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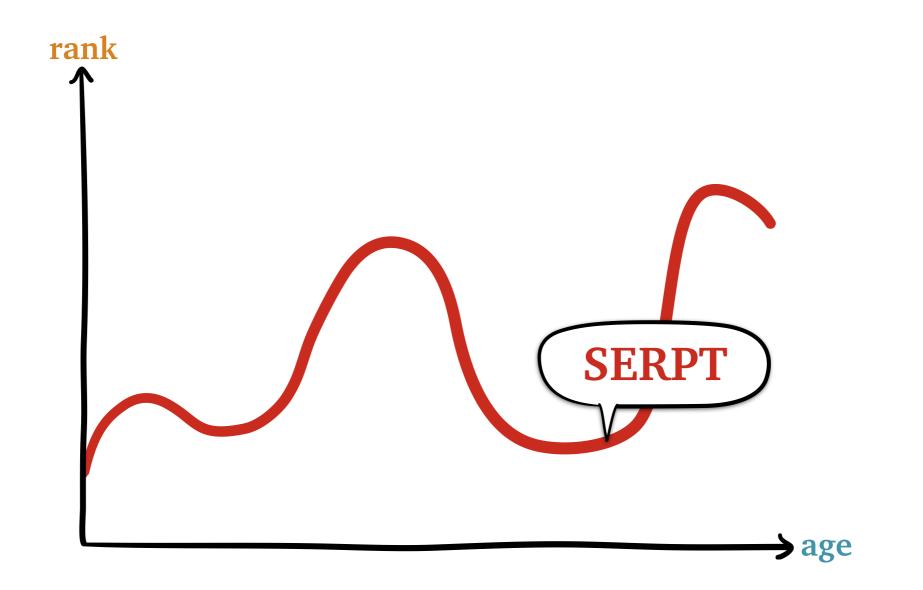
$$r_{\mathbf{M-SERPT}}(a) = \max_{0 \le b \le a} r_{\mathbf{SERPT}}(b)$$



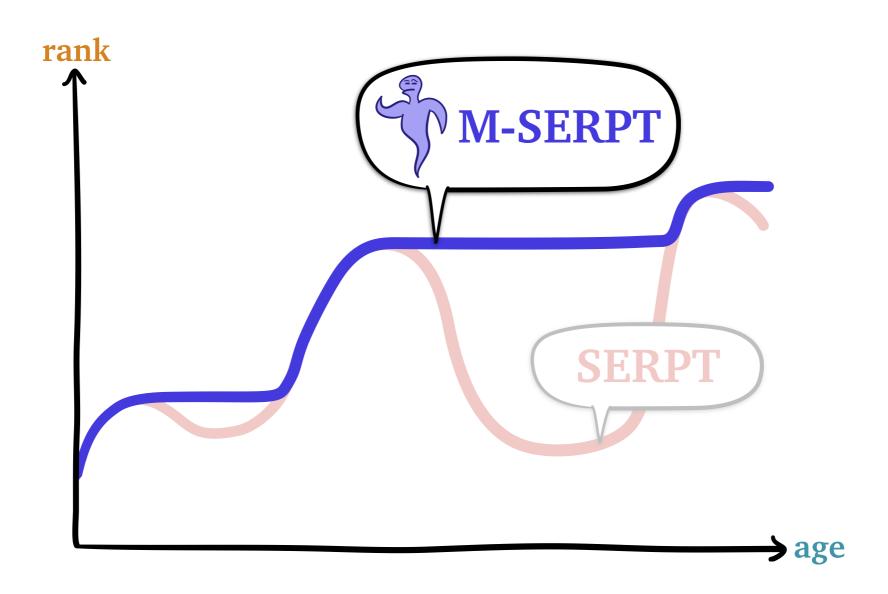
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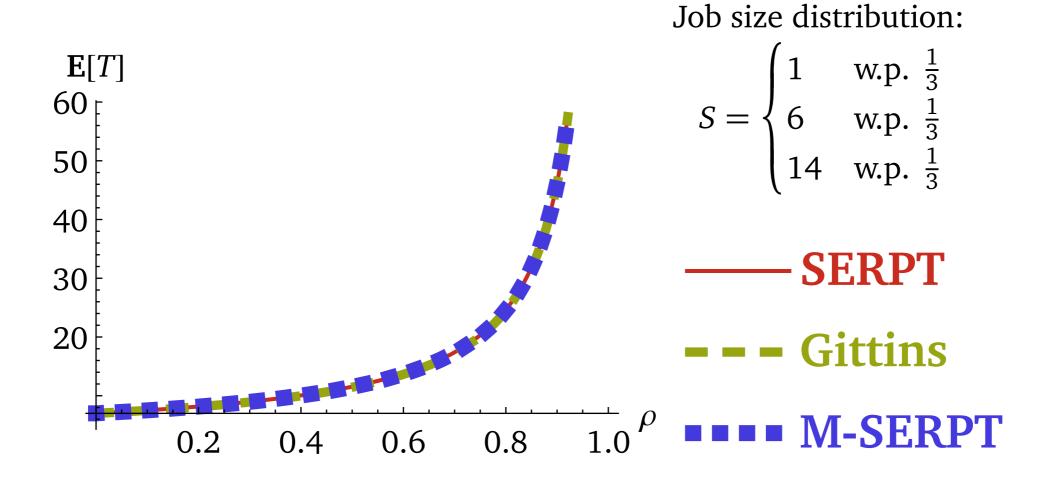


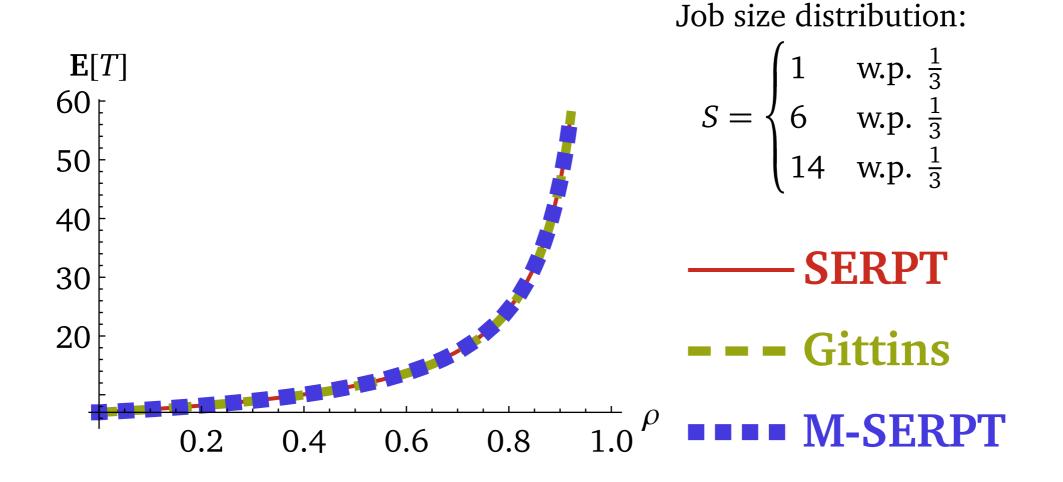
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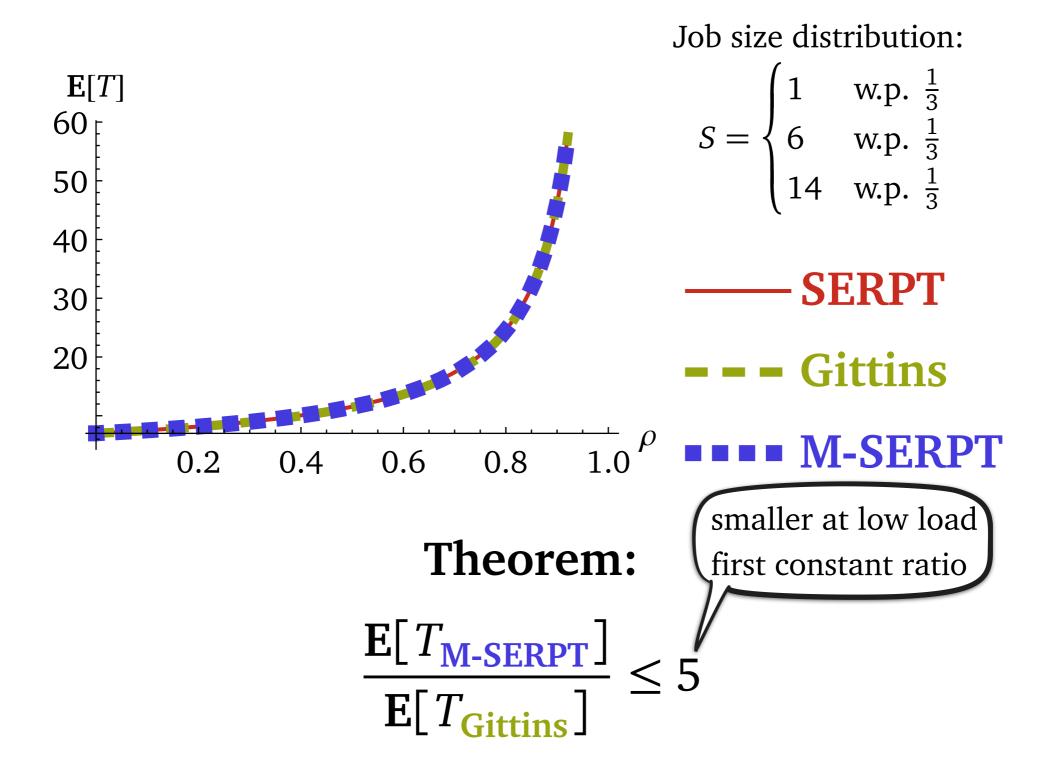


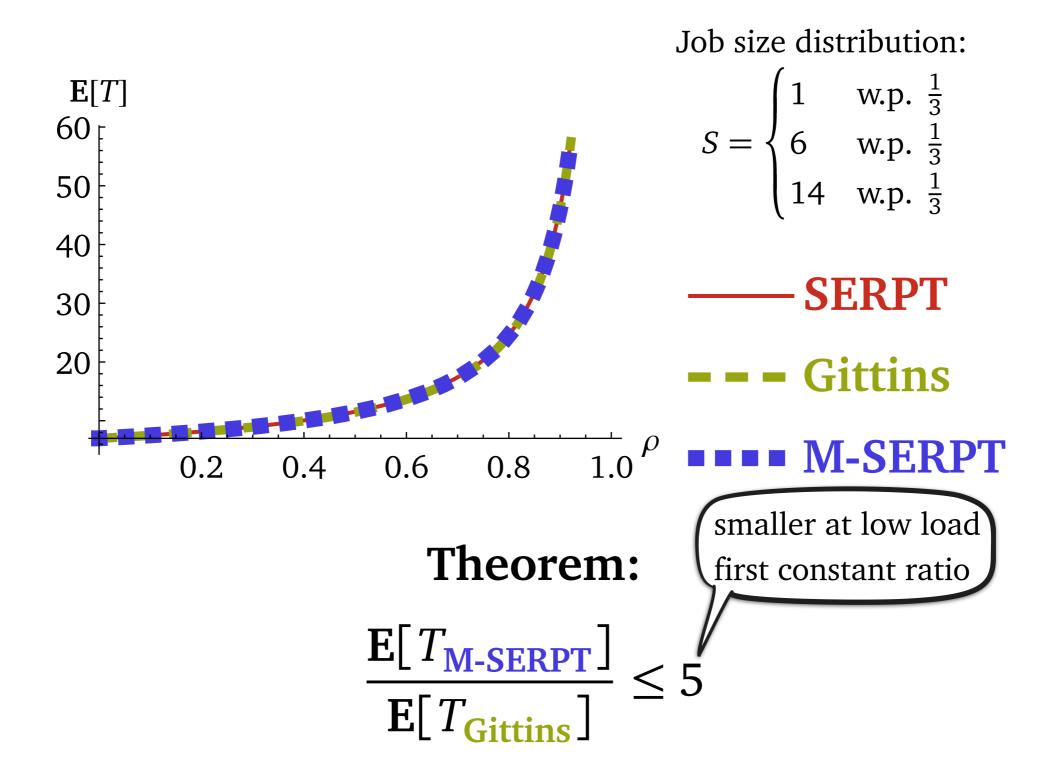




Theorem:

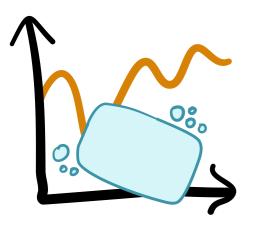
$$\frac{\mathbf{E}[T_{\mathbf{M-SERPT}}]}{\mathbf{E}[T_{\mathbf{Gittins}}]} \le 5$$





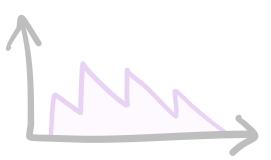
[Scully, Harchol-Balter, & Scheller-Wolf, SIGMETRICS 2020]

New Tools



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Goals

Multiple servers

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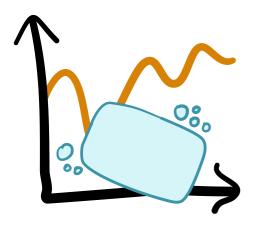


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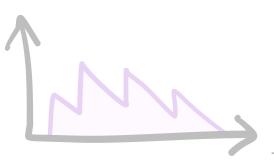
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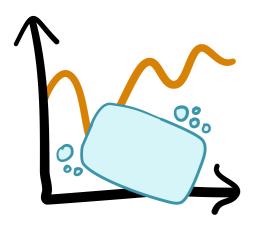


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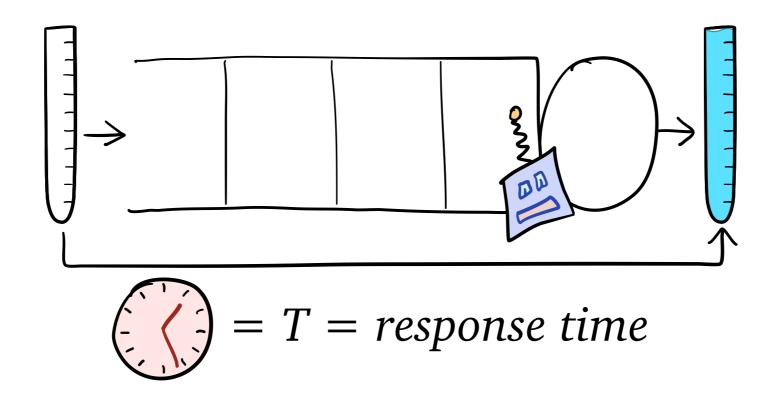
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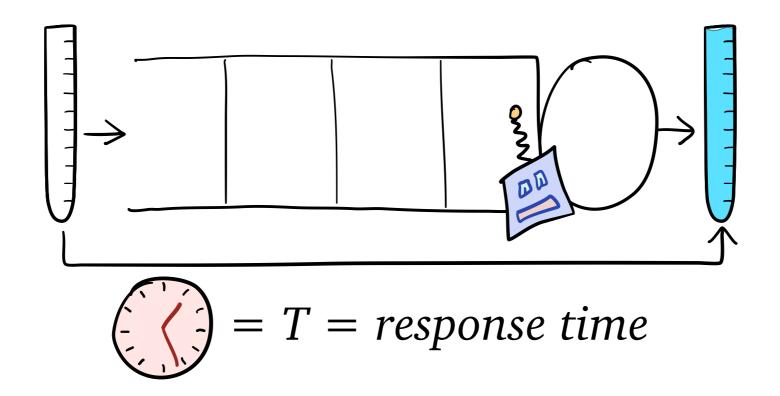


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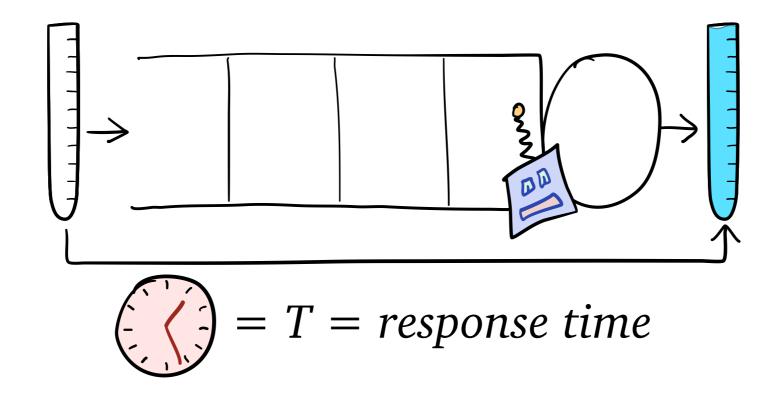


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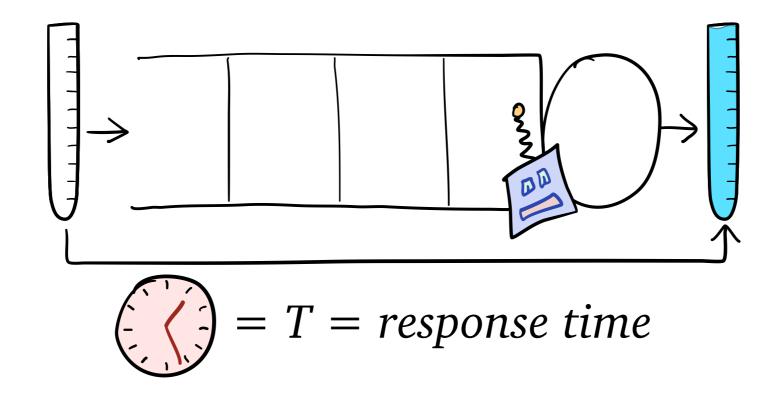


Goal: schedule to minimize two metrics



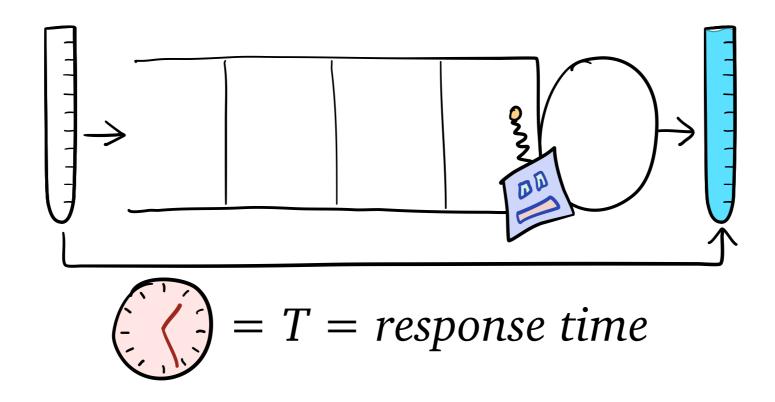
Goal: schedule to minimize two metrics

• *mean* response time **E**[*T*]



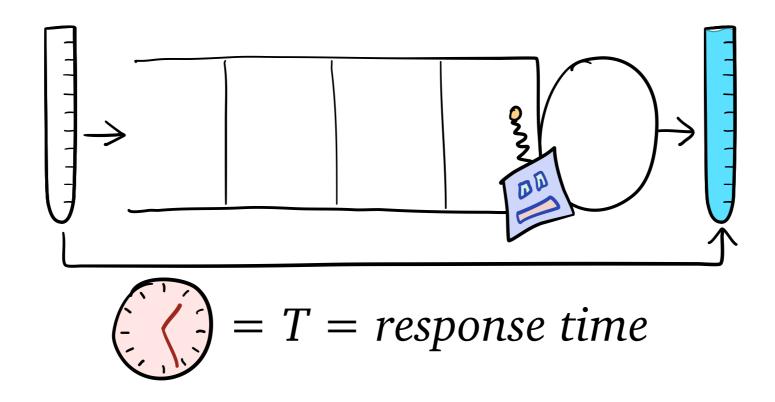
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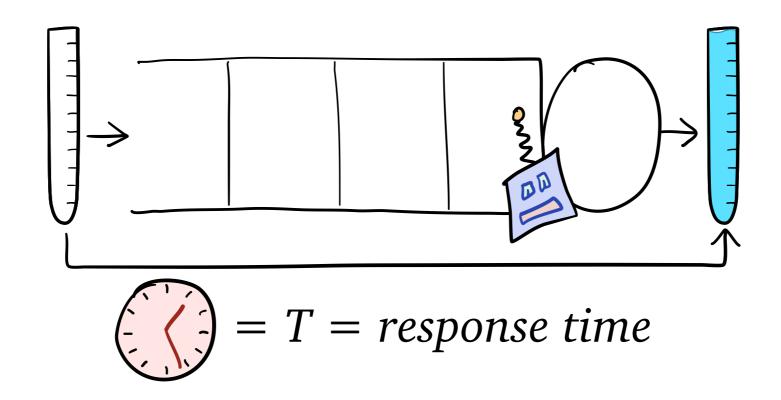
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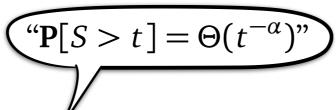
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Setting: *heavy-tailed* job size distribution *S*



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Scheduling with Heavy Tails

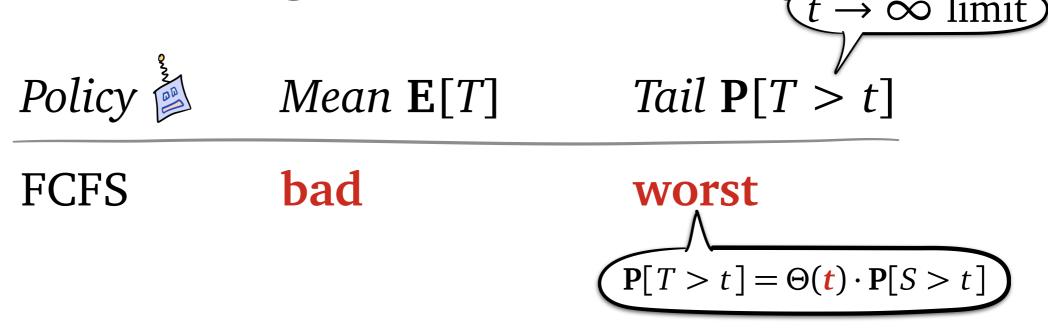
Policy 3

Mean $\mathbf{E}[T]$

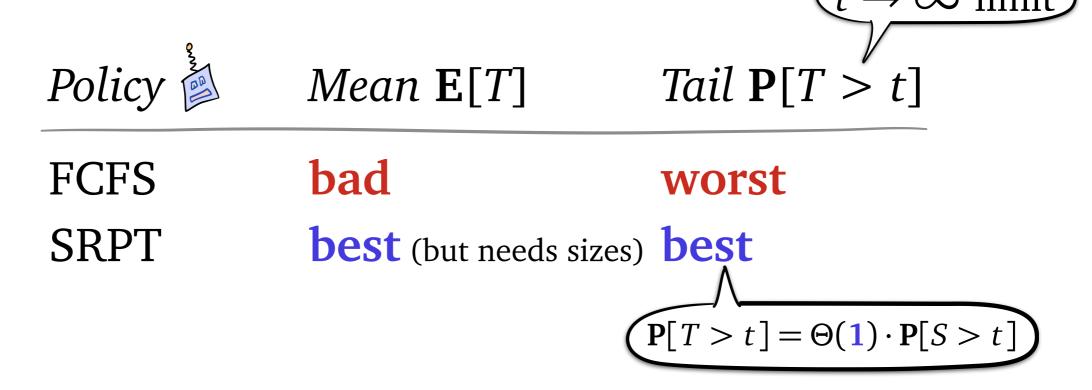
Tail P[T > t]

Policy $\stackrel{\$}{\bowtie}$ Mean $\mathbf{E}[T]$ Tail $\mathbf{P}[T > t]$ FCFS bad worst

Scheduling with Heavy Tails



Policy 🔋	Mean E [T]	Tail $P[T > t]$
FCFS	bad	worst
SRPT	best (but needs sizes)	best



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Policy

Mean $\mathbf{E}[T]$

Tail P[T > t]

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bad

worst

SRPT

best (but needs sizes) **best**

FB

good

best

Gittins

best

M-SERPT

5-approx.

Policy 🖟	Mean $\mathbf{E}[T]$	Tail $P[T > t]$
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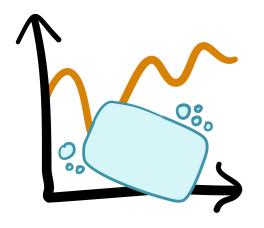
New sufficient condition for rank function to be tail-optimal

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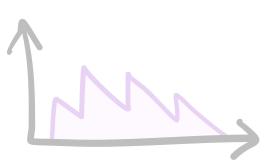
[Scully, van Kreveld, Boxma, Dorsman, & Wierman, SIGMETRICS 2020]

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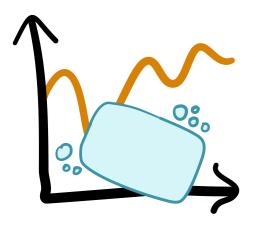


Preemption restricted and/or costly



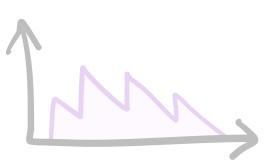
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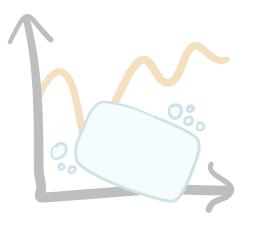
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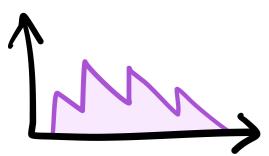


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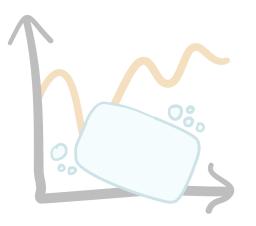
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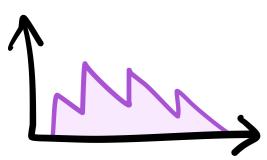


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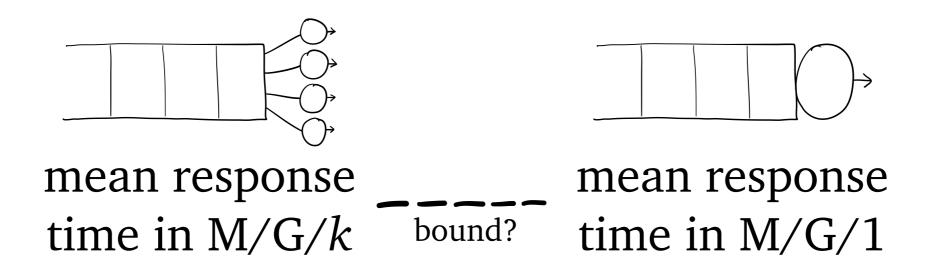


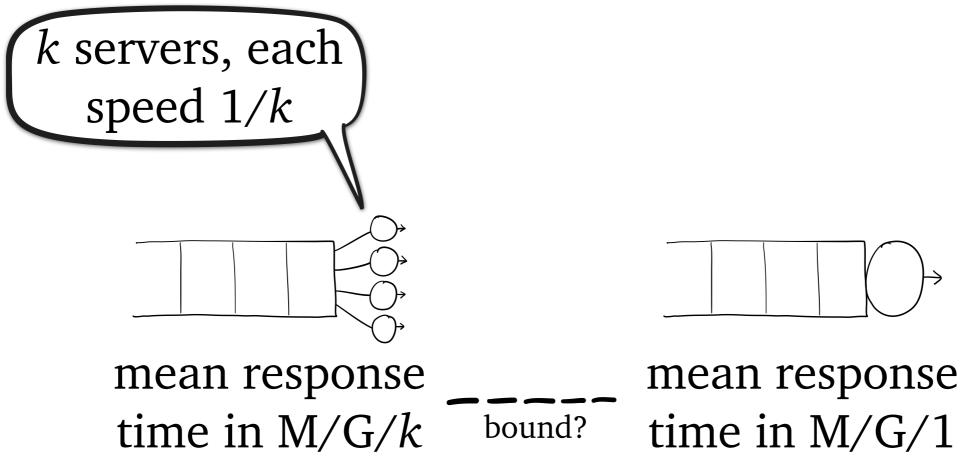
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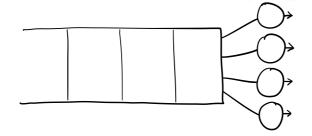


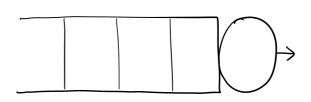




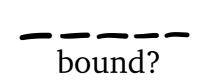
k servers, each speed 1/k

Gittins serves *k* jobs of *k* lowest **ranks**





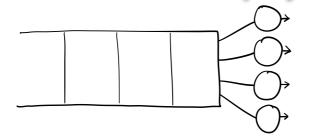
mean response time in M/G/k

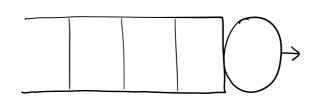


mean response time in M/G/1

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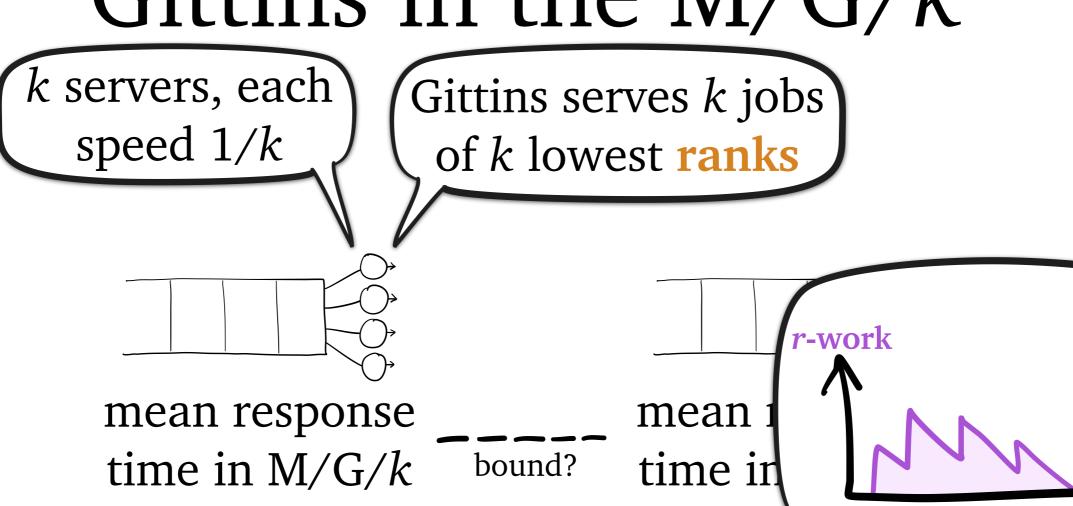
mean response time in M/G/k

bound?

mean response time in M/G/1

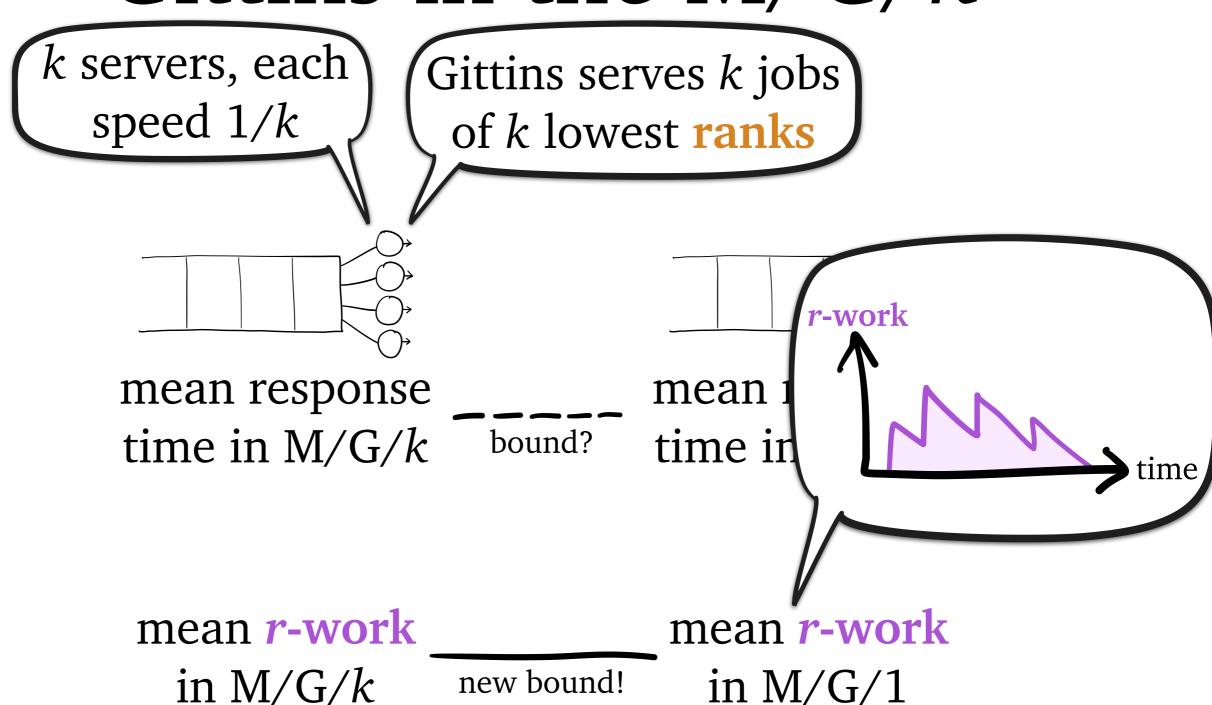
mean *r*-work in M/G/k

mean *r*-work in M/G/1



mean *r*-work in M/G/k

mean *r*-work in M/G/1

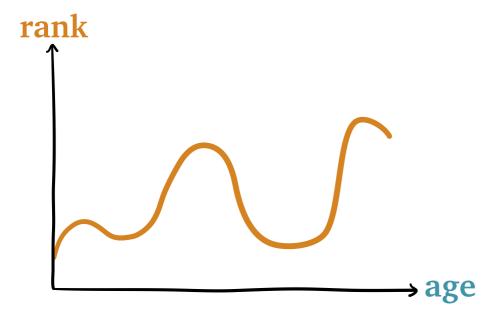


k servers, each Gittins serves *k* jobs speed 1/kof k lowest ranks mean response mean response bound? time in M/G/ktime in M/G/1 new connection! new connection! mean *r*-work mean *r*-work new bound! in M/G/kin M/G/1

r-Work: amount of service a job needs to either

- complete
- or reach Gittins rank $\geq r$

Gittins

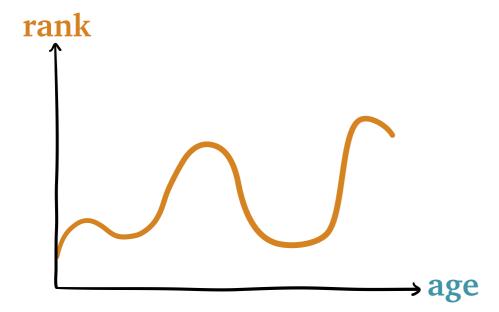


r-Work: amount of service a job needs to either

- complete
- or reach Gittins rank $\geq r$

depends on current age a and size s

Gittins

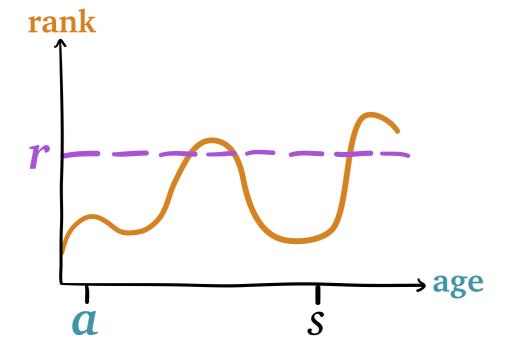


r-Work: amount of service a job needs to either

- complete
- or reach Gittins rank $\geq r$

depends on current age **a** and size s

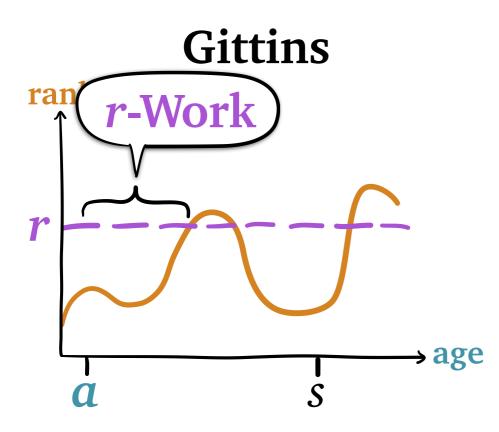
Gittins



r-Work: amount of service a job needs to either

- complete
- or reach Gittins $rank \ge r$

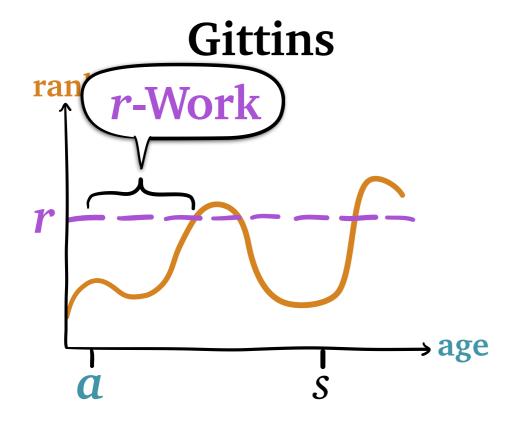
depends on current age **a** and size s

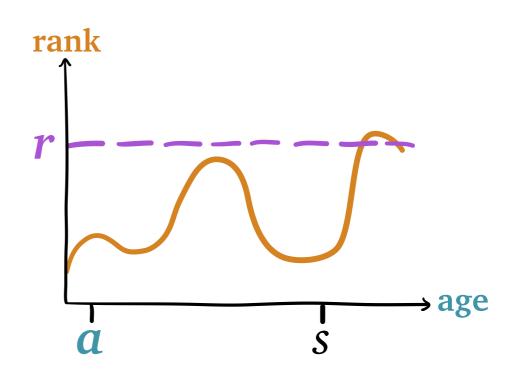


r-Work: amount of service a job needs to either

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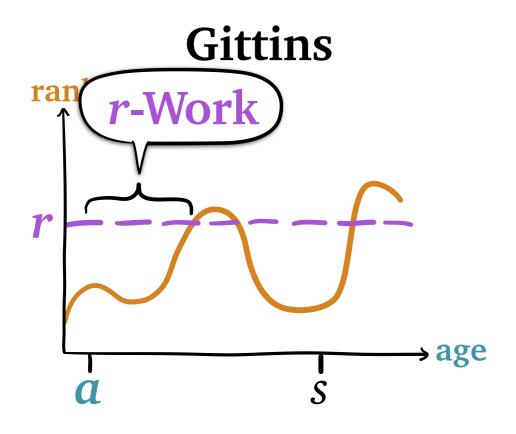


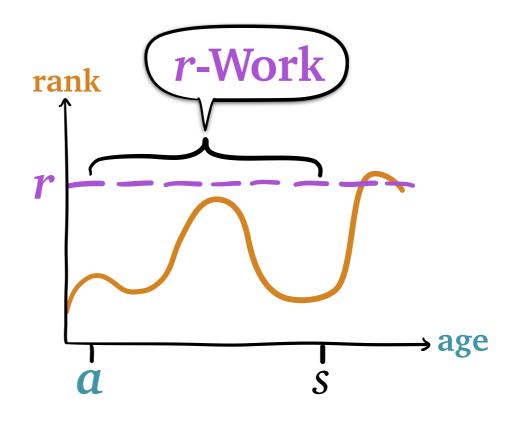
What is *r*-Work?

r-Work: amount of service a job needs to either

- complete
- or reach Gittins rank $\geq r$

depends on current age a and size s



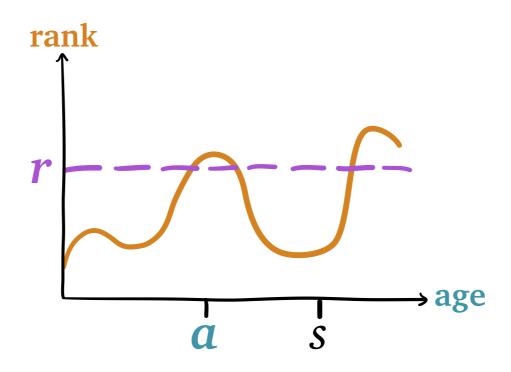


What is *r*-Work?

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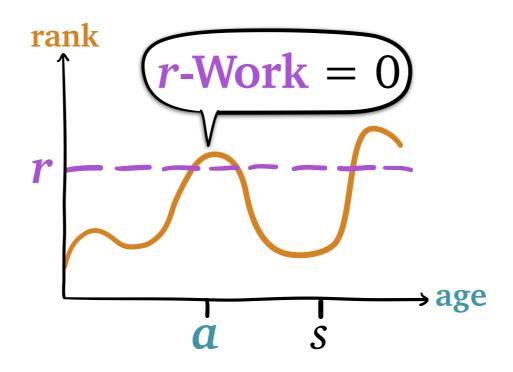


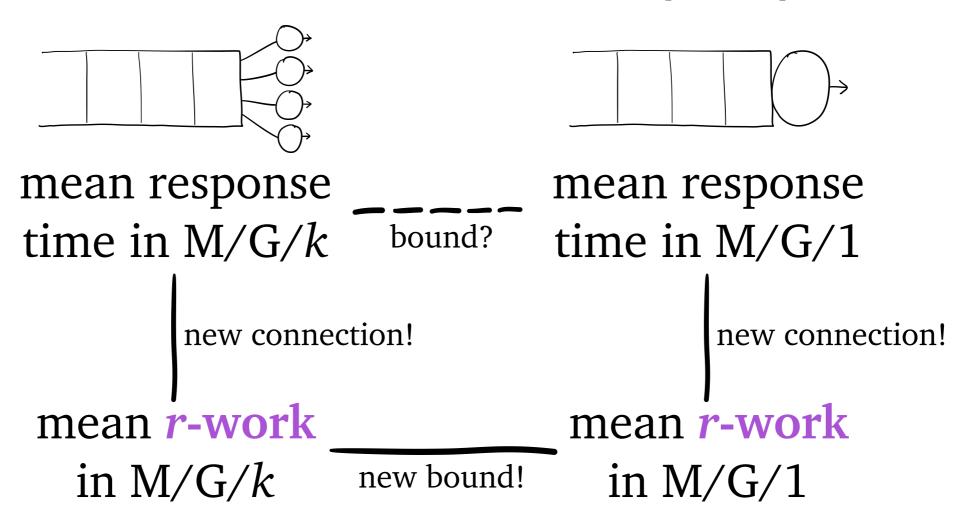
What is *r*-Work?

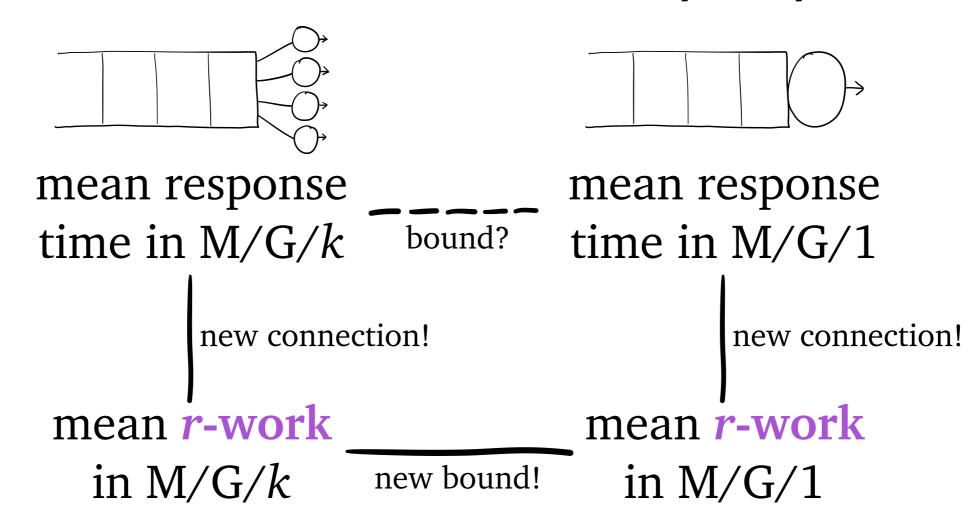
r-Work: amount of service a job needs to either

- complete
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depends on current age a and size s

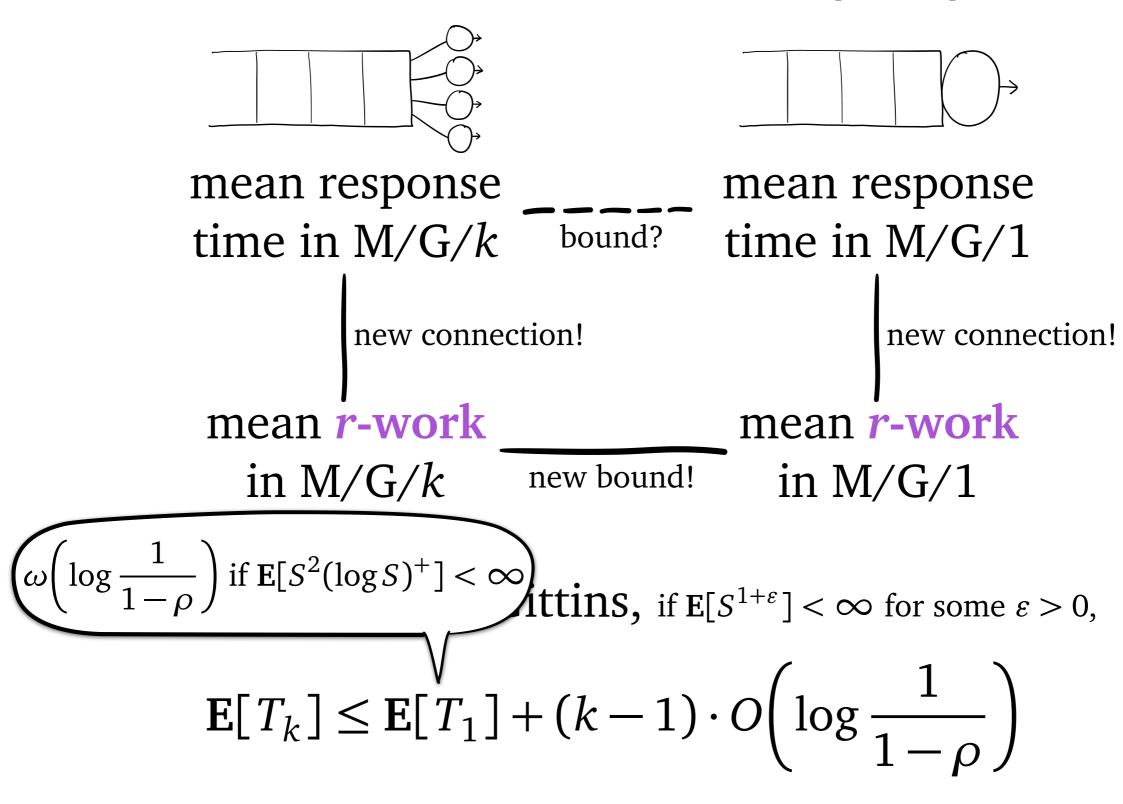


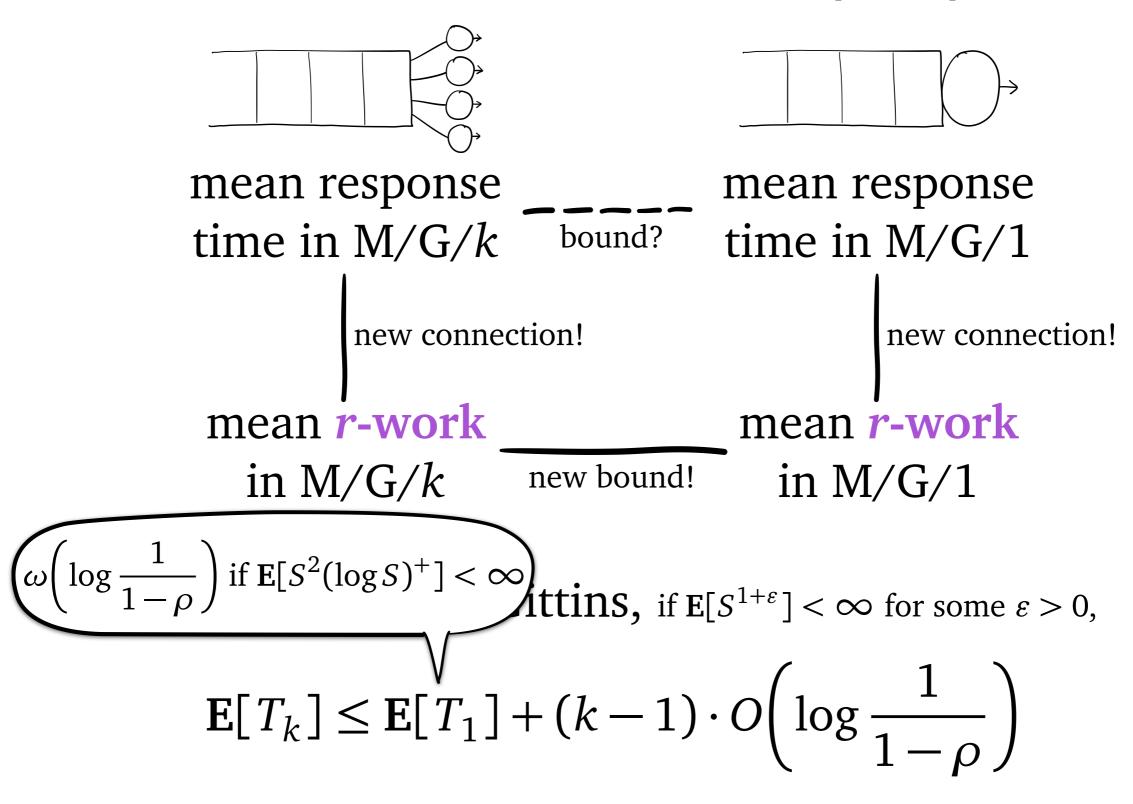




Theorem: under Gittins, if $E[S^{1+\varepsilon}] < \infty$ for some $\varepsilon > 0$,

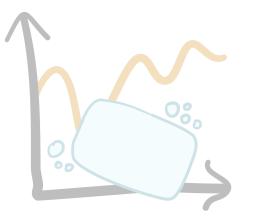
$$\mathbf{E}[T_k] \le \mathbf{E}[T_1] + (k-1) \cdot O\left(\log \frac{1}{1-\rho}\right)$$





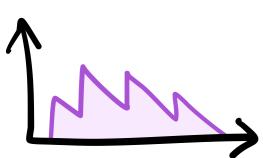
[Scully, Grosof, & Harchol-Balter, SIGMETRICS 2021]

New Tools



SOAP

analyzes a huge variety of scheduling heuristics



r-Work

> provides a new, deeper understanding of Gittins

Goals

Multiple servers



Simple implementation preferred



Preemption restricted and/or costly

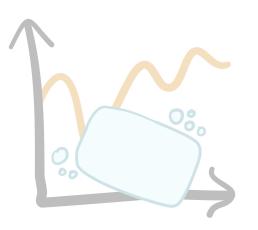


Limited number of priority levels



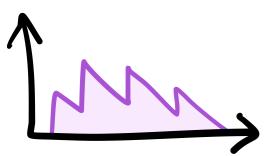
Want to optimize other response time metrics

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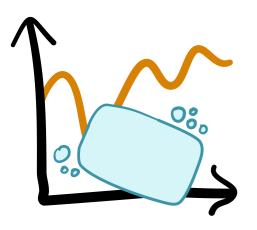


Limited number of priority levels



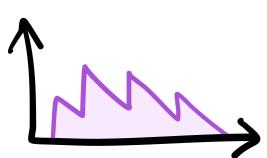
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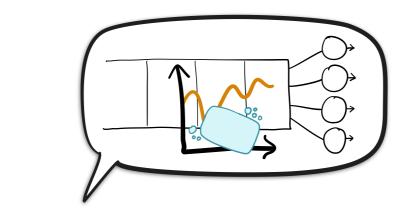
Preemption *restricted* and/or *costly*



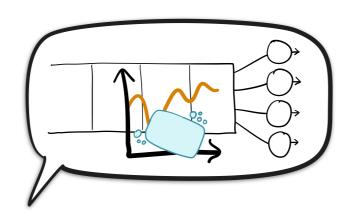
Limited number of priority levels



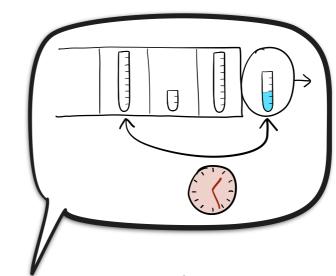
Want to optimize other response time metrics



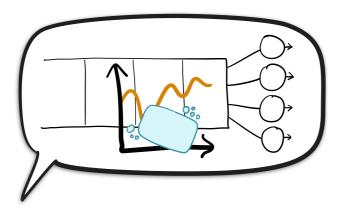
• **SOAP** for M/G/k



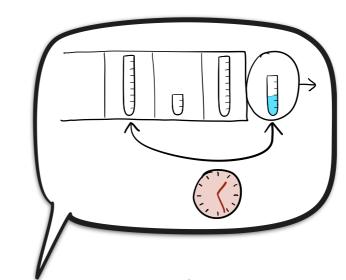




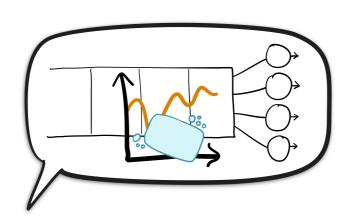
Preemption costs with unrestricted preemption timing



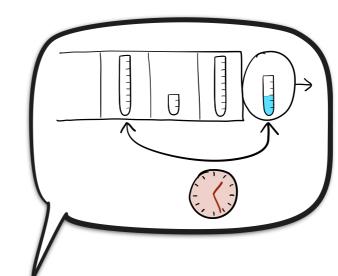




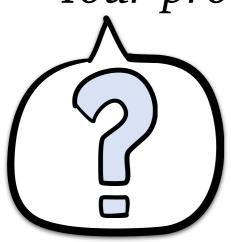
- Preemption costs with unrestricted preemption timing
- Simplifying Gittins for noisy size estimates

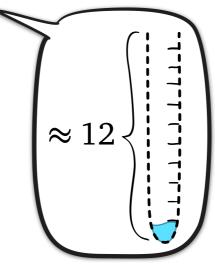






- Preemption costs with unrestricted preemption timing
- Simplifying Gittins for noisy size estimates
- Your problem here!





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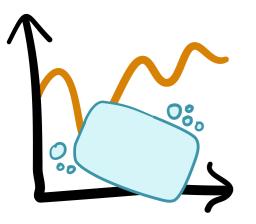
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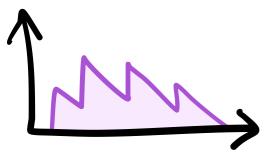
POMACS, 2020. To be presented at SIGMETRICS 2021.

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